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I, KIM MARSHALL, MANAGER EXAMINATION SUPPORT AND SALES hereby certify that annexed is a true copy of the Provisional specification in connection with Application No. PP 5686 for a patent by CANON KABUSHIKI KAISHA filed on 03 September 1998.



WITNESS my hand this
Twenty-seventh day of August 1999

**KIM MARSHALL
MANAGER EXAMINATION SUPPORT
AND SALES**

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PATENT APPLICATION

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

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| In re Application of: |) | |
| GEORGE POLITIS |) | Examiner: NYA |
| Application No.: 09/387,569 |) | Group Art Unit: 2772 |
| Filed: September 1, 1999 |) | |
| For: REGION BASED IMAGE |) | |
| COMPOSITING |) | November 9, 1999 |

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Assistant Commissioner for Patents
Washington, D.C. 20231

CLAIM TO PRIORITY

Sir:

Applicant hereby claims priority under the International Convention and all rights to which he is entitled under 35 U.S.C. § 119 based upon the following Australian Priority Application:

✓ PP 5686, filed September 3, 1998.

A certified copy of the priority document is enclosed.

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Applicant's undersigned attorney may be reached in our New York office by telephone at (212) 218-2100. All correspondence should continue to be directed to our address given below.

Respectfully submitted,


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AUSTRALIA

Patents Act 1990

PROVISIONAL SPECIFICATION FOR THE INVENTION ENTITLED:

Region-Based Image Compositing

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of Applicant:

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This invention is best described in the following statement:



Region-Based Image Compositing

Field of the Invention

The present invention relates to the creation of computer-generated images both in the form of still pictures and video imagery, and, in particular, relates to efficient process, apparatus, and system for creating an image made up by compositing multiple components.

Background

Computer generated images are typically made up of many differing components or graphical elements which are rendered and composited together to create a final image. In recent times, an "opacity channel" (also known as a "matte", an "alpha channel", or simply "opacity") has been commonly used. The opacity channel contains information regarding the transparent nature of each element. The opacity channel is stored alongside each instance of a colour, so that, for example, a pixel-based image with opacity stores an opacity value as part of the representation of each pixel. An element without explicit opacity channel information is typically understood to be fully opaque within some defined bounds of the element, and assumed to be completely transparent outside those bounds.

An expression tree offers a systematic means for representating an image in terms of its constituent elements and which facilitates later rendering. Expression trees typically comprise a plurality of nodes including leaf nodes, unary nodes and binary nodes. Nodes of higher degree, or of alternative definition may also be used. A leaf node, being the outer most node of an expression tree, has no descendent nodes and represents a primitive constituent of an image. Unary nodes represent an operation which modifies the pixel data coming out of the part of the tree below the unary operator. Unary nodes include such operations as colour conversions, convolutions (blurring etc) and operations such as red-eye removal. A binary node typically branches to left and right subtrees, wherein each subtree is itself an expression tree comprising at least one leaf node. Binary nodes represent an operation which combines the pixel data of its two children to form a single result. For example, a binary node may be one of the standard "compositing operators" such as OVER, IN, OUT, ATOP and alpha-XOR, examples of which and other are seen in Fig. 20.

Several of the above types of nodes may be combined to form a compositing tree. An

example of this is shown in Fig. 1. The result of the left-hand side of the compositing tree may be interpreted as a colour converted image being clipped to spline boundaries. This construct is composited with a second image.

Although the non-transparent area of a graphical element may of itself be of a certain size, it need not be entirely visible in a final image, or only a portion of the element may have an effect on the final image. For example, assume an image of a certain size is to be displayed on a display. If the image is positioned so that only the top left corner of the image is displayed by the display device, the remainder of the image is not displayed. The final image as displayed on the display device thus comprises the visible portion of the image, and the invisible portion in such a case need not be rendered.

Another way in which only a portion of an element may have an effect is when the portion is obscured by another element. For example, a final image to be displayed (or rendered) may comprise one or more opaque graphical elements, some of which obscure other graphical elements. Hence, the obscured elements have no effect on the final image.

A conventional compositing model considers each node to be conceptually infinite in extent. Therefore, to construct the final image, a conventional system would apply a compositing equation at every pixel of the output image. Interactive frame rates of the order greater than 15 frames per second can be achieved by relatively brute-force approaches in most current systems, because the actual pixel operations are quite simple and can be highly optimised. This highly optimised code is fast enough to produce acceptable frame rates without requiring complex code. However, this is certainly not true in a compositing environment.

The per-pixel cost of compositing is quite high. This is because typically an image is rendered in 24-bit colour in addition to an 8-bit alpha channel, thus giving 32 bits per pixel. Each compositing operator has to deal with each of the four channels. Therefore, the approach of completely generating every pixel of every required frame when needed is inefficient, because the per-pixel cost is too high.

Problems arise with prior art methods when rendering graphical objects which include transparent and partially-transparent areas. Further, such methods typically do not handle the full range of compositing operators.

Summary of the Invention

It is an object of the present invention to substantially overcome, or ameliorate, one or more of the deficiencies of the above mentioned methods by the provision of a method for creating an image made up by compositing multiple components.

5 According to one aspect of the present invention there is provided a method of creating an image, said image to be formed by rendering and compositing at least a plurality of graphical objects, each said object having a predetermined outline, said method comprising the steps of:

10 dividing a space in which said outlines are defined into a plurality regions, each said region being defined by at least one region outline substantially following at least one of said predetermined outlines or parts thereof and being substantially formed by segments of a virtual grid encompassing said space;

manipulating said regions to determine a plurality of further regions, wherein each said further region has a corresponding compositing expression;

15 classifying said further regions according to at least one attribute of said graphical objects within said further regions;

modifying each said corresponding compositing expression according to a classification of each said further region to form an augmented compositing expression for each said further region;

20 compositing said image using each of said augmented compositing expressions.

According to another aspect of the present invention there is provided a method of creating an image, said image to be formed by rendering and compositing at least a plurality of graphical objects, each said object having a predetermined outline, said method comprising the steps of:

25 dividing a space in which said outlines are defined into a plurality regions, each said region being defined by at least one region outline substantially following at least one of said predetermined outlines or parts thereof and being substantially formed by segments of a virtual grid encompassing said space, wherein each object has two region outlines arranged either side of said predetermined outline to thus define three regions for each said
30 object, and wherein each said region has a corresponding compositing expression;

classifying said regions according to at least one attribute of said graphical objects within said regions;

modifying each said corresponding compositing expression according to a classification of each said region to form an augmented compositing expression for each said region;

compositing said image using each of said augmented compositing expressions.

According to still another aspect of the present invention there is provided a method of creating an image, said image being formed by rendering at least a plurality of graphical objects to be composited according to a compositing expression, each said object having a predetermined outline, said method comprising the steps of:

dividing a space in which said outlines are defined into a plurality of mutually exclusive regions wherein each of said regions is defined by a region outline substantially following at least one of said predetermined outlines or parts thereof;

examining each said region to determine those said objects which contribute to said region;

modifying said compositing expression on the basis of the contribution of each of said objects within said region to form an optimized compositing expression for each said region; and

compositing said image using each of said optimized compositing expressions.

According to still another aspect of the present invention there is provided a method of creating an image, said image being formed by rendering at least a plurality of graphical objects to be composited according to a compositing expression, each said object having a predetermined outline, said method comprising the steps of:

dividing a space in which said outlines are defined into a plurality of mutually exclusive regions;

examining each said region to determine those said objects which contribute to said region;

modifying said compositing expression on the basis of the contribution of each of said objects within said region; and

compositing said image using said modified compositing expression.

According to still another aspect of the present invention there is provided a method of creating an image, said image comprising a plurality of graphical objects to be composited according to a compositing expression, said method comprising the steps of:

dividing a space in which said graphical objects are defined into a plurality of regions;

examining each said region to determine those said objects which contribute to said region;

modifying said compositing expression on the basis of said examination; and

compositing said image using said modified compositing expression.

According to still another aspect of the present invention there is provided a method of creating a series of images, each member of said series being related to a preceding member, said images being formed by rendering at least a plurality of graphical objects to be composited according to a hierarchical structure representing a compositing expression, said hierarchical structure including a plurality of nodes each representing a component of said image, each of said objects having a predetermined outline, said method comprising the steps of:

(a) for each said node:

(i) dividing a component image space in which said outlines are defined into at least one mutually exclusive region, each said region being related to at least one graphical object;

(ii) examining each said region to determine those objects that contribute to the region;

(b) creating an internodal dependency record for each of said regions identifying those said regions that will be affected by a change in any one of said regions;

(c) rendering a first image of said series by compositing all regions of said hierarchical structure;

(d) in response to at least one change to at least one of said nodes;

(i) examining said internodal dependency record to identify those of said regions affected by said at least one change;

(ii) for an affected node, updating the corresponding identified regions and

incorporating into said node those (any) new regions arising from the change;

(iii) updating said internodal dependency record to reflect changes to said hierarchical structure;

(iv) rendering a further image of said series by compositing (only) those regions affected by said at least one change; and

(e) repeating step (d) for further changes to at least one of said nodes.

According to still another aspect of the present invention there is provided a method of creating a series of images, said images being formed by rendering at least a plurality of graphical objects to be composited according to a hierarchical structure, said hierarchical structure including a plurality of nodes each representing a component of said image, each of said objects having a predetermined outline, said method comprising the steps of:

(a) for each said node:

(iii) dividing a space in which said outlines are defined into at least one mutually exclusive region;

(iv) examining each said region to determine those objects that contribute to the region;

(b) creating an internodal dependency record for each of said regions based on said examination;

(c) rendering a first image of said series utilising said hierarchical structure; and then, in response to at least one change to at least one of said nodes;

(d) examining said internodal dependency record;

(i) for an affected node, updating the corresponding regions;

(ii) updating said internodal dependency record;

(iii) rendering a further image of said series by compositing those regions affected by said at least one change; and

(e) repeating step (d) for further changes to at least one of said nodes.

According to still another aspect of the present invention there is provided a method of creating a series of images, said images being formed by rendering at least a plurality of graphical objects to be composited according to a hierarchical structure, said hierarchical structure including a plurality of nodes each representing a component of said image, said

method comprising the steps of:

(a) for each said node:

(i) dividing a component image space in which said graphical objects are defined into at least one region;

5 (ii) examining each said region;

(b) creating an internodal dependency record for each of said regions;

(c) rendering a first image of said series utilising said hierarchical structure;

and then, in response to at least one change to at least one of said nodes;

(d) examining said internodal dependency record;

10 (i) for an affected node, updating the corresponding regions;

(ii) updating said internodal dependency record;

(iii) rendering a further image of said series; and

(e) repeating step (d) for further changes to at least one of said nodes.

15 Brief Description of the Drawings

A preferred embodiment of the present invention will now be described with reference to the following drawings:

Fig. 1 is an example of a compositing tree;

20 Fig. 2 illustrates an image containing a number of overlapping objects and the corresponding compositing tree;

Fig. 3 shows the image of Fig. 2 illustrating the different regions which exist in the image and listing the compositing expression which would be used to generate the pixel data for each region;

25 Fig. 4 is the image of Fig. 3, illustrating the compositing operations after being optimised according to one example of the preferred embodiment;

Fig. 5 illustrates the result of combining two region descriptions using the Union operation according to the preferred embodiment;

Fig. 6 illustrates the result of combining two region descriptions using the Intersection operation according to the preferred embodiment;

30 Fig. 7 illustrates the result of combining two region descriptions using the Differ-

ence operation according to the preferred embodiment;

Figs. 8A to 8D illustrate the steps involved in combining two region groups using the Over operation according to the present invention;

Fig. 9 illustrates an image and compositing tree according to another example of the preferred embodiment;

Fig. 10 illustrates an image and compositing tree according to still another example of the preferred embodiment;

Fig. 11 illustrates the effect on the image of Fig. 10 of moving region A;

Fig. 12 illustrates an image and compositing tree according to still another example of the preferred embodiment;

Fig. 13 illustrates the effect on the image of Fig. 12 of moving region A;

Fig. 14 illustrates the effect on the image of Fig. 12 of moving region B; and

Fig. 15 illustrates those nodes in a compositing tree which need to have their region groups updated if leaf nodes B and H change;

Fig. 16 illustrates a region and its x and y co-ordinates;

Fig. 17 illustrates two regions and their x and y co-ordinates;

Fig. 18 illustrates an image and compositing tree according to still another example of the preferred embodiment;

Fig. 19 illustrates an apparatus upon which the preferred embodiment is implemented;

Fig. 20 depicts the result of a variety of compositing operators useful with the present invention;

Fig. 21 illustrates regions formed by combining two circles with non-grid-aligned regions;

Fig. 22 illustrates improved regions formed by combining two circles with grid-aligned regions; and

Appendix 1 is a listing of source code for implementing the preferred embodiment.

Detailed Description

1.0 Underlying Principles

The basic shape of operands to compositing operators in most current systems is the

rectangle, regardless of the actual shape of the object being composited. It is extremely easy to write an operator which composites within the intersection area of two bounding boxes. However, as a bounding box typically does not accurately represent the actual bounds of a graphical object, this method results in a lot of unnecessary compositing of completely transparent pixels over completely transparent pixels. Furthermore, when the typical make-up of a composition is examined, it can be noticed that areas of many of the objects are completely opaque. This opaqueness can be exploited during the compositing operation. However, these areas of complete opaqueness are usually non-rectangular and so are difficult to exploit using compositing arguments described by bounding boxes. If irregular regions are used for this purpose, then these regions could then be combined in some way to determine where compositing should occur. Furthermore, if any such region is known to be fully transparent or fully opaque, further optimisations are possible.

Most current systems fail to exploit similarities in composition between one frame and the next. It is rare for everything to change from frame to frame and therefore large areas of a compositing tree will remain unchanged. An example of this is where a cartoon type character comprising multiple graphical objects is rendered on a display. If, for example, the character spilt some paint on its shirt in the next frame, then it is not necessary to render the entire image again. For example, the head and legs of the character may remain the same. It is only necessary to render those components of the image that have been altered by the action. In this instance, the part of the shirt on which the paint has been spilt may be re-rendered to be the same colour as the paint, whilst the remainder of the character stays the same. Exploiting this principle may provide large efficiency improvements. If incremental changes are made to the compositing tree, then only a reduced amount of updating is necessary to affect the change.

Many current graphical systems use what is known as an *immediate mode* application program interface (API). This means that for each frame to be rendered, the complete set of rendering commands is sent to the API. This is somewhat inefficient in a compositing environment, as typically, large sections of the compositing tree will be unchanged from one frame to the next, but would be completely re-rendered anyway in immediate mode. The preferred embodiment, on the other hand, is considered by the present inventors to be

best described as a *retained mode* API. This means that instead of providing the complete compositing tree on a per-frame basis, the user provides an initial compositing tree, and then modifies it on a per-frame basis to effect change. Changes which can be made to the tree include geometrically transforming part or all of the tree, modifying the tree structure (unlinking and linking subtrees), and modifying attributes (eg: color) of individual nodes. Note that such modifications may not necessarily mean that the tree structure, for example as seen in Fig. 1, will change where only the attributes of an individual node have been modified.

The rendering operation of the preferred embodiment is a combination of a number of techniques and assumptions which combine to provide high quality images and high frame rates. Some of the contributing principles are:

(i) The use of irregular regions to minimise per-pixel compositing. For example, if one graphical object is on top of another, then pixel compositing is only needed inside the area where the two objects intersect. Having the ability to use irregular regions gives the ability to narrow down areas of interest much more accurately.

(ii) An assumption is made that in the transition from one frame to the next, only part of the tree will change. This can be exploited by caching away expensive-to-generate information regarding the composition so that it can be re-used from one frame to the next. Examples of expensive-to-generate information are - regions of interest (boundaries of areas of intersection between objects etc); pixel data (representing expensive composites etc); and topological relationships between objects.

(iii) If an opaque object is composited with another object using the OVER operator, then it completely obscures what it is composited onto (inside the opaque objects area). This is a very useful property because it means that no expensive pixel compositing is required to achieve the output pixel within the area of overlap. (The pixel value is the same as that at the equivalent spot on the opaque object). Opaque objects induce similar behaviour in most of the compositing operators. Therefore, the preferred embodiment attempts to exploit opaque areas as much as possible.

2.0 Basic Static Rendering

Static Rendering deals with the problem of generating a single image from a compos-

iting tree as quickly as possible. Some of the pixel compositing methods of the preferred embodiment will be explained using a static rendering example.

5 An example of a simple compositing tree which consists of leaf node objects and only using the "OVER" operator is shown in Fig. 2. Conventionally, each node is considered to be conceptually infinite in extent. One method to construct the final image is to apply the compositing equation $((D \text{ OVER } B) \text{ OVER } C) \text{ OVER } (A \text{ OVER } E)$ at every pixel of the output image. However, this is quite an inefficient method.

10 A composition can generally be subdivided into a number of mutually exclusive irregular regions. The above compositing expression may be simplified independently within each region. In the example of Fig. 2, A, C and E represent opaque objects. B and D, on the other hand are partially transparent. Fig. 3 shows the different regions (1-10) produced using the five objects which exist in the example, and the compositing expression which would be used to generate the pixel data for each specific region.

15 The compositing expressions provided in Fig. 3 make no attempt to exploit the properties of the object's opacity. If these properties are used to simplify the compositing expressions for each region, the expressions of Fig. 4 are obtained. This results in a simplification of the rendering of regions 2, 3, 5, 6, 7, 8 and 9 compared with Fig. 3. These simplified compositing expressions would result in far fewer pixel compositing operations being performed to produce the final picture.

20 Fig. 4 represents the region subdivision for the root of the compositing tree. However, every node in the compositing tree can itself be considered the root of a complete compositing tree. Therefore, every node in the compositing tree can have associated with it a group of regions which together represent the region subdivision of the subtree of which it is the root. Region subdivision provides a convenient means of managing the complexity of a compositing tree and an efficient framework for caching expensive data.

25 Using the principles noted above, a compositing expression may be simplified dependent upon whether the graphical objects being composited are wholly opaque, wholly transparent or otherwise (herewith deemed "ordinary").

30 Table 1 shows how the compositing operations of Fig. 20 can be simplified when one or both operands are opaque or transparent.

TABLE 1

| Expression | A's opacity | B's opacity | Optimised |
|------------|-------------|-------------|-----------|
| AoverB | Transparent | Transparent | neither |
| | Transparent | Ordinary | B |
| | Transparent | Opaque | B |
| | Ordinary | Transparent | A |
| | Ordinary | Ordinary | AoverB |
| | Ordinary | Opaque | AoverB |
| | Opaque | Transparent | A |
| | Opaque | Ordinary | A |
| | Opaque | Opaque | A |
| | Transparent | Transparent | neither |
| AroverB | Transparent | Ordinary | B |
| | Transparent | Opaque | B |
| | Transparent | Transparent | A |
| | Ordinary | Ordinary | BoverA |
| | Ordinary | Opaque | B |
| | Opaque | Transparent | A |
| | Opaque | Ordinary | BoverA |
| | Opaque | Opaque | B |
| | Transparent | Transparent | neither |
| | Transparent | Ordinary | neither |
| AinB | Transparent | Opaque | neither |
| | Transparent | Transparent | neither |
| | Ordinary | Ordinary | AinB |
| | Ordinary | Opaque | A |
| | Opaque | Transparent | neither |
| | Opaque | Ordinary | AinB |
| | Opaque | Opaque | A |
| | Transparent | Transparent | neither |
| | Transparent | Ordinary | neither |
| | Transparent | Opaque | neither |
| AinB | Ordinary | Transparent | neither |
| | Ordinary | Ordinary | BinA |
| | Ordinary | Opaque | BinA |
| | Opaque | Transparent | neither |
| | Opaque | Ordinary | B |
| | Opaque | Opaque | B |
| | Transparent | Transparent | neither |
| | Transparent | Ordinary | neither |
| | Transparent | Opaque | neither |
| | Ordinary | Ordinary | neither |
| AoutB | Ordinary | Opaque | neither |
| | Opaque | Transparent | neither |
| | Opaque | Ordinary | neither |
| | Opaque | Opaque | A |
| | Ordinary | Transparent | AoutB |
| | Ordinary | Ordinary | neither |
| | Ordinary | Opaque | A |
| | Opaque | Transparent | AoutB |
| | Opaque | Ordinary | neither |
| | Opaque | Opaque | neither |
| ArouB | Transparent | Transparent | neither |
| | Transparent | Ordinary | B |
| | Transparent | Opaque | B |
| | Ordinary | Transparent | neither |
| | Ordinary | Ordinary | BoutA |
| | Ordinary | Opaque | BoutA |
| | Opaque | Transparent | neither |
| | Opaque | Ordinary | neither |
| | Opaque | Opaque | neither |
| | Transparent | Transparent | neither |
| AatopB | Transparent | Ordinary | B |
| | Transparent | Ordinary | B |

| | | | |
|---------|-------------|-------------|---------|
| | Transparent | Opaque | B |
| | Ordinary | Transparent | neither |
| | Ordinary | Ordinary | AatopB |
| | Ordinary | Opaque | AatopB |
| | Opaque | Transparent | neither |
| | Opaque | Ordinary | AatopB |
| | Opaque | Opaque | A |
| AratopB | Transparent | Transparent | neither |
| | Transparent | Ordinary | neither |
| | Transparent | Opaque | neither |
| | Ordinary | Transparent | A |
| | Ordinary | Ordinary | BatopA |
| | Ordinary | Opaque | BatopA |
| | Opaque | Transparent | A |
| | Opaque | Ordinary | BatopA |
| | Opaque | Opaque | B |
| AxorB | Transparent | Transparent | neither |
| | Transparent | Ordinary | B |
| | Transparent | Opaque | B |
| | Ordinary | Transparent | A |
| | Ordinary | Ordinary | AxorB |
| | Ordinary | Opaque | AxorB |
| | Opaque | Transparent | A |
| | Opaque | Ordinary | AxorB |
| | Opaque | Opaque | neither |

2.1 Basic Data Model

Associated with every node in a compositing tree is a group of mutually exclusive regions which together represent the non-transparent area of the node. It should be noted that the region descriptions that the preferred embodiment uses are generally not pixel accurate. A region may in fact contain some transparent pixels. However, any point lying outside of all the regions at a node is certain to be transparent. The set of these regions at a node is known as a *region group*. A leaf node region group may contain only one or two regions. The region group at the root of the tree may contain hundreds of regions. Each region in a region group contains the following basic data:

(i) **A Region Description** is a low-level representation of the boundaries of the region. The region descriptions of all the regions in a region group must be mutually exclusive (non-intersecting). However, the preferred embodiment is not limited to using axis-parallel (ie: every side parallel or perpendicular to a scan line of an output device) region descriptions. The preferred embodiment allows region descriptions which more closely represent arbitrary shaped regions.

(ii) **A Proxy** is some means of caching the pixel data resulting from applying the operations specified by the compositing expression at every pixel inside the region

description. A proxy can be as simple as a 24-bit colour bitmap, or something much more complicated (such as a run-length encoded description). Fundamentally, a proxy simply has to represent pixel data in some way which makes it efficient to retrieve and use.

Every region group also contains a region description which is the union of all the region descriptions of the regions in the region group. It essentially represents the entire coverage of the region group.

2.2 Region Descriptions and Region Arithmetic

The region arithmetic and data structure of the preferred embodiment has the following properties:

-to allow the representation of *complex* regions, including convex regions, concave regions and regions with holes. This is necessary so that a region will be reasonably able to follow the geometry of the graphic object it represents;

-is space efficient. In a complicated composition there will be many regions. For memory efficiency, it is therefore preferable that the cost of storing these regions is reasonably

small;

-it should support basic set operations - Union, Intersection and Difference;

-the above-noted basic operations should be efficient in terms of speed. In a complex compositing tree, it is possible that a large amount of region arithmetic will be undertaken.

A poor implementation of region arithmetic could lead to the time taken by region arithmetic being greater than the time saved from the reduction in per-pixel compositing;

-it is advantageous if the region description can be geometrically translated efficiently. In cases where a graphic object is translated, it's associated regions can then be translated quickly; and

-it is sometimes helpful to be able to quickly compare two regions to determine if they are the same. It is not necessary to obtain any other statistics on their similarity, simple equality is all that is required.

Two conventional region description techniques were considered and rejected for the preferred embodiment. These were-

Polygons: A polygon can be used to represent almost any object, the disadvantage of using a polygon, however, is that its generality makes implementing the set

operations slow and inefficient.

Quadtrees: Using quadtrees, set operations are easy to implement and are quite efficient. In addition, they can represent a wide variety of regions given sufficient granularity (all edges in a quadtree have to be axis-parallel). Their major failing is that all quadtrees must be aligned on the same grid (granularity). This means that it is impossible to simply translate a quadtree by an arbitrary amount. Unless that amount is a multiple of the underlying grid size, the quadtree will need to be recalculated from the object it describes (otherwise it will keep growing). Therefore, quadtrees are not suitable in application domains where geometric translation is a frequent operation.

The region description data structure of the preferred embodiment can be understood by imagining that along a vertical line every coordinate has a state which is one of either inside or outside the region. The data structure stores those y co-ordinates at which some change of state between inside and outside occurs. For each such y co-ordinate, the data contains spans of coordinates each of which toggles the state every vertical line running through it. Each span of x co-ordinates is called a run. The sequence of runs associated with a y co-ordinate is called a row. For example, the region of Fig. 16 could be described by the following:

row y = 10 : x = 10, x = 100

row y = 100 : x = 10, x = 100

Similarly, the regions of Fig. 17 could be described by the following:

row y = 10 : x = 10, x = 100

row y = 30 : x = 30, x = 70

row y = 70 : x = 30, x = 70

row y = 100 : x = 10, x = 100

The data representing a region is represented by an array of integer values. There are two "special" values -

R_NEXT_IS_Y A beginning-of-row marker. Indicates that the next integer in the sequence will represent a y coordinate.

R_EOR Stands for End-of-Region. Indicates that the region description has finished.

All other values represent x or y coordinates. The x coordinates in a row represent runs. The first two co-ordinates represent a run, then the next two represent the next run and so on. Therefore, the x coordinates in a row should always be increasing. Also, there should always be an even number of x-coordinates in a row. The region data stream for
5 Fig. 17 is shown below.

```

R_NEXT_IS_Y 10 10 100
R_NEXT_IS_Y 30 30 70
R_NEXT_IS_Y 70 30 70
R_NEXT_IS_Y 100 10 100
10 R_EOR
```

The preferred embodiment also contains the bounding box of the region, as this is useful in certain set operations.

As seen in Fig. 6, if two region descriptions are combined using a Union operation, then the resultant region description will describe an area in which either region descrip-
15 tion is active.

As seen in Fig. 7, if two region descriptions are combined using the Intersection operation, then the resultant region description will describe an area in which both the region descriptions are active.

If two region descriptions are combined using the Difference operation, then the
20 resultant region will describe an area in which *only the first region is active*, as seen in Fig. 8.

<<< Implementation of Region operations is illustrated in the source code we will provide you. Presumably it will become an appendix>>>

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2.3 Constructing Region Groups:

2.3.1 Constructing Leaf Node Region Groups

A region group for a leaf node will typically contain one or more regions, which together fully contain the non-transparent area of the graphical object represented by the
30 leaf node. Typically, the non-transparent area is divided into regions where each region

has some property that facilitates optimization. For example, the non-transparent area of some graphical object may be divided into two regions, one fully opaque and the other with ordinary opacity. The above mentioned compositing optimizations would apply where the opaque region is composited.

Alternatively, the leaf node could be subdivided based on some other attribute. For example, it could be divided into two regions, one representing an area of constant colour, the other representing blended colour. Areas of constant colour may be composited more efficiently than areas with more general colour description.

2.3.1.1 Region Formation and Phasing

When creating regions, it is not always beneficial that region boundaries follow graphical object boundaries precisely. What is important is that any property that facilitates optimization is valid at all points within a region said to have that property. For example, an opaque circle could be covered exactly by one circular region which is classified as opaque, or by two approximate regions, one fully opaque octagonal region inscribed in the circle, and one annular octagonal region of ordinary opacity that includes the remainder of the circle plus some area exterior to the circle.

There is typically a trade-off between how closely region boundaries follow graphical object boundaries and the benefits obtained. If region boundaries follow object boundaries very closely, a lot of work is usually involved in creating the region boundaries and in performing intersections and differences of regions (the reasons for needing to perform such operations are explained in later sections). However, if region boundaries are too approximate, they may either include large areas that are outside the objects' boundaries, resulting in too much unnecessary compositing, or they may fail to include large areas where known properties lead to optimization.

One approach, as illustrated in the appendix, is to limit region boundaries to sequences of horizontal and vertical segments. Using this approach, the typical segment size is chosen so that there is neither too much detail so that the region operations are overburdened, nor too much approximation to result in wasted compositing or insufficient optimization.

One method to improve the efficiency of region operations is to choose as many as is practical of the horizontal and vertical segments of substantially all region boundaries to be in phase. In other words, they are to be chosen from the horizontal and vertical lines of the same grid. The grid need not be regularly spaced, nor have the same spacing horizon-

tally and vertically, although typically it will.

This method improves the efficiency of region operations by seeking to keep all region boundary detail to the level of detail contained in the underlying grid. Without constraining the majority of region boundary segments to a grid, region operators such as difference and intersection tend to produce a lot more fine detail. For example, in Figure 21, two circles 901 and 902 are shown with respective regions 903 and 904 that are not grid-aligned. These circles are overlapped yielding difference regions 905 and 907, and intersection region 906. In Figure 22, the same circles 901 and 902 have regions 913 and 914 that are aligned to grid 910. These circles are overlapped yielding difference regions 915 and 917 and intersection region 916. It can be seen in this example that the grid-aligned regions yield less detailed results at the expense of slightly less efficient region coverage. Regions 905, 906 and 907 together contain a total of sixty segments, while regions 915, 916 and 917 together contain only fifty-two.

2.3.2 Creating Binary Region Groups

The region groups of binary nodes in the compositing tree on the other hand are the result of combining the region groups of their child nodes. It will now be explained how region groups are combined to form new region groups. In this section, for simplicity only "OVER" and "IN" binary nodes will be dealt with. The operations required for binary nodes representing other compositing operators can easily be inferred from combining the "OVER" and "IN" cases in various ways.

For the sake of clarity, the method of the preferred embodiment is initially described without reference to optimization based properties such as opacity.

The following notation will be beneficial when considering binary region group creation:

| Notation | |
|-----------------------|---|
| RG1 | The region group of the binary node's left child |
| RG2 | The region group of the binary node's right child |
| RG | The region group of the binary node. It is this region group that is being initialised |
| RG1→urn | The region description representing the union of all RG1's region descriptions (RG1's coverage region). |
| RG2→urn | The region description representing the union of all RG2's region descriptions (RG2's coverage region). |
| RG→urn | The union of all RG's region descriptions (to be initialised) (RG's coverage region) |
| rg1 _i | The current region in RG1 |
| rg2 _j | The current region in RG2 |
| rg1 _i →rgn | rg1 _i 's region description |

| | |
|---------------------------|-------------------------------|
| $rg2_j \rightarrow rgn$ | $rg2_j$'s region description |
| $rg1_i \rightarrow proxy$ | $rg1_i$'s proxy |
| $rg2_j \rightarrow proxy$ | $rg2_j$'s proxy |

2.3.2.1 Constructing "OVER" Region Groups

When constructing "OVER" region groups, only areas where the contributing region groups intersect need to be composited. Areas where one operand does not overlap the other involve no compositing. The basic algorithm is broken into three iterative steps. First, the coverage region of the region group of the binary node that is being initialised ($RG \rightarrow urn$) is made equal to the union of the coverage regions of the binary nodes left child ($RG1 \rightarrow urn$) and the binary node's right child ($RG2 \rightarrow urn$). Then, for each region rg_i in $RG1$, the difference ($diff_rgn$) between that region and $RG2$'s coverage region ($RG2 \rightarrow urn$) is then calculated. If the difference ($diff_rgn$) is non-empty then a new region with $diff_rgn$ as its region description is added to RG . The proxy of this new difference region can be the same as the proxy $rg1_i$. No compositing is required to generate it. The difference regions between $RG2$'s regions and $RG1$'s coverage region are similarly constructed and added to RG . Finally, the intersection ($inter_rgn$) between each region $rg1_i$ in $RG1$ and each region $rg2_j$ in $RG2$ is calculated. If the result of this intersection is non-empty, then a new proxy (new_p) is created by compositing $rg1_i$'s proxy with $rg2_j$'s proxy using the over operation with the $inter_rgn$. A new region is then added to RG with $inter_rgn$ as its region description and new_p as its proxy. A method is described below using pseudo-code.

$RG \rightarrow urn = RG1 \rightarrow urn \text{ union } RG2 \rightarrow urn$

FOR $i = 0$ TO number of regions in $RG1$ DO

$diff_rgn = rg1_i \rightarrow rgn \text{ difference } RG2 \rightarrow urn$

 IF $diff_rgn$ is non-empty THEN

 ADD to RG a new region with $diff_rgn$ as its region description and

$rg1_i \rightarrow proxy$ as its proxy. (*)

 END IF

FOR $j = 0$ TO number of regions in $RG2$ DO

$inter_rgn = rg1_i \rightarrow rgn \text{ intersection } rg2_j \rightarrow rgn$

 IF $inter_rgn$ is non-empty THEN

 create new proxy new_p initialised to OVER of $rg1_i \rightarrow proxy$ and

$rg2_j \rightarrow proxy$ inside $inter_rgn$.

 ADD to RG a new region with $inter_rgn$ as its region description

```
and new_p as its proxy. (+)
    END IF
  END DO
END DO
FOR j = 0 TO number of regions in RG2 DO
  diff_rgn = rg2j→rgn difference RG1→rgn
  IF diff_rgn is non-empty THEN
    ADD to RG a new region with diff_rgn as its region description and
    rg2j→proxy as its proxy. (*)
  END IF
END DO
```

10 The regions added by the ADD operations marked with asterisks (*) above are termed difference regions. This is because their shape is the result of a difference operation. Such regions are very cheap computationally because their proxies require no compositing. The only work involved is the administrative overhead of adding a new region to the region group and the cost of the difference operation itself. In the preferred embodiment its proxy is inherited from the region (in one of the child region groups) on which it is based. 15 It can be seen that proxies which originate low in the compositing tree can be propagated upwards towards the root with minimal overhead (both in terms of speed and memory) by the use of difference regions.

The regions added by the ADD operation marked with the plus (+) are termed inter- 20 section regions. This is because their shape is the result of an intersection operation. The proxies of such regions are more expensive to generate than difference regions because they involve per-pixel compositing operations to be done within the area defined by the intersection. The more fidelity granted the region descriptions, the greater the saving in pixel processing costs, at the cost of a greater administrative overhead (more complex 25 regions require longer to intersect etc).

Figs. 8A to 8D provide a simple example of combining "OVER" region groups using the above method. The region group resulting from the combination contains 5 regions, 3 difference regions and 2 are intersection regions. Fig. 8A represents two region groups RG1 and RG2 which are to be combined. RG1 contains two regions 81 and 82, whereas 30 RG2 only contains a single region 83. As seen in Fig 8B, for each region in RG1, RG2's

region coverage is subtracted from it. If the resultant region is non-empty, it becomes a region in the new region group. In this example both regions 81 and 83 produce non-empty difference regions 84 and 85 respectively. For each region in RG2, RG1's region coverage is subtracted from it, as seen in Fig 8C. In this example difference region 86 is produced. Finally, every region in RG1 is intersected with every region in RG2, as seen in Fig 8D. Any non-empty region becomes a region in the new region group. In this example, regions 81 and 83 produce 87. Further, regions 82 and 83 produce 88.

2.3.2.2 Constructing "IN" Region Groups

The properties of the "IN" operator lead to the fact that an "IN" binary region group only produces pixel data in the region of intersection between the two contributing region groups. Essentially, when compared to the algorithm used for "OVER" region groups, only intersection regions are generated. Therefore, for each region $rg1_i$ of RG1, and for each region $rg2_j$ of RG2 the intersection ($inter_rgn_{ij}$) between $rg1_i$ and $rg2_j$ is calculated. If the intersection is non-empty then a new proxy (new_p) is created by compositing $rg1_i$'s proxy with $rg2_j$'s proxy using the "in" operation within $inter_rgn_{ij}$. A new region is then added to RG with $inter_rgn$ as its region description and new_p as its proxy. The pseudocode for the algorithm is provided below:

```
RG→urn = RG1→urn intersection RG2→urn
FOR i = 0 TO number of regions in RG1 DO
  10   FOR j = 0 TO number of regions in RG2 DO
        inter_rgn =  $rg1_i$ →rgn intersection  $rg2_j$ →rgn
        IF inter_rgn is non-empty THEN
            create new proxy new_p initialised to IN of  $rg1_i$ →proxy and
             $rg2_j$ →proxy inside inter_rgn.
            ADD to RG a new region with inter_rgn as its region description
            and new_p as its proxy. (+)
        15   END IF
    END DO
END DO
```

The major difference between the "IN" and the "OVER" cases is that the "OVER" case generates difference regions while "IN" does not. In the example demonstrated by Figs. 8A to 8D, only new regions 97 and 98 would be generated, as these are intersection

regions. Difference regions 94, 95 and 96 would not be generated using "IN".

Using **Table 2** below and the pseudocode examples of "OVER" and "IN", the relevant code for other compositing operators can be derived.

2.3.2.3 Constructing Region Groups of Other Compositing Operators

Other compositing operators typically generate the same intersection regions as the "OVER" and "IN" cases do. However, they typically differ from one another (as indeed from "OVER" and "IN") in what difference regions they generate. This is dependent on the particular properties of each compositing operator. Table 2 summarises which difference regions are generated for some commonly used compositing operators.

TABLE 2

| Compositing Operator | Generate Diff Rgns from RG1 ? | Generate Diff Rgns from RG2 ? |
|----------------------|-------------------------------|-------------------------------|
| Over | Yes | Yes |
| In | No | No |
| Out | Yes | No |
| Atop | No | Yes |
| Xor | Yes | Yes |
| Plus | Yes | Yes |

2.4 Optimising using Opaque Areas

The preferred embodiment stores within each region a flag indicating whether the pixel data in the region proxy is completely opaque. It is therefore possible to reduce the number of per-pixel compositing operations by exploiting the effect opaque operands have on the compositing operators.

2.4.1 Opaque Area Optimisation for "Over" Region Groups

If an opaque region is "OVER" another region, then there is no need to compute the result of the composite, as no part of the right operand region's proxy is visible through the left operand's opaque proxy. In the preferred embodiment, the resultant region is made to reference the right operand's proxy, which has the same effect as actually doing the composite.

The method in this case is a slightly modified version of the "OVER" region group construction method provided previously. The only difference is that when calculating the intersection region of the current region in RG1 and each region of RG2, a check is carried out to see whether the current region in RG1 is opaque. If this is the case, then the proxy

of the newly calculated region (new_p) will be the proxy of the current region in RG1.

The technique is illustrated using the following pseudocode :

```
RG→urn = RG1→urn union RG2→urn
FOR i = 0 TO number of regions in RG1 DO
5   diff_rgn = rg1→rgn difference RG2→urn
   IF diff_rgn is non-empty THEN
       ADD to RG a new region with diff_rgn as its region description and
       rg1→proxy as its proxy. (*)
   END IF
   FOR j = 0 TO number of regions in RG2 DO
10      inter_rgn = rg1→rgn intersection rg2→rgn
       IF inter_rgn is non-empty THEN
           IF rg1i is OPAQUE THEN
               new_p = rg1→proxy
           ELSE
               create new proxy new_p initialised to OVER of rg1→proxy
               and rg2→proxy inside inter_rgn.
15          END IF
           ADD to RG a new region with inter_rgn as its region description
           and new_p as its proxy. (+)
       END IF
   END DO
END DO
20 FOR j = 0 TO number of regions in RG2 DO
   diff_rgn = rg2→rgn difference RG1→urn
   IF diff_rgn is non-empty THEN
       ADD to RG a new region with diff_rgn as its region description and
       rg2→proxy as its proxy. (*)
   END IF
END DO
```

25

2.4.2 Opaque Area Optimisation for "IN" Region Groups

If a region is "IN" an opaque region, then according to the properties of the "IN" operator, the resultant pixel data is the same as that of the left operand. This can be achieved by having the resultant region simply reference the proxy of the left operand.

30 The method of the preferred embodiment is a slightly modified version of the "IN" region

group construction method provided previously. The only difference is that when calculating the intersection region of the current region in RG1 and each region of RG2, a check is carried out to see whether the current region in RG2 is opaque. If this is the case, then the proxy of the newly calculated region (new_p) will be the proxy of the current region in RG1.

The technique is illustrated using the following pseudocode:

```
RG→urn = RG1→urn intersection RG2→urn
FOR i = 0 TO number of regions in RG1 DO
  FOR j = 0 TO number of regions in RG2 DO
    inter_rgn = rg1i→rgn intersection rg2j→rgn
    IF inter_rgn is non-empty THEN
      IF rg2j is OPAQUE THEN
        new_p = rg1i→proxy
      ELSE
        create new proxy new_p initialised to IN of rg1i→proxy and
        rg2j→proxy inside inter_rgn.
      END IF
      ADD to RG a new region with inter_rgn as its region description
      and new_p as its proxy. (+)
    END IF
  END DO
END DO
```

2.5 Initialising the Entire Tree

The entire compositing tree can be initialised by using the above-described method of the preferred embodiment on every binary region group in the tree. A node cannot be initialised until its children have been initialised. Therefore the process simply starts at the bottom of the tree and works its way up towards the root. The process first checks to see if the current node is a leaf node. If this is the case, then a leaf node region group is constructed. However, in the case that the current node is a binary node then a binary node region group is constructed using the method of the preferred embodiment outlined in sections 2.4.1 and 2.4.2. The following pseudocode outlines a method for initialising all the region groups of the tree. It utilises a recursive function, which is called passing the root of the tree as an argument.


```
tree_init(node : tree ptr)
BEGIN
    IF node is a leaf node THEN
        CONSTRUCT leaf node region group
    ELSE
        tree_init(node→left)
        tree_init(node→right)
        CONSTRUCT binary node region group by combining region groups of
the left and right children
    END IF
END tree_init
```

10 2.6 Constructing the Resultant Image

Once the compositing tree has been initialised, the region group at the root of the tree contains a group of zero or more regions which together represent the partitioning of the resultant image into areas which differ in the way the image data was generated. Some of the regions' proxies may refer to image data directly from leaf nodes of the tree, having not
15 required any compositing. Other regions, on the other hand, may have proxies which are the result of compositing operations. If a single resultant image is required, such as an image stored in a pixel buffer, this can be achieved by copying the image data from each region's proxy to the pixel buffer within the area corresponding to the region. This process is demonstrated in the pseudocode provided below, which is generalised and able to
20 restrict the construction of the final image to any nominated update region.

```
construct_image
(
    output_image : pixel data ptr,
    urn : region description
25 )
BEGIN
    FOR i = 0 TO number of region in RG DO
        int_rgn = rgi→rgn intersection urn
        IF int_rgn is non-empty THEN
            COPY image data from rgi→proxy to output_image inside int_rgn
30        END IF
```

END DO
END construct_image

3.0 Dynamic Rendering

Dynamic Rendering refers to the problem of generating multiple successive images. Given a compositing tree, it is possible to generate it's region groups (containing regions and proxies) using the method described above. Enhancements to this method required to support dynamic rendering are described below. The compositing tree represents an image. Changes to this tree can be made to make it represent a new image. The tree's region groups (and tree region description and proxies) are updated to reflect this modified tree. Performance is improved by exploiting commonality between the two images. A simple example will illustrate the basic techniques and terminology the preferred embodiment uses.

Fig. 3 shows the region subdivision and the respective compositing expressions (advantage is not taken of opacity) for the simple compositing tree. Consider therefore the situation in which object A moves by a small amount relative to the other objects. Some regions in the region group at the root of the tree will be affected by A moving.

If opaque case optimisations are ignored, the regions with compositing expressions which include A will be significantly affected by A moving. The region numbers which are so affected are 2, 3, 5 and 6. When updating the region group at the root of the tree, those regions will need both their region descriptions and their proxies completely recalculated. This situation is known in the preferred embodiment as primary damage. Any region whose compositing equation includes an object which has changed in some way, may be said to suffer primary damage.

Regions that abut regions which have A in their compositing expression are also effected by A moving, though not as severely as those regions with primary damage. In the example, these other affected regions are 1, 4, 7 and 8. When updating the region group at the root of the tree, these regions will need their region descriptions recalculated. However, their proxies will only need to be recalculated in areas of the new region which were not included in the corresponding earlier region. This situation is known in the preferred embodiment as secondary damage. Generally, secondary damage is incurred if an

object upon which a region's boundary (but not content) depends, changes in some way.

In order to reduce the per-frame update cost, it is important to reduce, as far as is practicable, the amount of work necessary, both in terms of per-pixel operations, but also in terms of region group operations. The concepts of primary and secondary damage are a way of facilitating this. If the preferred embodiment is able to accurately determine the minimum set of regions throughout all the compositing tree which have some kind of damage, then obviously the amount of work being done is reduced. The following sections describe how this is achieved.

3.1 Basic Data Model

The data model used for static rendering, consisting as it does of a region description and a proxy, is insufficient for use in dynamic rendering. This is because, for primary and secondary damage to be determined, it must be possible to associate regions of the same content between frames. To support this, some extra information is required in each region in a region group. Therefore, each region in a region group now contains the following data:

(i) A Region Description: A low-level representation of the boundaries of the region. The region descriptions of all the regions in a region group must be mutually exclusive (non-intersecting, non-overlapping).

(ii) A Proxy: Some means of caching the pixel data resulting from applying the operation specified by the compositing expression at every pixel inside the region description. A proxy can be as simple as a 24-bit colour bit-map, or something much more complicated (such as a run-length encoded description). Fundamentally, a proxy simply has to represent pixel data in some way which makes it efficient to retrieve and use.

(iii) A Contents Label: A contents label represents a unique symbolic expression that describes the method of construction of image data. The terms in the symbolic expression distinguish between different categorisations of a source of image data. Therefore, the region groups of two distinct leaf nodes in the compositing tree will contain regions which are labelled with distinct contents labels even if their actual image data is equivalent. The preferred embodiment uses a system of unique integers to represent its contents labels. For example "23" could represent "(A over B) over C".

(iv) A Flag Register: A general-purpose flag register used to store state during the region group update process. The exact flags stored here will be outlined in a later section.

3.2 Contents Labels

Leaf node region groups may contain multiple regions, with each region naturally having a unique contents label. For example, the region group of a leaf node in a compositing tree could contain a single region (tagged with a single contents label) representing the non-transparent area of the leaf node. Alternatively, the leaf node region group could contain two regions (each tagged with a different contents label), one representing the area of the leaf node which is completely opaque, the other representing the remaining non-transparent area. If it offers an advantage, a leaf node may be categorised even further, into an arbitrary number of regions (and associated contents labels).

One way a contents label may be created is by assigning a new one to a region of a leaf node region group. Another way is taking other contents labels and combining them to create a new contents label that represents the symbolic expression that represents the combination of the contributing expressions. For example, if the contents label representing ((A comp B) comp C) is combined with the contents label representing (D comp E) then a new contents label will be created which represents (((A comp B) comp C) comp (D comp E)).

As well as contents labels, dependency information is also required. Dependency information indicates how a given contents label is related to other contents labels, both in terms of how the contents of one region contribute to contents of other regions, and how change of a region boundary affect the boundary of other regions. The preferred embodiment associates the following data with each contents label.

(i) Primary Dependency List: Each primary dependency is a contents label L' to which a contents label L directly contributes. In other words, a "primary dependency" is a contents label L' representing an expression which has been constructed by combining L and some other contents label. Each time contents labels are combined, the contents label for the combination is added to the primary dependencies of all contributors.

(ii) Secondary Dependency List: Each secondary dependency is a contents label L" which may be indirectly affected if the image represented by the contents label L has

changed in some way that affects it's boundary. Whenever contents labels are combined, a contributing contents label is added to the secondary steps of the continuation if and only if the compositing operator yields a difference region with said contributing contents label. Table 2 shows which of some commonly used operators yield difference regions for their left and right operands. In addition, for a combination of (A comp B), the secondary dependencies of the combination contents labels for all (A comp b_i) and all (a_j comp B) are added, where a_j are the secondary dependencies of A and b_i are the secondary dependencies of B.

(iii) Property Information: Each contents label may represent contents which have properties which the compositing engine may be able to exploit. An example is that of opaqueness. If a contents label represents opaque content, then compositing that content could be much faster, as for certain operators, no per-pixel compositing operations would be required.

3.3 Contents Label Implementation

The preferred embodiment uses unique integers as contents labels, and stores a number representing the number of contents labels which currently exist. When a new contents label is created, this number is incremented and becomes the unique integer representing the contents label. This technique of assigning a contents label by monotonically incrementing an integer means that the contents labels' associated data structures may be stored in a one dimensional array which grows as more contents labels are added. A content label's data structure can be referenced simply by using the contents label as an index. When a leaf node contents label is created, it is initialised to have no primary or secondary dependencies. If the current leaf node contents label is opaque, then a flag is set in content label i's properties.

The pseudocode below illustrates the basic techniques used to create a new contents label which is not dependent on other contents labels (ie: a leaf node region group contents label):

Notation

| | |
|----------------------|---|
| opaque | A flag passed to the function which indicates whether or not the contents label represents opaque content or not. |
| cur_clab | A global integer which stores the last contents label created. |
| clabs | A global array which stores the associated data structures of the contents label. |
| clabs[i]->pri_deps | A pointer to the head of content label i's primary dependency list. |
| clabs[i]->sec_deps | A pointer to the head of content label i's secondary dependency list. |
| clabs[i]->properties | A flag register representing contents label i's properties. |

```

create_new_contents_label
(
    opaque : boolean
) : RETURNS unsigned int
BEGIN
    INCREMENT cur_clab.
    clabs[cur_clab]→pri_deps = NULL.
    clabs[cur_clab]→sec_deps = NULL.
    IF opaque THEN
        clabs[cur_clab]→properties = OPAQUE.
    ELSE
        clabs[cur_clab]→properties = 0.
    END IF
    RETURN cur_clab.
END create_new_contents_label.

```

Contents labels can also be created to represent the combination of existing contents labels. This is achieved in the preferred embodiment by a hash table which maps an operation and the contents labels of its operands (hashed together to create a key) to a single contents label representing the result.

When a region is created which represents an intersection between two other regions (each with its own contents label), a new contents label is generated which is used to tag the new region. When this new contents label is generated, it must be added to the primary dependency lists of both its contributing operands. A secondary dependency list which depends on the secondary dependencies of the two contributing contents labels as well as the properties of the compositing operator must also be generated.

The process is recursive and begins by adding the newly created contents label

(new_cl) to the primary dependency lists of the contributing contents labels. Then, depending on the properties of the compositing operator, none, either or both of the contributing contents labels are added to the secondary dependency list. Then every contents label representing (clab1 op sd2_i) and (sd1_i op tab2) are added to the secondary dependency list.

Notation

| | |
|------------------|--|
| clab1 | The first contributing contents label. |
| clab2 | The second contributing contents label. |
| sd1 _i | The i'th element of clab1's secondary dependency list. |
| sd2 _i | The i'th element of clab2's secondary dependency list. |

create_binary_contents_label

(

clab1 : contents label,

clab2 : contents label,

op: compositing operator

)

BEGIN

IF the hash table already contains an entry representing clab1 op clab2

THEN

RETURN the existing contents label representing the combination.

END IF

Generate a new entry in the hash table representing clab1 op clab2, mapping to new_cl.

(Add the new contents label to the primary dependency lists of the contributing contents labels if the compositing op requires it)

add_to_primary_dep_list(clab1, new_cl)

add_to_primary_dep_list(clab2, new_cl)

(Generate the secondary dependencies)

IF op generates left diff rgns THEN

add clab1 to secondary deps

END IF

IF op generates right diff rgns THEN

add clab2 to secondary deps

END IF

FOR i = 0 TO number of elements in sd1 DO

```

        add_to_secondary_dep_list
        (
            new_cl,
            create_binary_contents_label(sd1i, clab2)
        )
    END DO

    FOR i = 0 TO number of elements in sd2 DO
        add_to_secondary_dep_list
        (
            new_cl,
            create_binary_contents_label(clab1, sd2j)
        )
    END DO
END construct_binary_contents_label

```

3.4 Combining Region Groups for Dynamic Rendering

Before any incremental updates can be made to a compositing tree, the compositing tree must be constructed to be in a consistent initial state. The basic technique for achieving this is the same as that used for static rendering, except that support for contents labels is included.

Leaf node region groups are initialised essentially as before, except that each region in each leaf node region group is tagged with a unique contents label. Each contents label may in turn be tagged with various categorisation properties which may help the renderer to be more efficient. For example, a contents label may be tagged as being completely opaque.

The initialisation of binary nodes is also similar to a static rendering case. By way of example, the way in which the region group for an “OVER” binary node is constructed will now be explained. The techniques for constructing the region groups of the other compositing operators can easily be inferred from the “OVER” case.

When a difference region between rg_i of one operand and the coverage region of the other operand is calculated, the difference region inherits the contents label rg_i . When an intersection region is created, on the other hand, a new contents label is created by combining the contents labels of the two contributing regions. This is because the two contrib-

uting regions had their proxies composited into a new proxy which means new content. The pseudocode for constructing an “OVER” region group which includes contents label management is provided below:

| Notation | | |
|----------|-------------------------|--|
| 5 | RG1 | The region group of the binary node's left child |
| | RG2 | The region group of the binary node's right child |
| | RG | The region group of the binary node. It is this region group |
| | | that we are initialising |
| | RG1→urn | The region description representing the union of all RG1's |
| | | region descriptions (RG1's coverage region). |
| | RG2→urn | The region description representing the union of all RG2's |
| | | region descriptions (RG2's coverage region). |
| 10 | RG→urn | The union of all RG's region descriptions. |
| | rg1 _i | The current region in RG1 |
| | rg2 _j | The current region in RG2 |
| | rg1 _i →rgn | rg1 _i 's region description |
| | rg2 _j →rgn | rg2 _j 's region description |
| | rg1 _i →proxy | rg1 _i 's proxy |
| | rg2 _j →proxy | rg2 _j 's proxy |

```
15  RG→urn = RG1→urn union RG2→urn
    FOR i = 0 TO number of regions in RG1 DO
        diff_rgn = rg1i→rgn difference RG2→urn
        IF diff_rgn is non-empty THEN
            ADD to RG a new region with diff_rgn as its region description,
            rg1i→proxy as its proxy and rg1i→clab as its contents label.
        END IF
20  FOR j = 0 TO number of regions in RG2 DO
        inter_rgn = rg1i→rgn intersection rg2j→rgn
        IF inter_rgn is non-empty THEN
            new_clab = GENERATE a new unique contents label as a result
            of combining rg1i→clab and rg2j→clab.
            IF rg1i→clab is OPAQUE THEN
25  new_p = rg1i→proxy
            ELSE
                create new proxy new_p initialised to OVER of rg1i→proxy
                and rg2j→proxy inside inter_rgn.
            END IF
            ADD to RG a new region with inter_rgn as its region description,
30  new_p as its proxy and new_clab as its contents label.
```

```
        END IF
    END DO
END DO
FOR j = 0 TO number of regions in RG2 DO
    diff_rgn = rg2j→rgn difference RG1→urn
    IF diff_rgn is non-empty THEN
        ADD to RG a new region with diff_rgn as its region description,
        rg2j→proxy as its proxy and rg2j→clab as its contents label.
    END IF
END DO
```

3.5 Secondary Dependencies and Over

The rationale behind the algorithm used for generating secondary dependencies requires more explanation. Secondary dependencies are only generated when a new contents label is created by combining two other contents labels. As can be seen in the above pseudocode, this only occurs when an intersection region is generated. Essentially, the preferred embodiment uses contents labels generated for intersection regions as triggers - the regions tagged with two contents labels cannot indirectly affect one another unless they intersect. The secondary dependency list for a particular contents label is dependent on the compositing operator the composite contents label represents, the two contributing contents labels and their secondary dependency lists.

The method of the preferred embodiment of generating a secondary dependency list for a new contents label (C) which represents one contents label (A) composited over another contents label (B) using the "OVER" operator will now be explained. Elements of A's and B's secondary dependency lists are referred to as A_i and B_i respectively. First, both A and B are added to C's secondary dependency list. This is because if the region tagged with C changes its boundary, then it is likely that any regions tagged with A and B will need to be recalculated (because their regions are likely to abut C's region). Next, for each element of B's secondary dependency list, each contents labels representing (A OVER B_i) is added. A mapping representing A OVER B_i may not currently exist in the system and needs to be created. A secondary dependency list may contain contents labels which are not represented by any region in a region group. They could come into existence by changes in region boundaries. The rationale is that A intersects B, and there-

fore it is likely that it also intersects regions tagged with contents labels which exist in B's secondary dependency list. Similarly, for each element of A's secondary dependency list, each contents label representing (A_i OVER B) is added.

3.6 Contents Labels and Damage

The concepts of primary and secondary damage were introduced with reference to Fig. 3 to demonstrate that it is not always necessary to regenerate an entire image as a result of a change to the compositing tree. By keeping track of dependencies between regions of different content, it only becomes necessary to regenerate image data in regions whose contents have become damaged. The following explanation outlines the dependencies and damage for simple compositing tree changes. "Simple" means that only leaf nodes are modified. More complex change scenarios such as tree structure changes etc will be outlined in later sections.

If a leaf node is modified, the contents labels of its affected regions are said to be "primary damaged". Primary-damaging a contents label involves recursively primary-damaging all its primary dependencies. Whenever a contents label is primary-damaged, all its secondary dependencies are non-recursively marked with secondary damage. The process begins by flagging the contents label to be damaged. The following pseudocode demonstrates how contents labels can be damaged:

| Notation | |
|-----------------|---|
| clab | The contents label to be damaged |
| pd _i | The i'th element of clab's primary dependency list. |
| sd _i | The i'th element of clab's secondary dependency list. |

```
damage_contents_label
(
    clab : contents label,
)
BEGIN
    FLAG clab with PRIMARY damage

    FOR i = 0 TO number of elements in sd DO
        FLAG sdi with SECONDARY damage
    END DO
```

```
FOR i = 0 TO number of elements in pd DO
    damage_contents_label(pdi)
END DO
END damage_contents_label
```

When a tree update occurs, any region with its contents label marked as having primary damage will need to recalculate both its region boundaries and its proxy. Any region with its contents label marked as having secondary damage will need to recalculate its region description but will only need to recalculate its proxy in areas of the new region that were not included in the earlier region.

3.7 Examples of Contents Labels and Dependencies

In order to clarify the concepts of contents labels and damage, some examples of varying complexity will be presented.

3.7.1 Example 1

Fig. 9 will result in the following contents label table after the compositing tree is initially constructed (Note: in the following table contents labels are represented as unique strings not as integers where "over" has been abbreviated to "o". This is simply for readability.):

| Contents Label | Primary Deps. | Secondary Deps. |
|----------------|---------------|-----------------|
| A | AoB | |
| B | AoB | |
| AoB | | A, B |

If A moves, then AoB will have primary damage, resulting in B having secondary damage.

3.7.2 Example 2

Fig. 10 will result in the following contents label table after the compositing tree is initially constructed:

| Contents Label | Primary Deps. | Secondary Deps. |
|----------------|---------------|-----------------|
| A | AoB, AoC | |

| | | |
|---------|-------------------|------------------|
| B | AoB, BoC | |
| AoB | AoBoC | A, B |
| C | AoC, BoC, (AoB)oC | |
| AoC | | A, C |
| BoC | | B, C |
| (AoB)oC | | AoB, C, AoC, BoC |

In this example, every object intersects every other object, so if something changes, everything will be damaged in some way - everything which is a primary dependency of the changed object will have primary damage, whereas everything else will have secondary damage.

Fig. 11 illustrates the effect of A moving in a subsequent frame. As can be seen, if A is damaged, the regions defined by A, AoB, AoC and (AoB)oC will each have primary damage. The regions defined by B, C and BoC will each have secondary damage.

3.7.3 Example 3

Fig. 12 will result in the following contents label table after the compositing tree is initially constructed:

| Contents Label | Primary Deps. | Secondary Deps. |
|----------------|---|--|
| A | AoB, AoC, AoE, Ao(DoE), AoD | |
| B | AoB, BoC, BoE | |
| AoB | AoBoE | A, B |
| D | DoE, AoD, CoD, (AoC)oD | |
| E | DoE, AoE, (AoB)oE, BoE, CoE, (BoC)oE, (AoC)oE | |
| DoE | Ao(DoE), (AoC)o(DoE), Co(DoE) | D, E |
| C | AoC, BoC, Co(DoE), CoE, CoD | |
| AoC | AoCoE, (AoC)o(DoE), (AoC)oD | A, C |
| BoC | (BoC)oE | B, C |
| AoE | | A, E |
| (AoB)oE | | AoB, E, AoE, BoE |
| BoE | | B, E |
| CoE | | C, E |
| (BoC)oE | | BoC, E, BoE, CoE |
| AoD | | A, D |
| CoD | | C, D |
| (AoC)oE | | AoC, E, AoE, CoE |
| Ao(DoE) | | A, DoE, AoD, AoE |
| Co(DoE) | | C, DoE, CoD, CoE |
| (AoC)o(DoE) | | AoC, DoE, Ao(DoE), Co(DoE), (AoC)oD, (AoC)oE |
| (AoC)oD | | AoC, D, AoD, CoD |

Since A intersects every other object, if it moves, a large amount of the compositing tree will need to be recomputed. Fig. 13 shows that the only part left alone is the area corresponding to BoC and its dependent BoCoE. To summarise:

- Primary Damage - A, AoB, AoC, AoE, Ao(DoE), (AoB)oE, (AoC)oE, (AoC)o(DoE), AoD, (AoC)oD
- Secondary Damage - B, C, E, DoE, BoE, CoE, DoE, CoDoE

On the other hand, if B moves, the amount of damage is less than if A moved. This is because B doesn't intersect D. DoE, Ao(DoE), (AoC)o(DoE), Co(DoE) and (AoC)oE (and their ancestors) are not damaged when B moves. This is shown in Fig. 14. The rest of the damage is summarised as:

- Primary Damage - B, AoB, BoC, BoE, (AoB)oE, (BoC)oE
- Secondary Damage - A, E, C, AoE, CoE

The examples presented so far are very simple, but they are sufficient to demonstrate that the dependencies techniques presented so far will damage those contents labels which are affected when a particular contents label/s is(are) damaged. In a typical complex composite, it is rare for large numbers of objects to intersect a large number of other objects, meaning that large areas of the compositing tree should be untouched during updates using the above technique.

3.8 Example of Secondary Dependencies and Compositing Operators

Consider a modified version of Example 3 above. Fig. 18, will result in the following contents label table after the compositing tree is initially constructed. Note that AaB represents A ATOP B and AiB represents A IN B etc:

| Contents Label | Primary Deps | Secondary Deps |
|----------------|--------------|----------------|
| A | AaB | |
| B | AaB, BoC | |
| AaB | | B |
| C | BoC, Co(DiE) | |
| BoC | | B, C |
| D | DiE | |
| E | DiE | |
| DiE | Co(DiE) | |
| Co(DiE) | | C, DiE |

As seen in Fig. 18, the ATOP operator clips A to B's bounds, meaning that intersections

between A and any of C, D or E never occur. Similar things occur with the IN operator. This means that the objects in this scene are less tightly coupled. For example, if A is changed, then only B and AaB are immediately damaged. Similarly, if E is damaged, it is only possible for it to damage DiE.

3.9 Updating Region Groups

The preferred embodiment uses the contents label and damage framework to reduce the amount of work it has to do to make a binary region group consistent with its updated operands during an update. It does this by only updating those regions in a region group whose contents labels have primary or secondary damage, adding any new region which comes into existence as a result of the changes made to the compositing tree, and deleting any region in the right group whose contact no longer exists.

Each different binary operator has a different updating function which deals with the specific requirement of that operator. The process of updating region groups is a two-pass process. The first pass updates any intersection regions that have been primary damaged and adds any new intersection regions generated due to the damage. Each region of one operand's region group is intersected with each region of the other operand's region group whenever one or both of their corresponding contents labels are primary damaged. If the intersection is non-empty, then the preferred embodiment determines if a contents label representing the combination exists. If it doesn't, one is created and primary damaged.

Note that primary damaging a contents label will mark all it's secondary dependencies with secondary damage.

If a region in the region group is currently tagged with this contents label, its boundary and proxy are updated. If no such region exists in this region group, then a new region keyed by this contents label is added to the region group. A new proxy is generated and assigned to this region along with the right description relating from the intersection operation.

A difference between each region group of one operand and the coverage region of the other operand is calculated whenever the regions contents label has primary or secondary damage. If the difference is non-empty and a region tagged with the contents label exists in the region group, then its region description and proxy reference are updated. If such a region doesn't exist then a region keyed by the contents label is added to the region group.

It is assigned as a coverage region of the difference result and references the proxy of current region.

Each region of one operand's region group is interacted with each region of the other operand's region group whenever the contents label representing their combination has secondary damage and no primary damage. If the intersection is non-empty, the region group is searched looking for a region keyed by the contents label. If such a region exists its region description is updated and it's proxy is updated as the difference between the new and old regions. If such a region doesn't exist, then a region keyed by the contents label is created. It's region description is assigned the result of the interaction operation and it's proxy generated.

Pseudocode which illustrates a simple algorithm for updating a binary "OVER" region group is provided below.

| Notation | |
|-------------------------|---|
| RG1 | The region group of the binary node's left child |
| RG2 | The region group of the binary node's right child |
| RG | The region group of the binary node. It is this region group that is being initialised. |
| RG1→urn | The region description representing the union of all RG1's region descriptions (RG1's coverage region). |
| RG2→urn | The region description representing the union of all RG2's region descriptions (RG2's coverage region). |
| RG→urn | The union of all RG's region descriptions. |
| rg1 _i | The current region in RG1 |
| rg2 _j | The current region in RG2 |
| rg1 _i →rgn | rg1 _i 's region description |
| rg2 _j →rgn | rg2 _j 's region description |
| rg1 _i →proxy | rg1 _i 's proxy |
| rg2 _j →proxy | rg2 _j 's proxy |
| rg1 _i →clab | rg1 _i 's contents label |
| rg2 _j →clab | rg2 _j 's contents label |

RG→urn = RG1→urn union RG2→urn

(First Pass - this pass is used to deal with primary damage of intersection regions and any new intersection regions generated)

FOR i = 0 TO number of regions in RG1 DO

 FOR j = 0 TO number of regions in RG2 DO

 IF rg1_i→clab has PRIMARY damage OR rg2_j→clab has PRIMARY DAMAGE THEN

 inter_rgn = rg1_i→rgn intersection rg2_j→rgn

IF inter_rgn is non-empty THEN
 comp_clab = SEARCH for an existing contents label which
 represents (rg1_i→clab comp rg2_j→clab).

 IF a region tagged with comp_clab already exists in RG
THEN

 IF rg1_i→clab is OPAQUE THEN
 new_p = rg1_i→proxy
 ELSE
 create new proxy new_p initialised to OVER of
 rg1_i→proxy and rg2_j→proxy inside inter_rgn.

 END IF
 MODIFY the existing region to have inter_rgn as its
 region description and new_p as its proxy.

 ELSE
 new_clab = create_binary_contents_label(rg1_i→clab,
 rg2_j→clab).

 IF rg1_i→clab is OPAQUE THEN
 new_p = rg1_i→proxy
 ELSE
 create new proxy new_p initialised to OVER of
 rg1_i→proxy and rg2_j→proxy inside inter_rgn.

 END IF
 damage_contents_label(new_clab)
 ADD to RG a new region with inter_rgn as its region
 description, new_p as its proxy and new_clab as its contents label. (+)

 END IF
 FLAG the region as being 'RETAIN AFTER UPDATE'

 END IF

 END IF

 END DO

END DO

(Second Pass - this pass is used to deal with primary and secondary damage of
difference regions and secondary damage of intersection regions)

FOR i = 0 TO number of regions in RG1 DO

 IF rg1_i→clab has PRIMARY or SECONDARY damage THEN

 diff_rgn = rg1_i→rgn difference RG2→urgn

 IF diff_rgn is non-empty THEN

 IF a region tagged with rg1_i→clab already exists in RG THEN

MODIFY it to have diff_rgn as its region description and
rg1_i→proxy as its proxy.

ELSE

ADD to RG a new region with diff_rgn as its region description, rg1_i→proxy as its proxy and rg1_i→clab as its contents label. (*)

END IF

FLAG the region as being 'RETAIN AFTER UPDATE'

END IF

END IF

FOR j = 0 TO number of regions in RG2 DO

comp_clab = SEARCH for an existing contents label which represents
(rg1_i→clab comp rg2_j→clab).

IF comp_clab exists AND comp_clab has SECONDARY damage but
NO PRIMARY damage THEN

inter_rgn = rg1_i→rgn intersection rg2_j→rgn

IF inter_rgn is non-empty THEN

GET a reference to the existing region tagged in this region
group with comp_clab which MUST exist in this region group

IF rg1_i→clab is OPAQUE THEN

existing regions proxy = rg1_i→proxy

ELSE

update_rgn = inter_rgn difference the region's previous
region description.

update existing regions proxy to include OVER of
rg1_i→proxy and rg2_j→proxy inside update region.

END IF

MODIFY the existing region to have inter_rgn as its region
description and new_p as its proxy.

FLAG the region as being 'RETAIN AFTER UPDATE'

END IF

END IF

END DO

END DO

FOR j= 0 TO number of regions in RG2 DO

IF rg2_j→clab has PRIMARY or SECONDARY damage THEN

diff_rgn = rg2_j→rgn difference RG1→urn

IF diff_rgn is non-empty THEN

IF a region tagged with rg2_j→clab already exists in RG THEN

MODIFY it to have diff_rgn as its region description and
rg2j→proxy as its proxy.

ELSE

ADD to RG a new region with diff_rgn as its region description,
rg2j→proxy as its proxy and rg2j→clab as its contents label. (*)

END IF

FLAG the region as being 'RETAIN AFTER UPDATE'

END IF

END IF

END DO

DELETE all regions of RG which are not marked RETAIN AFTER UPDATE but
whose contents labels have damage, and CLEAR flag in retained regions.

4.0 Tree Modifications (Linking and Unlinking)

More complex changes to a compositing tree can be achieved by changing its structure.

Most typical tree structure changes can be made by using two low level operations, link
and unlink.

The unlink operation is used to separate a child node from its parent. After the operation is
completed, the child node has no parent (meaning it can be linked in somewhere else), and
the parent has a link available (meaning that some other node can be linked there instead).

Nodes in the compositing tree above the unlinked child contain content which is depend-
ent on the unlinked child. Therefore, at the time of the next update, the contents label

present in the unlinked child at the time of unlinking must be damaged to ensure that the
dependent region groups higher in the tree are appropriately updated. This is achieved by
the parent node caching away those contents label existing in its unlinked child. If another
subtree is linked in its place and subsequently unlinked without the region group of the
parent being updated, it is not necessary to cache the contents labels of this new subtree.

Pseudocode for the unlink operation is provided below. Note that the UNLINKED_LEFT
or UNLINKED_RIGHT flag is set so that the contents labels of a newly linked subtree
may be damaged when region groups (including their proxies) higher in the tree must then
be updated.

unlink

(

node : compositing tree node

```
)
BEGIN
    parent = node →parent.
    node →parent = NULL.
    IF node is parent's left child THEN
        parent →left = NULL.
        IF parent doesn't have UNLINKED_LEFT set THEN
            SET the UNLINKED_LEFT flag in parent.
        ELSE.
            RETURN.
        END IF
    ELSE IF node is parent's right child THEN
        parent →right = NULL.
        IF parent doesn't have UNLINKED_RIGHT set THEN
            SET the UNLINKED_RIGHT flat in parent.
        ELSE
            RETURN
        END IF
    END IF
    COPY all the contents labels in node's region group into an array stored in
    parent →unlinked_clabs.
END unlink
```

The link operation involves linking a node with no parent to a free link of a parent node.

Pseudocode for the operation is provided below.

```
link
(
    child : compositing tree node,
    parent : compositing tree node,
    which_link : either LEFT or RIGHT
)
BEGIN
    child →parent = parent
    IF which_link is LEFT THEN
        parent →left = child.
    ELSE
        parent → right = child.
    END IF
END LINK
```

4.1 Updating the Entire Compositing Tree

If a leaf node in the compositing tree changes, the region group of every node in a direct

line from the leaf node to the root of tree must be updated using the techniques described above. Fig. 15 shows circled those nodes which need to have their region groups updated if leaf nodes B and H change in some way.

5 Pseudocode for the tree updating routine is provided below:

```
update_tree
(
    node : compositing tree node
)
BEGIN
    IF node is leaf node THEN
10      Render the leaf node and update its region group.
    ELSE
        IF unlinking occurred in left subtree or left subtree contains dirty leaf
nodes THEN
            update_tree(node →left).
        END IF.
15      IF unlinking occurred in right subtree or right subtree contains dirty leaf
nodes THEN
            update_tree(node →right).
        END IF.
        IF node has UNLINKED_LEFT or UNLINKED_RIGHT flags set THEN
            CALL damage_contents_label on every element of
node→unlinked_clabs.
20      IF node has UNLINKED_LEFT set THEN
            CALL damage_contents_label on every contents label exist-
ing in node→left's region group.
            CLEAR the UNLINKED_LEFT flag in node.
        END IF
        IF node has UNLINKED_RIGHT set THEN
            CALL damage_contents_label on every contents label exist-
25      ing in node→right's region group.
            CLEAR the UNLINKED_RIGHT flag in node.
        END IF
        END IF
        CALL the region group update routine appropriate for node's composit-
ing operator.
        END IF
30      END update_tree
```

The embodiments of the invention can be implemented using a conventional general-purpose computer system 2100, such as that shown in Fig. 19, wherein the process described with reference to Fig. 1 to Fig. 18 is implemented as software recorded on a computer readable medium that can be loaded into and carried out by the computer. The computer system 2100 includes a computer module 2101, input devices 2102, 2103 and a display device 2104.

With reference to Fig 19, the computer module 2101 includes at least one processor unit 2105, a memory unit 2106 which typically includes random access memory (RAM) and read only memory (ROM), input/output (I/O) interfaces including a video interface 2107, keyboard 2118 and mouse 2120 interface 2108 and an I/O interface 2110. The storage device 2109 can include one or more of the following devices: a floppy disk, a hard disk drive, a CD-ROM drive or similar a non-volatile storage device known to those skilled in the art. The components 2105 to 2110 of the computer module 2101, typically communicate via an interconnected bus 2114 and in a manner which results in a usual mode of operation of the computer system 2100 known to those in the relevant art. Examples of computer systems on which the embodiments can be practised include IBM-PC/ATs and compatibles, Sun Sparcstations or alike computer system. In particular, the pseudocode described herein can be programmed into any appropriate language and stored for example on the HDD and executed in the RAM 2106 under control of the processor 2105 with the results being stored in RAM within the video interface 2107 and reproduced on the display 2116. The programs may be supplied to the system 2100 on a pre-programmed floppy disk or CD-ROM or accessed via a connection with a computer network, such as the Internet.

The foregoing describes only one embodiment of the present invention, and modifications, obvious to those skilled in the art, can be made thereto without departing from the scope of the present invention.

Aspects of the Invention: (O-SCRN01)

1. A method of creating an image, said image to be formed by rendering and compositing at least a plurality of graphical objects, each said object having a predetermined outline, said method comprising the steps of:

dividing a space in which said outlines are defined into a plurality regions, each said region being defined by at least one region outline substantially following at least one of said predetermined outlines or parts thereof and being substantially formed by segments of a virtual grid encompassing said space;

manipulating said regions to determine a plurality of further regions, wherein each said further region has a corresponding compositing expression;

classifying said further regions according to at least one attribute of said graphical objects within said further regions;

modifying each said corresponding compositing expression according to a classification of each said further region to form an augmented compositing expression for each said further region;

compositing said image using each of said augmented compositing expressions.

2. A method according to paragraph 1, wherein said attribute is selected from the group consisting of colour, opacity and object outline.

3. A method according to paragraphs 1 or 2, wherein said manipulating said regions comprises applying set operations to said regions.

4. A method according to paragraph 3, wherein said set operations include difference and/or intersection operations.

5. A method according to any one of the preceding paragraphs, wherein said grid is regularly spaced and preferably orthogonally based.

6. A method according to any one of paragraphs 1 to 4, wherein said grid is irregularly shaped.

7. A method according to any one of paragraphs 1 to 6, wherein the compositing expression is a hierarchically structured representation of the image.

5 8. A method according to any one of paragraphs 1 to 7, wherein said image is at least in part a pixel-based image.

9. A method according to any one of the preceding paragraphs, wherein a flag is stored to indicate whether data of an object is opaque or ordinary.

10

10. A method according to paragraph 9, wherein said compositing expression is optimized based on a value of said flag for contributing objects.

11. A method according to any one of the preceding paragraphs, wherein a
15 wholly opaque object in said region acts to eliminate one or more objects within said region from said compositing expressions.

12. A method according to any one of the preceding paragraphs, wherein a
20 wholly transparent object in said region eliminates at least itself from said compositing expression.

25

13. A method according to any one of paragraphs 1 to 12, wherein said modifying comprises modifying a manner in which said compositing expression is evaluated without modifying said hierarchically structured representation.

30

14. A method of creating an image, said image to be formed by rendering and compositing at least a plurality of graphical objects, each said object having a predetermined outline, said method comprising the steps of:

dividing a space in which said outlines are defined into a plurality regions, each said region being defined by at least one region outline substantially following at least one of said predetermined outlines or parts thereof and being substantially formed

by segments of a virtual grid encompassing said space, wherein each object has two region outlines arranged either side of said predetermined outline to thus define three regions for each said object, and wherein each said region has a corresponding compositing expression;

5 classifying said regions according to at least one attribute of said graphical objects within said regions;

 modifying each said corresponding compositing expression according to a classification of each said region to form an augmented compositing expression for each said region;

10 compositing said image using each of said augmented compositing expressions.

15. A method of displaying an image substantially as described herein with reference to any one of the embodiments or Examples.

15 16. Apparatus configured to perform the method of any one of the preceding paragraphs.

20 17. A computer program product including a computer readable medium having a plurality of software modules configured to perform the method of any one of paragraphs 1 to 15.

Dated 3 September, 1998
Canon Kabushiki Kaisha
Patent Attorneys for the Applicant/Nominated Person
SPRUSON & FERGUSON

APPENDIX 1

```
* region.cpp
*
* The implementation of the region manipulation functionality in
* the Screen OpenPage prototype.
*/

#include "protos.h"

static int    union_tot = 0;
static int    int_tot = 0;
static int    diff_tot = 0;
static int    union_full = 0;
static int    int_full = 0;
static int    diff_full = 0;

/*
 * Some #defs which are used to control the optimisations used in the region
 * builder implementation..
 */
#define R_USE_NEW_IMP
#define RB_FAST_SHIFT_AND_DUP_LOOPS
#define RB_USE_LOOKUP
#define R_NEW_IMP_CONSTRUCTION_LOOP

/*
 * The global variables used to store the temporary results needed during
 * region manipulation operations. Two statically allocated
 * R_RegionBuilder structures are used. This is to allow data to be
 * read from one of them whilst the data required for the next operation
 * is written into the other one. Two pointers r_PrevRB and R_CurRB are
 * used to swap access to the two static structures. The
 * r_grow_region_builder function is called to grow a R_RegionBuilder
 * structure if required.
 */
static R_RegionBuilderr_RB1 = {0, 0, NULL, NULL};
static R_RegionBuilderr_RB2 = {0, 0, NULL, NULL};
static R_RegionBuilder*r_PrevRB = &r_RB1;
R_RegionBuilder                *R_CurRB  = &r_RB2;

/*
 * r_shift_and_dup
 *
 * A 16-byte lookup table which when provided with an unsigned char
 * of the following form xxyy, simply produces xxxx. This lookup table
 * _assumes_ that R_STATE_SIZE is 2. It _won't_ work (and will die
 * horribly) if this isn't the case.
 */
unsigned char r_shift_and_dup[16] = {

0x00, 0x00, 0x00, 0x00,

0x05, 0x05, 0x05, 0x05,

0x0A, 0x0A, 0x0A, 0x0A,

0x0F, 0x0F, 0x0F, 0x0F
```

};

```
/*
 * A buffer is required to store the new region data whilst a region is
 * being constructed. This buffer is expanded when required.
 */
static R_Int  *r_RgnBuf = NULL;
static int      r_RgnBufSize = 0;

/*
 * A buffer of IntXYMinMax structures is required to store the rectangles
 * generated during R_rects_from_region. This buffer is expanded when
 * required.
 */
static IntXYMinMax*r_RectBuf = NULL;
static int      r_RectBufSize = 0;

/*
 * R_FREE_LIST_GROWTH_SIZE
 *
 * This macro defines the number of elements which will be added to
 * free list whenever it is grown.
 */
#define R_FREE_LIST_GROWTH_SIZE100
/*
 * r_free_list
 *
 * A linked list of unused R_RgnGrowItems which may be used during
 * region construction.
 */
R_RgnGrowItem *r_free_list = NULL;
/*
 * r_growth_list
 *
 * A linked list of R_RgnGrowItems which represents the current
 * state during region construction.
 */
R_RgnGrowItem *r_growth_list = NULL;

/*
 * r_grow_region_builder
 *
 * This function simply checks to see if a R_RegionBuilder structure
 * is of the required size. If it isn't the size of both
 * arrays in the R_RegionBuilder structure are doubled.
 *
 * Parameters:
 *          rb          The region builder
 *          size        The required size of the arrays in the R_RegionBuilder.
 * Returns:
 *          TRUE on success, FALSE on failure.
 */
static int
r_grow_region_builder
(
    R_RegionBuilder      *rb,
```

```

                                R_Int                                size
)
{
    unsigned char *new_state_data;
    R_Int                                *new_rgn_data;
    int
new_size;

    new_size = max(size, rb->rrb_Size * 2);
    new_state_data = (unsigned char *)malloc(new_size *
sizeof(unsigned char));
    if (new_state_data == NULL)
        return FALSE;
    new_rgn_data = (R_Int *)malloc(new_size * sizeof(R_Int));
    if (new_rgn_data == NULL)
    {
        free(new_state_data);
        return FALSE;
    }
    if (rb->rrb_StateData != NULL)
    {
        memcpy
        (
            new_state_data,
            rb->rrb_StateData,
            rb->rrb_Size * sizeof(unsigned
char)

        );
        free(rb->rrb_StateData);
    }
    if (rb->rrb_RgnData != NULL)
    {
        memcpy
        (
            new_rgn_data,
            rb->rrb_RgnData,
            rb->rrb_Size * sizeof(R_Int)
        );
        free(rb->rrb_RgnData);
    }
    rb->rrb_StateData = new_state_data;
    rb->rrb_RgnData = new_rgn_data;
    rb->rrb_Size = new_size;
    return TRUE;
}

/*
 * r_swap_region_builders
 *
 * This function simply swaps the static pointers to the r_RB1 and r_RB2
 * region builders.
 *
 * Parameters:
 *
 * None.
 *
 * Returns:
 *
 * Nothing.
 */
inline static void
```

```
r_swap_region_builders()
{
    R_RegionBuilder*tmp;
    tmp = R_CurRB;
    R_CurRB = r_PrevRB;
    r_PrevRB = tmp;
}

/*
 * R_add_row_to_region_builder
 *
 * This function adds a row from a R_Region to a R_RegionBuilder structure.
 * The region from which the row comes is passed as an argument. "Adding"
 * has the following conditions...
 *
 * If a pixel run in the row does not exist in the
region builder
 *
it is added and it's current state is tagged
with the region
 *
to which the row belongs. The previous state
is set to 0,
 *
indicating that it did not exist before.
 *
* If a pixel run in the row did exist before, but
it's present state
 *
indicates that it came from the other region
then the run
 *
is retained but it's state is modified to
indicate that
 *
both regions are active at this point.
 *
* If a pixel run in the row did exist before, and
it's present
 *
state indicates that the current region then
the region is
 *
removed and it's state is modified to indicate
that the run is
 *
now empty.
 *
* If a pixel run in the row did exist before, and
it's present state
 *
indicates that both regions are currently
active then the run
 *
is retained, but its state is modified to
indicate that only the
 *
other region is active in this run.
 *
 * Parameters:
 *
row_ptr                A R_Int ** pointer
to the row in the region. Used
 *
to
return the updates row pointer.
 *
rgn_mask                A mask for the region the row
comes from. Must
 *
be
either 1 or 2.
 *
first                Whether this is the
first region to be processed
 *
on
the current scanline.
 * Returns:
 *
TRUE on success, FALSE on failure.
```

```
    */
    #if 1
    int
    R_add_row_to_region_builder
    (
        R_Int          **row_ptr,
        int             rgn_mask,
        int             first
    )
    {
        R_Int          *row;
        int
src_index;
        int
dest_index;
        R_Int          rb_run_start;
        unsigned char rb_prev_run_state;
        int
row_on;

        ASSERT(rgn_mask == 1 || rgn_mask == 2);

        row = *row_ptr;
        r_swap_region_builders();
        /*
         * Skip over the row's y value at the beginning.
         */
        ASSERT(*row == R_NEXT_IS_Y);
        row += 2;
        ASSERT(*row != R_NEXT_IS_Y && *row != R_EOR);

        if (r_PrevRB->rrb_Nels == 0)
        {
            /*
             * If the current region builder's src region
data array is empty, then
             * we are dealing with an empty region builder.
We simply convert
             * the input row to the region builder format.
             */
            row_on = TRUE;
            dest_index = 0;
            while (R_NOT_END_OF_ROW(*row))
            {
                if (++dest_index > R_CurRB-
>rrb_Size)
                {
                    if
(!r_grow_region_builder(R_CurRB, dest_index))
return FALSE;
                }
                R_CurRB->rrb_RgnData[dest_index -
1] = *row;
                if (row_on)
                    R_CurRB-
>rrb_StateData[dest_index - 1] =
```

```
(rgn_mask << RB_STATE_SIZE);

else
    R_CurRB->rrb_StateData[dest_index - 1] = 0;
    row_on = !row_on;
    row++;
}
*row_ptr = row;
R_CurRB->rrb_Nels = dest_index;
return TRUE;
}
/*
 * Firstly, we copy any runs from the region builder which
 * precede this run from the region. We are checking the
 * starting row against the start of each pixel run. Therefore
 * we start checking against the 1st region builder data
 * element.
 */
ASSERT(r_PrevRB->rrb_Nels >= 2);
src_index = 0;
while (src_index < r_PrevRB->rrb_Nels && *row > r_PrevRB->rrb_RgnData[src_index])
    src_index++;
dest_index = src_index;
if (src_index > 0)
{
    if (src_index > R_CurRB->rrb_Size)
    {
        if
        (!r_grow_region_builder(R_CurRB, src_index))
            return FALSE;
    }
    memcpy
    (
        R_CurRB->rrb_RgnData,
        r_PrevRB->rrb_RgnData,
        src_index * sizeof(R_Int)
    );
    if (!first)
    {
        memcpy
        (
            R_CurRB->rrb_StateData,
            r_PrevRB->rrb_StateData,
            src_index *
            sizeof(unsigned char)
        );
    }
    else
    {
        int i = 0;
        unsigned char *src;
        unsigned char *dest;
    }
}
```

```
src_index;
src_index;

i > 10; i--)

#ifndef RB_USE_LOOKUP
*(dest - i) = (*(src - i) & RB_CUR_STATE_MASK) |
(*(src - i) >> RB_STATE_SIZE);
#else
*(dest - i) = r_shift_and_dup(*(src - i));
#endif

#ifndef RB_USE_LOOKUP
*(src - 10) & RB_CUR_STATE_MASK |
(*(src - 10) >> RB_STATE_SIZE);
#else
r_shift_and_dup(*(src - 10));
#endif

#ifndef RB_USE_LOOKUP
*(src - 9) & RB_CUR_STATE_MASK |
(*(src - 9) >> RB_STATE_SIZE);
#else
r_shift_and_dup(*(src - 9));
#endif

#ifndef RB_USE_LOOKUP
*(src - 8) & RB_CUR_STATE_MASK |
(*(src - 8) >> RB_STATE_SIZE);
#else
r_shift_and_dup(*(src - 8));
#endif

#ifndef RB_USE_LOOKUP
*(src - 7) & RB_CUR_STATE_MASK |

src = r_PrevRB->rrb_StateData +
dest = R_CurRB->rrb_StateData +
switch (src_index)
{
default:
    for (i = src_index;
    {

}
/* FALLTHROUGH!! */
case 10:
    *(dest - 10) =
    *(dest - 10) =
    /* FALLTHROUGH!! */
case 9:
    *(dest - 9) =
    *(dest - 9) =
    /* FALLTHROUGH!! */
case 8:
    *(dest - 8) =
    *(dest - 8) =
    /* FALLTHROUGH!! */
case 7:
    *(dest - 7) =
```



```
(* (src - 7) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 7)];
#endif

#ifdef R_USE_LOOKUP

(* (src - 6) & RB_CUR_STATE_MASK) |

(* (src - 6) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 6)];
#endif

#ifdef RB_USE_LOOKUP

(* (src - 5) & RB_CUR_STATE_MASK) |

(* (src - 5) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 5)];
#endif

#ifdef RB_USE_LOOKUP

(* (src - 4) & RB_CUR_STATE_MASK) |

(* (src - 4) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 4)];
#endif

#ifdef RB_USE_LOOKUP

(* (src - 3) & RB_CUR_STATE_MASK) |

(* (src - 3) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 3)];
#endif

#ifdef RB_USE_LOOKUP

(* (src - 2) & RB_CUR_STATE_MASK) |

(* (src - 2) >> RB_STATE_SIZE);
#else
```

```
*(dest - 7) =

/* FALLTHROUGH!! */

case 6:

*(dest - 6) =

*(dest - 6) =

/* FALLTHROUGH!! */

case 5:

*(dest - 5) =

*(dest - 5) =

/* FALLTHROUGH!! */

case 4:

*(dest - 4) =

*(dest - 4) =

/* FALLTHROUGH!! */

case 3:

*(dest - 3) =

*(dest - 3) =

/* FALLTHROUGH!! */

case 2:

*(dest - 2) =
```

```

r_shift_and_dup[* (src - 2)];
#endif

/* FALLTHROUGH!! */
case 1:

#ifdef RB_USE_LOOKUP
    (* (src - 1) & RB_CUR_STATE_MASK) |
    (* (src - 1) >> RB_STATE_SIZE);
#else
    r_shift_and_dup[* (src - 1)];
#endif

/* FALLTHROUGH!! */
case 0:
    ;
    /* FALLTHROUGH!! */
}

}
if (src_index == r_PrevRB->rrb_Nels)
{
    /*
     * We've already exhausted the previous region
     * of the next pixel run to be the max. possible
     * to be 0.
     */
    rb_run_start = R_INT_MAX_VALUE - 2;
    rb_prev_run_state = 0;
}
else
{
    /*
     * We are still within the previous region
     * the run info appropriately.
     */
    rb_run_start = r_PrevRB->rrb_RgnData[src_index];
    if (src_index == 0)
        rb_prev_run_state = 0;
    else
        rb_prev_run_state = r_PrevRB->rrb_StateData[src_index - 1];
}

/*
 * We can now start dealing with the elements in the row.
 */
row_on = 1;
while (R_NOT_END_OF_ROW(*row))
{
    if (*row < rb_run_start)
    {

```

```

>rrb_Size)
    if (dest_index + 1 > R_CurRB-
    {
        if
        (!r_grow_region_builder(R_CurRB, dest_index + 1))
        return FALSE;
    }
    R_CurRB->rrb_RgnData[dest_index]
    = *row;
    if (first)
    {
        /*
        * We are
        * copy the current
        * RB_STATE_SIZE
        */
        R_CurRB-
        >rrb_StateData[dest_index] = (rb_prev_run_state & RB_CUR_STATE_MASK) |
        (rb_prev_run_state >> RB_STATE_SIZE);
    }
    else
    {
        /*
        * We are
        * has already been
        * just copy the
        */
        R_CurRB-
        >rrb_StateData[dest_index] = rb_prev_run_state;
    }
    /*
    * Now, if the row for the current
    * xor the region mask with the
    * slot. This gives the desired
    * if it is not there already, but
    */
    if (row_on)
        R_CurRB-
        >rrb_StateData[dest_index] ^= (rgn_mask << RB_STATE_SIZE);
    dest_index++;
    /*
    * We now move onto the next row
    */
    row++;
    row_on = !row_on;

```

processing the first region. Therefore, we state of the run to the lowest bits.

processing the second region. Therefore, the state data copied to the previous state area so we state.

region is active at this transition, we current contents of the new region builder behaviour of making that region active turns it off if it is...

element.

```

                                continue;
                                }
                                /*
in x position to the current      * If the current row transition point is equal
                                * previous region builder transition point, we
advance the row counter to        * the next position.
                                */
                                if (*row == rb_run_start)
                                {
                                    row++;
                                    row_on = !row_on;
                                }
                                /*
transition region. We do similar  * Output the previous regions builder's
                                * things as for the region transition stuff
above.. Firstly though, we        * advance the rb_prev_run_state variable to the
next element. We know             * we can do this because if we were on the last
element, we wouldn't             * have hit this section of code.
                                */
                                rb_prev_run_state = r_PrevRB-
>rrb_StateData[src_index];
                                if (dest_index + 1 > R_CurRB->rrb_Size)
                                {
                                    if
(!r_grow_region_builder(R_CurRB, dest_index + 1))
                                        return FALSE;
                                }
                                R_CurRB->rrb_RgnData[dest_index] =
rb_run_start;
                                if (first)
                                {
                                    R_CurRB-
>rrb_StateData[dest_index] = (rb_prev_run_state & RB_CUR_STATE_MASK) |
                                (rb_prev_run_state >> RB_STATE_SIZE);
                                }
                                else
                                {
                                    R_CurRB-
>rrb_StateData[dest_index] = rb_prev_run_state;
                                }
                                if (!row_on)
                                {
                                    R_CurRB-
>rrb_StateData[dest_index] ^= (rgn_mask << RB_STATE_SIZE);
dest_index++;
                                }
                                /*
transitions. We now move          * We've output the previous region builder's
                                * over onto the next transition. If the previous
src_index increment              * has moved us onto the last element, we declare
that we have run
                                */

```

```

    * out of previous region builder data.
    */
    ASSERT(rb_run_start != R_EOR);
    if (++src_index >= r_PrevRB->rrb_Nels)
    {
        /*
         * We've run out of data..
         */
        rb_run_start = R_INT_MAX_VALUE -
2;
        continue;
    }
    /*
     * Otherwise, we still have stuff left to do, so
we move onto
     * the next run in the previous region builder.
     */
    rb_run_start = r_PrevRB-
>rrb_RgnData[src_index];
    }
    /*
     * Now, we simply blast out any remaining region builder
transition
     * points.
     */
    if (r_PrevRB->rrb_Nels - src_index > 0)
    {
        R_Int      nels_to_copy;
        nels_to_copy = r_PrevRB->rrb_Nels - src_index;
        if (dest_index + nels_to_copy > R_CurRB-
>rrb_Size)
        {
            if
(!r_grow_region_builder(R_CurRB, dest_index + nels_to_copy))
                return FALSE;
        }
        memcpy
        (
            R_CurRB->rrb_RgnData +
dest_index,
            r_PrevRB->rrb_RgnData +
src_index,
            nels_to_copy * sizeof(R_Int)
        );
        if (!first)
        {
            memcpy
            (
                R_CurRB-
>rrb_StateData + dest_index,
                r_PrevRB-
>rrb_StateData + src_index,
                nels_to_copy *
sizeof(unsigned char)
            );
            dest_index += nels_to_copy;
        }
    }
}
```

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```
#else

r_shift_and_dup[* (src - 8)];
#endif

#ifdef RB_USE_LOOKUP

    (* (src - 7) & RB_CUR_STATE_MASK) |

    (* (src - 7) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 7)];
#endif

#ifdef RB_USE_LOOKUP

    (* (src - 6) & RB_CUR_STATE_MASK) |

    (* (src - 6) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 6)];
#endif

#ifdef RB_USE_LOOKUP

    (* (src - 5) & RB_CUR_STATE_MASK) |

    (* (src - 5) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 5)];
#endif

#ifdef RB_USE_LOOKUP

    (* (src - 4) & RB_CUR_STATE_MASK) |

    (* (src - 4) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 4)];
#endif

#ifdef RB_USE_LOOKUP

    (* (src - 3) & RB_CUR_STATE_MASK) |

    (* (src - 3) >> RB_STATE_SIZE);
#else

r_shift_and_dup[* (src - 3)];

/* FALLTHROUGH!! */

case 7:

    *(dest - 7) =

/* FALLTHROUGH!! */

    *(dest - 7) =

/* FALLTHROUGH!! */

case 6:

    *(dest - 6) =

    *(dest - 6) =

/* FALLTHROUGH!! */

case 5:

    *(dest - 5) =

    *(dest - 5) =

/* FALLTHROUGH!! */

case 4:

    *(dest - 4) =

    *(dest - 4) =

/* FALLTHROUGH!! */

case 3:

    *(dest - 3) =

    *(dest - 3) =
```

```
#endif

case 2: /* FALLTHROUGH!! */

#ifdef RB_USE_LOOKUP
    *(dest - 2) =
    (*(src - 2) & RB_CUR_STATE_MASK) |
    (*(src - 2) >> RB_STATE_SIZE);
#else
    *(dest - 2) =
    r_shift_and_dup[*(src - 2)];
#endif

case 1: /* FALLTHROUGH!! */

#ifdef RB_USE_LOOKUP
    *(dest - 1) =
    (*(src - 1) & RB_CUR_STATE_MASK) |
    (*(src - 1) >> RB_STATE_SIZE);
#else
    *(dest - 1) =
    r_shift_and_dup[*(src - 1)];
#endif

case 0: /* FALLTHROUGH!! */
;
/* FALLTHROUGH!! */
}

}
/*
 * Finally, we set the number of elements of the latest region
 * builder. We also return the updates row variable.
 */
R_CurRB->rrb_Nels = dest_index;
*row_ptr = row;
return TRUE;
}
#else
int
R_add_row_to_region_builder
(
    R_Int      **row_ptr,
    int        rgn_mask,
    int        first
)
{
    R_Int      *row;
    R_Int      rb_run_start;
    unsigned char rb_prev_run_state;
    int
row_on;

    int
dest_index;
    unsigned char *src_state_ptr;
    unsigned char *dest_state_ptr;
    unsigned char *src_state_end_ptr;
```



```
register R_Int                                     *src_rgn_ptr;
R_Int                                              *src_rgn_end_ptr;
register R_Int                                     *dest_rgn_ptr;
R_Int                                              *dest_rgn_end_ptr;
int                                                inc;
int                                                i;

ASSERT(rgn_mask == 1 || rgn_mask == 2);

row = *row_ptr;
r_swap_region_builders();
/*
 * Skip over the row's y value at the beginning.
 */
ASSERT(*row == R_NEXT_IS_Y);
row += 2;
ASSERT(*row != R_NEXT_IS_Y && *row != R_EOR);

if (r_PrevRB->rrb_Nels == 0)
{
    /*
     * If the current region builder's src region
data array is empty, then
     * we are dealing with an empty region builder.
We simply convert
     * the input row to the region builder format.
     */
    row_on = TRUE;
    dest_index = 0;
    while (R_NOT_END_OF_ROW(*row))
    {
        if (++dest_index > R_CurRB->rrb_Size)
        {
            if
            (!r_grow_region_builder(R_CurRB, dest_index))
            return FALSE;
        }
        R_CurRB->rrb_RgnData[dest_index - 1] = *row;
        if (row_on)
            R_CurRB->rrb_StateData[dest_index - 1] =
            (rgn_mask << RB_STATE_SIZE);
        else
            R_CurRB->rrb_StateData[dest_index - 1] = 0;
        row_on = !row_on;
        row++;
    }
    *row_ptr = row;
    R_CurRB->rrb_Nels = dest_index;
    return TRUE;
}
/*
 * Firstly, we copy any runs from the region builder which
```

```
* precede this run from the region. We are checking the
* starting row against the start of each pixel run. Therefore
* we start checking against the 1st region builder data
* element.
*/
src_state_ptr = r_PrevRB->rrb_StateData;
src_rgn_ptr = r_PrevRB->rrb_RgnData;
src_state_end_ptr = src_state_ptr + r_PrevRB->rrb_Nels;
src_rgn_end_ptr = src_rgn_ptr + r_PrevRB->rrb_Nels;
dest_state_ptr = R_CurRB->rrb_StateData;
dest_rgn_ptr = R_CurRB->rrb_RgnData;
dest_rgn_end_ptr = dest_rgn_ptr + R_CurRB->rrb_Size;

ASSERT(r_PrevRB->rrb_Nels >= 2);
while (src_rgn_ptr != src_rgn_end_ptr && *row > *src_rgn_ptr)
    src_rgn_ptr++;
inc = src_rgn_ptr - r_PrevRB->rrb_RgnData;

if (inc > 0)
{
    src_state_ptr += inc;
    dest_state_ptr += inc;
    dest_rgn_ptr += inc;
    if (dest_rgn_ptr > dest_rgn_end_ptr)
    {
        if
        (!r_grow_region_builder(R_CurRB, inc))
            return FALSE;
        dest_state_ptr = R_CurRB-
        >rrb_StateData;
        dest_rgn_ptr = R_CurRB-
        >rrb_RgnData;
        dest_rgn_end_ptr = dest_rgn_ptr +
        R_CurRB->rrb_Size;
    }
    #if 1
    const R_Int * const
    src_rgn_ptr2 = src_rgn_ptr;
    R_Int
    *dest_rgn_ptr2 = dest_rgn_ptr;
    switch(inc)
    {
    default:
        for (i = inc; i > 10; i--)
        {
            *(dest_rgn_ptr2 -
            i) = *(src_rgn_ptr - i);
        }
        /* FALLTHROUGH!! */
        case 10:
            *(dest_rgn_ptr2 - 10) =
            *(src_rgn_ptr2 - 10);
        /* FALLTHROUGH!! */
        case 9:
            *(dest_rgn_ptr2 - 9) =
            *(src_rgn_ptr2 - 9);
        /* FALLTHROUGH!! */
        case 8:
    }
```

```

    *(src_rgn_ptr2 - 8);
    case 7:
        *(dest_rgn_ptr2 - 8) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 7) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 6) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 5) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 4) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 3) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 2) =
        /* FALLTHROUGH!! */
        *(dest_rgn_ptr2 - 1) =
        /* FALLTHROUGH!! */
        ;
        /* FALLTHROUGH!! */
    }
    memcpy
    (
        R_CurRB->rrb_RgnData,
        r_PrevRB->rrb_RgnData,
        inc * sizeof(R_Int)
    );
    if (!first)
    {
        memcpy
        (
            R_CurRB->rrb_StateData,
            r_PrevRB->rrb_StateData,
            inc *
            sizeof(unsigned char)
        );
    }
    else
    {
        switch (inc)

```

```

{
default:
    for (i = inc; i >
10; i--)
    {
#ifdef RB_USE_LOOKUP
        *(dest_state_ptr - i) = (*(src_state_ptr - i) & RB_CUR_STATE_MASK) |
        (*(src_state_ptr - i) >> RB_STATE_SIZE);
    }
#else
        *(dest_state_ptr - i) = r_shift_and_dup[*(src_state_ptr - i)];
    }
#endif
    /* FALLTHROUGH!! */
case 10:
#ifdef RB_USE_LOOKUP
        *(dest_state_ptr -
10) = (*(src_state_ptr - 10) & RB_CUR_STATE_MASK) |
        (*(src_state_ptr - 10) >> RB_STATE_SIZE);
    }
#else
        *(dest_state_ptr -
10) = r_shift_and_dup[*(src_state_ptr - 10)];
    }
#endif
    /* FALLTHROUGH!! */
case 9:
#ifdef RB_USE_LOOKUP
        *(dest_state_ptr -
9) = (*(src_state_ptr - 9) & RB_CUR_STATE_MASK) |
        (*(src_state_ptr - 9) >> RB_STATE_SIZE);
    }
#else
        *(dest_state_ptr -
9) = r_shift_and_dup[*(src_state_ptr - 9)];
    }
#endif
    /* FALLTHROUGH!! */
case 8:
#ifdef RB_USE_LOOKUP
        *(dest_state_ptr -
8) = (*(src_state_ptr - 8) & RB_CUR_STATE_MASK) |
        (*(src_state_ptr - 8) >> RB_STATE_SIZE);
    }
#else
        *(dest_state_ptr -
8) = r_shift_and_dup[*(src_state_ptr - 8)];
    }
#endif
    /* FALLTHROUGH!! */
case 7:
#ifdef RB_USE_LOOKUP
        *(dest_state_ptr -
7) = (*(src_state_ptr - 7) & RB_CUR_STATE_MASK) |
        (*(src_state_ptr - 7) >> RB_STATE_SIZE);
    }
#else
        *(dest_state_ptr -
7) = r_shift_and_dup[*(src_state_ptr - 7)];
    }
#endif
}
```



```
#ifndef RB_USE_LOOKUP
1) = (*(src_state_ptr - 1) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 1) >> RB_STATE_SIZE);
#else
1) = r_shift_and_dup(*(src_state_ptr - 1)];
#endif

case 0:
    /* FALLTHROUGH!! */
    ;
    /* FALLTHROUGH!! */
}

}
if (src_state_ptr == src_state_end_ptr)
{
    /*
    * We've already exhausted the previous region
    * of the next pixel run to be the max. possible
    * to be 0.
    */
    rb_run_start = R_INT_MAX_VALUE - 2;
    rb_prev_run_state = 0;
}
else
{
    /*
    * We are still within the previous region
    * the run info appropriately.
    */
    rb_run_start = *src_rgn_ptr;
    if (src_state_ptr == r_PrevRB->rrb_StateData)
        rb_prev_run_state = 0;
    else
        rb_prev_run_state =
*(src_state_ptr - 1);
}

/*
* We can now start dealing with the elements in the row.
*/
row_on = 1;
while (R_NOT_END_OF_ROW(*row))
{
    if (*row < rb_run_start)
    {
        if (dest_rgn_ptr + 1 >
dest_rgn_end_ptr)
        {
            if
(!r_grow_region_builder(R_CurRB, dest_rgn_ptr - R_CurRB->rrb_RgnData + 1))
return FALSE;

```

```
R_CurRB->rrb_StateData;
R_CurRB->rrb_RgnData;
dest_rgn_ptr + R_CurRB->rrb_Size;

}
*dest_rgn_ptr = *row;
if (first)
{
    /*
     * We are
     * copy the current
     * RB_STATE_SIZE
     */
    *dest_state_ptr =
        (rb_prev_run_state & RB_CUR_STATE_MASK) |
        (rb_prev_run_state >> RB_STATE_SIZE);
}
else
{
    /*
     * We are
     * has already been
     * just copy the
     */
    *dest_state_ptr =
        rb_prev_run_state;
}
/*
 * Now, if the row for the current
 * xor the region mask with the
 * slot. This gives the desired
 * if it is not there already, but
 */
if (row_on)
    *dest_state_ptr ^=
        (rgn_mask << RB_STATE_SIZE);
dest_state_ptr++;
dest_rgn_ptr++;
/*
 * We now move onto the next row
 */
row++;
row_on = !row_on;
continue;
```

processing the first region. Therefore, we state of the run to the lowest bits.

processing the second region. Therefore, the state data copied to the previous state area so we state.

region is active at this transition, we current contents of the new region builder behaviour of making that region active turns it off if it is...

element.

```

    }
    /*
    * If the current row transition point is equal
in x position to the current
    * previous region builder transition point, we
advance the row counter to
    * the next position.
    */
    if (*row == rb_run_start)
    {
        row++;
        row_on = !row_on;
    }
    /*
    * Output the previous regions builder's
transition region. We do similiar
    * things as for the region transition stuff
above.. Firstly though, we
    * advance the rb_prev_run_state variable to the
next element. We know
    * we can do this because if we were on the last
element, we wouldn't
    * have hit this section of code.
    */
    rb_prev_run_state = *src_state_ptr;
    if (dest_rgn_ptr + 1 > dest_rgn_end_ptr)
    {
        if
(!r_grow_region_builder(R_CurRB, dest_rgn_ptr - R_CurRB->rrb_RgnData + 1))
            return FALSE;
        dest_state_ptr = R_CurRB->rrb_StateData;
        dest_rgn_ptr = R_CurRB->rrb_RgnData;
        dest_rgn_end_ptr = dest_rgn_ptr +
R_CurRB->rrb_Size;
    }
    *dest_rgn_ptr = rb_run_start;
    if (first)
    {
        *dest_state_ptr =
(rb_prev_run_state & RB_CUR_STATE_MASK) |
(rb_prev_run_state >> RB_STATE_SIZE);
    }
    else
    {
        *dest_state_ptr =
rb_prev_run_state;
    }
    if (!row_on)
        *dest_state_ptr ^= (rgn_mask <<
RB_STATE_SIZE);
    dest_state_ptr++;
    dest_rgn_ptr++;
    /*
    * We've output the previous region builder's
transitions. We now move

```



```
src_index increment
that we have run

2;

we move onto

transition

src_state_ptr;
dest_rgn_end_ptr)

!r_grow_region_builder

R_CurRB,

dest_rgn_ptr - R_CurRB->rrb_RgnData + nels_to_copy

>rrb_StateData;
>rrb_RgnData;

* over onto the next transition. If the previous
* has moved us onto the last element, we declare
* out of previous region builder data.
*/
ASSERT(rb_run_start != R_EOR);
++src_rgn_ptr;
if (++src_state_ptr == src_state_end_ptr)
{
    /*
    * We've run out of data..
    */
    rb_run_start = R_INT_MAX_VALUE -
    continue;
}
/*
* Otherwise, we still have stuff left to do, so
* the next run in the previous region builder.
*/
rb_run_start = *src_rgn_ptr;
}
/*
* Now, we simply blast out any remaining region builder
* points.
*/
if (src_state_ptr != src_state_end_ptr)
{
    R_Int      nels_to_copy;
    nels_to_copy = src_state_end_ptr -
    if (dest_rgn_ptr + nels_to_copy >
    {
        if
        (
            (
                R_CurRB,
                dest_rgn_ptr - R_CurRB->rrb_RgnData + nels_to_copy
            )
        )
            return FALSE;
        dest_state_ptr = R_CurRB-
        dest_rgn_ptr = R_CurRB-
    }
    memcpy
    (
        dest_rgn_ptr,
        src_rgn_ptr,
```

```

                                nels_to_copy * sizeof(R_Int)
                                );
                                dest_rgn_ptr += nels_to_copy;
                                if (!first)
                                {
                                    memcpy
                                    (
                                        dest_state_ptr,
                                        src_state_ptr,
                                        nels_to_copy *
sizeof(unsigned char)
                                    );
                                }
                                else
                                {
                                    i = nels_to_copy;
                                    src_state_ptr =
                                        dest_state_ptr += nels_to_copy;
                                    switch (i)
                                    {
                                        default:
                                            for (; i > 10; i--)
                                            {
#ifdef RB_USE_LOOKUP
*(dest_state_ptr - i) = (*(src_state_ptr - i) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - i) >> RB_STATE_SIZE);
#else
*(dest_state_ptr - i) = r_shift_and_dup[*(src_state_ptr - i)];
#endif
                                            }
                                            /* FALLTHROUGH!! */
case 10:
#ifdef RB_USE_LOOKUP
10) = (*(src_state_ptr - 10) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 10) >> RB_STATE_SIZE);
#else
10) = r_shift_and_dup[*(src_state_ptr - 10)];
#endif
                                            /* FALLTHROUGH!! */
case 9:
#ifdef RB_USE_LOOKUP
9) = (*(src_state_ptr - 9) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 9) >> RB_STATE_SIZE);
#else
9) = r_shift_and_dup[*(src_state_ptr - 9)];
#endif
                                            /* FALLTHROUGH!! */
case 8:

```

```
#ifndef RB_USE_LOOKUP
8) = (*(src_state_ptr - 8) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 8) >> RB_STATE_SIZE);
#else
8) = r_shift_and_dup[*(src_state_ptr - 8)];
#endif

case 7:
#ifndef RB_USE_LOOKUP
7) = (*(src_state_ptr - 7) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 7) >> RB_STATE_SIZE);
#else
7) = r_shift_and_dup[*(src_state_ptr - 7)];
#endif

case 6:
#ifndef R_USE_LOOKUP
6) = (*(src_state_ptr - 6) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 6) >> RB_STATE_SIZE);
#else
6) = r_shift_and_dup[*(src_state_ptr - 6)];
#endif

case 5:
#ifndef RB_USE_LOOKUP
5) = (*(src_state_ptr - 5) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 5) >> RB_STATE_SIZE);
#else
5) = r_shift_and_dup[*(src_state_ptr - 5)];
#endif

case 4:
#ifndef RB_USE_LOOKUP
4) = (*(src_state_ptr - 4) & RB_CUR_STATE_MASK) |
    (*(src_state_ptr - 4) >> RB_STATE_SIZE);
#else
4) = r_shift_and_dup[*(src_state_ptr - 4)];
#endif

case 3:
#ifndef RB_USE_LOOKUP
3) = (*(src_state_ptr - 3) & RB_CUR_STATE_MASK) |
```

```
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
/* FALLTHROUGH!! */
*(dest_state_ptr -
```

```
(* (src_state_ptr - 3) >> RB_STATE_SIZE);
#else

3) = r_shift_and_dup[* (src_state_ptr - 3)];
#endif

case 2:

#ifndef RB_USE_LOOKUP

2) = (* (src_state_ptr - 2) & RB_CUR_STATE_MASK) |

(* (src_state_ptr - 2) >> RB_STATE_SIZE);
#else

2) = r_shift_and_dup[* (src_state_ptr - 2)];
#endif

case 1:

#ifndef RB_USE_LOOKUP

1) = (* (src_state_ptr - 1) & RB_CUR_STATE_MASK) |

(* (src_state_ptr - 1) >> RB_STATE_SIZE);
#else

1) = r_shift_and_dup[* (src_state_ptr - 1)];
#endif

case 0:

;

}

}

/*
 * Finally, we set the number of elements of the latest region
 * builder. We also return the updates row variable.
 */
R_CurRB->rrb_Nels = dest_rgn_ptr - R_CurRB->rrb_RgnData;
*row_ptr = row;
return TRUE;
}
#endif

/*
 * r_check_rgn_buf_len
 *
 * This function checks to see if the static region buffer is large enough.
 * If it isn't then it is reallocated to make it large enough.
 *
 * Parameters:
 *
 * size The required size of the r_RegBuf
array.
 * Returns:
 *
 * TRUE on success, FALSE on failure.
 */
static int
```

```
r_check_rgn_buf_len
(
    int size
)
{
    ASSERT(size >= 0);
    if (size > r_RgnBufSize)
    {
        int new_buf_size;
        R_Int *new_buf;
        new_buf_size = max(size, r_RgnBufSize * 2);
        new_buf = (R_Int *)malloc(new_buf_size *
sizeof(R_Int));

        if (new_buf == NULL)
            return FALSE;
        if (r_RgnBuf != NULL)
        {
            memcpy(new_buf, r_RgnBuf,
r_RgnBufSize * sizeof(R_Int));
            free(r_RgnBuf);
        }
        r_RgnBuf = new_buf;
        r_RgnBufSize = new_buf_size;
    }
    return TRUE;
}

/*
 * R_init_region_with_rect
 *
 * This function initialises a R_Region structure to represent a rectangular
 * region. It is assumed that the region is currently uninitialised.
 *
 * Parameters:
 *     rgn      A pointer to the R_Region to be initialised.
 *     rect     A pointer to an
IntXYMinMax structure representing
 *             the
rectangular area requiring an equivalent region
 *             description.
 * Returns:
 *             TRUE on success, FALSE on failure.
 */
int
R_init_region_with_rect
(
    R_Region *rgn,
    IntXYMinMax *rect
)
{
    R_Int *rgn_data;

    ASSERT(rect->X.Min <= rect->X.Max);
    ASSERT(rect->Y.Min <= rect->Y.Max);

    rgn->rr_BBox = *rect;
    rgn_data = (R_Int *)malloc(9 * sizeof(R_Int));
```

```
        if (rgn_data == NULL)
        {
            return FALSE;
        }
        rgn_data[0] = R_NEXT_IS_Y;
        rgn_data[1] = rect->Y.Min;
        rgn_data[2] = rect->X.Min;
        rgn_data[3] = rect->X.Max + 1;
        rgn_data[4] = R_NEXT_IS_Y;
        rgn_data[5] = rect->Y.Max + 1;
        rgn_data[6] = rect->X.Min;
        rgn_data[7] = rect->X.Max + 1;
        rgn_data[8] = R_EOR;
        rgn->rr_RgnData = rgn_data;
        rgn->rr_RgnDataSize = 9;
        return TRUE;
    }

/*
 * R_region_with_region
 *
 * This function initialises a R_Region structure to represent a the region
 * passed as an argument. It is assumed that the region is currently
 * uninitialised.
 *
 * Parameters:
 *     rgn          A pointer to the R_Region to be initialised.
 *     src_rgn      A pointer to an
R_Region structure representing
 *
 * the
region to which this region is to be initialised.
 * Returns:
 *
 *     TRUE on success, FALSE on failure.
 */
int
R_init_region_with_region
(
    R_Region    *rgn,
    R_Region    *src_rgn
)
{
    R_Int        *rgn_data;

    rgn->rr_BBox = src_rgn->rr_BBox;
    rgn_data = (R_Int *)malloc(src_rgn->rr_RgnDataSize *
sizeof(R_Int));
    if (rgn_data == NULL)
    {
        return FALSE;
    }
    memcpy
    (
        rgn_data,
        src_rgn->rr_RgnData,
        src_rgn->rr_RgnDataSize * sizeof(R_Int)
    );
    rgn->rr_RgnData = rgn_data;
    rgn->rr_RgnDataSize = src_rgn->rr_RgnDataSize;
}
```

```
        return TRUE;
    }

/*
 * R_region_with_translated_region
 *
 * This function initialises a R_Region structure to represent a the region
 * passed as an argument translated by delta. It is assumed that the region
 * is currently uninitialised.
 *
 * Parameters:
 *     rgn          A pointer to the R_Region to be initialised.
 *     src_rgn      A pointer to an
R_Region structure representing
 *
 * the
region to which this region is to be initialised.
 *     delta        A pointer to a IntXY structure representing the
 *                  translation required.
 * Returns:
 *
 *                 TRUE on success, FALSE on failure.
 */
int
R_init_region_with_translated_region
(
    R_Region    *rgn,
    R_Region    *src_rgn,
    IntXY       *delta
)
{
    R_Int        *rgn_data;
    R_Int        *src_data;

    rgn->rr_BBox.X.Min = src_rgn->rr_BBox.X.Min + delta->X;
    rgn->rr_BBox.X.Max = src_rgn->rr_BBox.X.Max + delta->X;
    rgn->rr_BBox.Y.Min = src_rgn->rr_BBox.Y.Min + delta->Y;
    rgn->rr_BBox.Y.Max = src_rgn->rr_BBox.Y.Max + delta->Y;
    rgn_data = (R_Int *)malloc(src_rgn->rr_RgnDataSize *
sizeof(R_Int));
    if (rgn_data == NULL)
    {
        return FALSE;
    }
    src_data = src_rgn->rr_RgnData;
    for (int i = 0; i < src_rgn->rr_RgnDataSize; i++)
    {
        if (src_data[i] == R_NEXT_IS_Y)
        {
            rgn_data[i] = src_data[i];
            i++;
            rgn_data[i] = src_data[i] + delta->Y;
            continue;
        }
        else if (src_data[i] == R_EOR)
            rgn_data[i] = src_data[i];
        else
            rgn_data[i] = src_data[i] + delta->X;
    }
}
```

```
    rgn->rr_RgnData = rgn_data;
    rgn->rr_RgnDataSize = src_rgn->rr_RgnDataSize;
    return TRUE;
}

/*
 * R_empty_region
 *
 * Deallocates the region data allocated for a region. Only the
 * data is freed. The R_Region structure itself is not.
 *
 * Parameters:
 *          rgn                The region whose
region data is to be deallocated.
 * Returns:
 *          Nothing.
 */
void
R_empty_region
(
    R_Region    *rgn
)
{
    if (rgn != NULL && rgn->rr_RgnData != NULL)
    {
        free(rgn->rr_RgnData);
        rgn->rr_RgnData = NULL;
    }
}

#ifdef R_USE_NEW_IMP
/*
 * R_union
 *
 * This function inits a R_Region structure to represent the union
 * of it's two arguments.
 *
 * Parameters:
 *          rgn                The R_Region to be initialised.
 *          r1                A R_Region ptr
representing the first region.
 *          r2                A R_Region ptr
representing the second region.
 * Returns
 *          TRUE on success, FALSE on failure.
 */
int
R_union
(
    R_Region    *rgn,
    R_Region    *r1,
    R_Region    *r2
)
{
    R_Int        *r1_dat;
    R_Int        *r2_dat;
    int          overlap_flags;

    union_tot++;
}
```



```
        if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
        {
            /*
            * The bounding boxes don't intersect. This means
            * union very easily, simply by copying data from
            * We malloc a new region data array of size r1-
            * r2->rr_RgnDataSize - 1. This is the maximum
            * resulting region. Not all of this memory will
            * the two regions being combined have rows with
            * (R_NEXT_IS_Y marker is not duplicated).
            */
            rgn->rr_RgnDataSize = r1->rr_RgnDataSize + r2-
            rgn->rr_RgnData = (R_Int *)malloc(rgn-
            sizeof(R_Int));

            if (rgn->rr_RgnData == NULL)
            {
                return FALSE;
            }
            /*
            * Now, check to see if the regions overlap in
            y...
            */
            if (!(overlap_flags & BB_INTERSECT_OVERLAP_Y))
            {
                /*
                * The regions don't overlap in y.
                * and then another into the array
                * that r1 points to the region
                */
                if (r2->rr_BBox.Y.Min < r1-
                {
                    R_Region
                    tmp = r1;
                    r1 = r2;
                    r2 = tmp;
                }
                memcpy
                (
                    rgn->rr_RgnData,
                    r1->rr_RgnData,
                    (r1-
                    >rr_RgnDataSize - 1) * sizeof(R_Int)
                )
            }
        }
    }
    *tmp;
    >rr_BBox.Y.Min)
    *tmp;
    >rr_RgnDataSize - 1) * sizeof(R_Int)
```

```

r1->rr_RgnDataSize - 1,

* sizeof(R_Int)

>rr_RgnDataSize - 1] == R_EOR);
    }
    else
    {

r1_done;

r1_consumed;

r2_consumed;

num_written;

in x. We simply go row
memcpy the individual rows as
points to the region with

>rr_BBox.X.Min)

*tmp;

!= R_EOR)

R_NEXT_IS_Y);

R_NEXT_IS_Y);

```

```

);
memcpy
(
    rgn->rr_RgnData +
    r2->rr_RgnData,
    r2->rr_RgnDataSize
);
ASSERT(rgn->rr_RgnData[rgn-

R_Int      *r1_tmp;
R_Int      *r2_tmp;
R_Int      *dest;
R_Int      min_row;
int

int

int

int

/*
 * The bboxes overlap in y but not
 * by row through each region and
 * appropriate. We ensure that r1
 * the smallest x coordinate.
 */
if (r2->rr_BBox.X.Min < r1-
{
    R_Region

    tmp = r1;
    r1 = r2;
    r2 = tmp;
}
r1_dat = r1->rr_RgnData;
r2_dat = r2->rr_RgnData;
dest = rgn->rr_RgnData;
rgn->rr_RgnDataSize = 0;
r1_consumed = 0;
r2_consumed = 0;
while (*r1_dat != R_EOR && *r2_dat
{
    ASSERT(*r1_dat ==
    ASSERT(*r2_dat ==

```

```
min(r1_dat[1], r2_dat[1]);
```

```
min_row)
```

need to emit r1. We therefore need to find where the next row (if any) starts. When we do this we recall that a y value _must be followed by at least two x values..

```
r1_tmp = r1_dat + 4;
```

```
while (*r1_tmp != R_NEXT_IS_Y && *r1_tmp != R_EOR)
```

```
r1_tmp++;
```

```
num_written = r1_tmp - r1_dat;
```

```
memcpy(dest, r1_dat, num_written * sizeof(R_Int));
```

```
+= num_written;
```

```
r1_consumed += num_written;
```

```
>rr_RgnDataSize += num_written;
```

```
r1_dat = r1_tmp;
```

```
r1_done = TRUE;
```

```
min_row)
```

need to emit r1. We therefore need to find where the next row (if any) starts. When we do this we recall that a y value _must be followed by at least two x values. If r1's current row has already been emitted for this y value, we do _not_ emit the R_NEXT_IS_Y marker or the y value itself.

```
(r1_done)
```

```
min_row =
```

```
r1_done = FALSE;
```

```
if (r1_dat[1] ==
```

```
{
```

```
/*  
 * We  
 *  
 *  
 *  
 */
```

```
dest
```

```
rgn-
```

```
}
```

```
if (r2_dat[1] ==
```

```
{
```

```
/*  
 * We  
 *  
 *  
 *  
 *  
 */  
if  
  
{
```

```

    r2_dat += 2;

    r2_tmp = r2_dat + 2;

    r2_consumed += 2;

}
else
{

    r2_tmp = r2_dat + 4;

}

while (*r2_tmp != R_NEXT_IS_Y && *r2_tmp != R_EOR)

    r2_tmp++;

num_written = r2_tmp - r2_dat;

memcpy(dest, r2_dat, num_written * sizeof(R_Int));

+= num_written;

r2_consumed += num_written;

>rr_RgnDataSize += num_written;

r2_dat = r2_tmp;

}

if (*r1_dat != R_EOR)
{
    /*
     * r1 is the last
     * the remainder of
     * R_EOR marker) to
     */
    ASSERT(r2_consumed
    memcpy
    (
        dest,
        r1_dat,
        >rr_RgnDataSize - r1_consumed) * sizeof(R_Int)
    );
    rgn-
    >rr_RgnDataSize += (r1->rr_RgnDataSize - r1_consumed);
}
else
{
    /*

```

```
region left standing. We memcpy
the region (including the
the destination.

== r1->rr_RgnDataSize - 1);

dest,
r2_dat,
>rr_RgnDataSize - r2_consumed) * sizeof(R_Int)

>rr_RgnDataSize += (r2->rr_RgnDataSize - r2_consumed);
    }
    ASSERT
    (
    rgn-
>rr_RgnData[rgn->rr_RgnDataSize - 1] == R_EOR
    );
    }
    }
    else
    {
        R_Int
        min_row;
        dest_size;
        unsigned char *rgn_bld_stat;
        R_Int
        *rgn_bld_dat;
        int
        int
        in_run;
        int
        done_r1_in_row;

        union_full++;

        /*
        * The two regions _do_ overlap in x _and_ y. We
        * to do a bit more work in calculating the union
        * regions. We use the R_RegionBuilder struct to
        * regarding the currently active regions as we
        * the rows of each region. After any rows

therefore have
of the two
store state
progress through
relevent to a y-coord
```

```
== min row)
```

```

is active on this y coord. We add this
current region builder.

(!R_add_row_to_region_builder(&r2_dat, 0x2, !done_r1_in_row))
return FALSE;

for the input rows.

(!r_check_rgn_buf_len(dest_size + 2))

R_NEXT_IS_Y;

>rrb_StateData;

>rrb_RgnData;

i--)

*rgn_bld_stat > 0

(*rgn_bld_stat & RB_CUR_STATE_MASK) == 0

(*rgn_bld_stat & RB_PREV_STATE_MASK) == 0

have to emit a run here, if we're not already
one..

(!in_run)

(!r_check_rgn_buf_len(dest_size + 1))

return FALSE;

* The first region
* row to the
*/
if
}
/*
* Now, we generate the output row
*/
if
{
    return FALSE;
}
r_RgnBuf[dest_size++] =
r_RgnBuf[dest_size++] = min_row;
rgn_bld_stat = R_CurRB-
rgn_bld_dat = R_CurRB-
in_run = FALSE;
for (i = R_CurRB->rrb_Nels; i > 0;
{
    if
    (
        &&
        (
            ||
        )
    )
    {
        /*
        * We
        * in
        */
        if
        {
            if
            {

```

[0:\cisra\openscrn] Region.cpp


```

        ASSERT(rgn->rr_RgnDataSize >= 9);
    }
    /*
    * We now do a bounding box union of the two component bboxes
and place
    * the result in the new region.
    */
    BB_union(&r1->rr_BBox, &r2->rr_BBox, &rgn->rr_BBox);
    /*
    * Done! We can get out..
    */
    return TRUE;
}

/*
* R_union_equals
*
* This function basically implements a r1 union= r2 type operation. Ie
* r1 union r2 is calculated and the result returned in r1.
*
* Parameters:
*
*           r1                A pointer to an
R_Region. This represents
*
*           the first half of
the union, and is also used to return
*
*           the eventual
result.
*
*           r2                A pointer to an
R_Region. This represents the second
*
*           half of the union.
* Returns:
*
*           TRUE on success, FALSE on failure.
*/
int
R_union_equals
(
    R_Region    *r1,
    R_Region    *r2
)
{
    R_Region    new_rgn;
    if (r1->rr_RgnData == NULL)
        return R_init_region_with_region(r1, r2);
    if (!R_union(&new_rgn, r1, r2))
        return FALSE;
    R_empty_region(r1);
    *r1 = new_rgn;
    return TRUE;
}

/*
* R_intersection
*
* This function inits a R_Region structure to represent the intersection
* of it's two arguments.
*
* Parameters:
*
*   rgn        A R_Region ptr to the R_Region structure to be initialised.

```

```

*                               r1                               A R_Region ptr
representing the first region.
*                               r2                               A R_Region ptr
representing the second region.
* Returns
*                               TRUE on success, FALSE on failure.
*/
int
R_intersection
(
    R_Region    *rgn,
                R_Region    *r1,
                R_Region    *r2
)
{
    R_Int      *r1_dat;
    R_Int      *r2_dat;
    int
                                overlap_flags;

    int_tot++;

    rgn->rr_RgnData = NULL;
    if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
    {
        /*
        * The bounding boxes don't intersect. This means
        * don't intersect. Therefore, we simply set rgn-
        * (signifying an empty region) and get out..
        */
        return TRUE;
    }
    R_Int
min_row;
    int
dest_size;
    unsigned char    *rgn_bld_stat;
    R_Int
*rgn_bld_dat;
    int
                                i;
    int
in_run;
    int
done_r1_in_row;
    IntXYMinMax
                                new_bbox;

    int_full++;

    /*
    * The two regions _do_ overlap in x_and y. We therefore have
    * to do a bit more work in calculating the intersection of the
two
    * regions. We use the R_RegionBuilder struct to store state
    * regarding the currently active regions as we progress
through

```

```
* the rows of each region. After any rows relevent to a y-coord
* are added to the region builder, we examine the state of each
* pixel run in the region builder. If the addition of the
row(s)
* for the y-coord have caused a transition to or from 0x3, then
* the pixel run is emitted.
*/
/*
* Initialise the new_bbox structure for determining the new
bounding box.
*/
new_bbox.X.Min = R_INT_MAX_VALUE;
new_bbox.Y.Min = R_INT_MAX_VALUE;
new_bbox.X.Max = R_INT_MIN_VALUE;
new_bbox.Y.Max = R_INT_MIN_VALUE;

/*
* The next thing we do is ensure the current region builder
is empty,
* and set up pointers into the region data of the two regions.
*/
r1_dat = r1->rr_RgnData;
r2_dat = r2->rr_RgnData;
R_CurRB->rrb_Nels = 0;
dest_size = 0;
/*
* We are now ready to loop through the data from both regions.
Notice
* that we only keep looping whilst _both_ regions have some
data left
* to give. As soon as either of the region's data has been
exhausted,
* then we stop as the intersection region has already been
calculated
* and is sitting in the rgn_buf.
*/
while (*r1_dat != R_EOR && *r2_dat != R_EOR)
{
    ASSERT(*r1_dat == R_NEXT_IS_Y || *r1_dat ==
R_EOR);
    ASSERT(*r2_dat == R_NEXT_IS_Y || *r2_dat ==
R_EOR);

    if (*r1_dat == R_EOR)
        min_row = r2_dat[1];
    else if (*r2_dat == R_EOR)
        min_row = r1_dat[1];
    else
        min_row = min(r1_dat[1],
r2_dat[1]);

    done_r1_in_row = FALSE;
    if (*r1_dat != R_EOR && r1_dat[1] == min_row)
    {
        /*
        * The first region is active on
this y coord. We add this
builder.
        * row to the current region
        */

```

```

        if
(!R_add_row_to_region_builder(&r1_dat, 0x1, TRUE))
            return FALSE;
        done_r1_in_row = TRUE;
    }
    if (*r2_dat != R_EOR && r2_dat[1] == min_row)
    {
        /*
        * The first region is active on
        * row to the current region
        */
        if
(!R_add_row_to_region_builder(&r2_dat, 0x2, !done_r1_in_row))
            return FALSE;
    }
    /*
    * Now, we generate the output row for the input
    rows.
    */
    if (!r_check_rgn_buf_len(dest_size + 2))
    {
        return FALSE;
    }
    r_RgnBuf[dest_size++] = R_NEXT_IS_Y;
    r_RgnBuf[dest_size++] = min_row;
    rgn_bld_stat = R_CurRB->rrb_StateData;
    rgn_bld_dat = R_CurRB->rrb_RgnData;
    in_run = FALSE;
    for (i = R_CurRB->rrb_Nels; i > 0; i--)
    {
        if
        (
            *rgn_bld_stat != (3
| (3 << RB_STATE_SIZE))
            &&
            (
                (*rgn_bld_stat & RB_PREV_STATE_MASK) == 3
                ||
                (*rgn_bld_stat & RB_CUR_STATE_MASK) == (3 << RB_STATE_SIZE))
            )
        {
            /*
            * We have to emit
            * in one..
            */
            if (!in_run)
            {
                if
                {
                    return FALSE;
                }
            }
        }
    }

```

```

r_RgnBuf[dest_size++] = *rgn_bld_dat;

in_run = TRUE;

new_bbox.X.Min = min(new_bbox.X.Min, *rgn_bld_dat);
    }
    else
    {
        if (in_run)
        {
            /*
            *
            * We've come to the end of a run. We output the next element to end it.
            */
            if
            (!r_check_rgn_buf_len(dest_size + 1))
            {
                return FALSE;
            }

            r_RgnBuf[dest_size++] = *rgn_bld_dat;

            new_bbox.X.Max = max(new_bbox.X.Max, *rgn_bld_dat);
        }
        in_run = FALSE;
    }
    rgn_bld_stat++;
    rgn_bld_dat++;
}
if (r_RgnBuf[dest_size - 2] == R_NEXT_IS_Y)
{
    /*
    * We didn't output anything for
    */
    dest_size -= 2;
}
else
{
    if (min_row < new_bbox.Y.Min)
        new_bbox.Y.Min =
min_row;

    else if (min_row >
new_bbox.Y.Max)
        new_bbox.Y.Max =
min_row;
}
}
/*
* We've completed constructing the data for the region.
Firstly
* we check to see if we've emitted anything at all. If we have
* then dest_size must be > 0. If it isn't we simply free the
* region we created and get out, as the regions don't really
* intersect, in spite of their intersecting bounding boxes.

```

```
        */
        if (dest_size == 0)
        {
                return TRUE;
        }
        /*
        * We make a copy the constructed data from the permanent buffer
to
        * an exactly fitting buffer.
        */
        rgn->rr_RgnData = (R_Int *)malloc(++dest_size *
sizeof(R_Int));
        if (rgn->rr_RgnData == NULL)
        {
                return FALSE;
        }
        memcpy(rgn->rr_RgnData, r_RgnBuf, (dest_size - 1) *
sizeof(R_Int));
        rgn->rr_RgnData[dest_size - 1] = R_EOR;
        rgn->rr_RgnDataSize = dest_size;
        ASSERT(rgn->rr_RgnDataSize >= 9);
        /*
        * Now, copy across the bounding box.. Before we do this, we
subtract
        * 1 from X.Max and Y.Max because of the region format.
        */
        new_bbox.X.Max--;
        new_bbox.Y.Max--;
        rgn->rr_BBox = new_bbox;
        /*
        * Done! We can get out..
        */
        return TRUE;
}

/*
 * R_difference
 *
 * This function inits a R_Region structure to represent the difference of
 * it's two arguments. It essentially calculates r1 - r2
 *
 * Parameters:
 *      rgn      A R_Region ptr representing the R_Region to be initied.
 *              r1      A R_Region ptr
representing the first region.
 *              r2      A R_Region ptr
representing the second region.
 * Returns
 *
 *              TRUE on success, FALSE on failure.
 */
int
R_difference
(
        R_Region      *rgn,
        R_Region      *r1,
        R_Region      *r2
)
{
```

```
R_Int          *r1_dat;
R_Int          *r2_dat;
int            overlap_flags;

diff_tot++;

rgn->rr_RgnData = NULL;
if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
{
    /*
    * The bounding boxes don't intersect. This means
    that r1 - r2
    * simply equals r1. We make a copy of the
    relevant bits and get out..
    */
    rgn->rr_BBox = r1->rr_BBox;
    rgn->rr_RgnDataSize = r1->rr_RgnDataSize;
    rgn->rr_RgnData = (R_Int *)malloc(r1-
>rr_RgnDataSize * sizeof(R_Int));
    if (rgn->rr_RgnData == NULL)
    {
        return FALSE;
    }
    memcpy
    (
        rgn->rr_RgnData,
        r1->rr_RgnData,
        r1->rr_RgnDataSize *
sizeof(R_Int)
    );
    return TRUE;
}
R_Int          min_row;
int
dest_size;
unsigned char *rgn_bld_stat;
R_Int          *rgn_bld_dat;
int            i;
int
in_run;
int
done_r1_in_row;
unsigned char m_high;
unsigned char m_low;
IntXYMinMax   new_bbox;

diff_full++;

/*
* The two regions _do_ overlap in x _and_ y. We therefore have
* to do a bit more work in calculating the difference of the
two
* regions. We use the R_RegionBuilder struct to store state
* regarding the currently active regions as we progress
through
* the rows of each region. After any rows relevent to a y-coord
```

```

* are added to the region builder, we examine the state of each
* pixel run in the region builder. If the addition of the
row(s)
* for the y-coord have caused the following transitions -
*
*           r1      -> 0
*           0      -> r2
*           r1 + r2 -> r2
*           r2      -> r1 + r2
* ..then the relevent runs are emitted. Firstly, though,
* we ensure the current region builder is empty,
* and set up pointers into the region data of the two regions.
*/
r1_dat = r1->rr_RgnData;
r2_dat = r2->rr_RgnData;
R_CurRB->rrb_Nels = 0;
dest_size = 0;

/*
* Initialise the new_bbox structure for determining the new
bounding box.
*/
new_bbox.X.Min = 32767;
new_bbox.Y.Min = 32767;
new_bbox.X.Max = -32768;
new_bbox.Y.Max = -32768;

/*
* We are now ready to loop through the data from both regions.
Notice
* that we only keep looping whilst r1 has data outstanding.
When
* r1's data is consumed, then any transitions made by r2 are
* irrelevant.
*/
while (*r1_dat != R_EOR)
{
    ASSERT(*r1_dat == R_NEXT_IS_Y || *r1_dat ==
R_EOR);
    ASSERT(*r2_dat == R_NEXT_IS_Y || *r2_dat ==
R_EOR);

    if (*r1_dat == R_EOR)
        min_row = r2_dat[1];
    else if (*r2_dat == R_EOR)
        min_row = r1_dat[1];
    else
        min_row = min(r1_dat[1],
r2_dat[1]);

    done_r1_in_row = FALSE;
    if (*r1_dat != R_EOR && r1_dat[1] == min_row)
    {
        /*
        * The first region is active on
this y coord. We add this
builder.
        * row to the current region
        */
        if
(!R_add_row_to_region_builder(&r1_dat, 0x1, TRUE))

```



```

                                return FALSE;
                                done_r1_in_row = TRUE;
                                }
                                if (*r2_dat != R_EOR && r2_dat[1] == min_row)
                                {
                                    /*
                                    * The first region is active on
this y coord. We add this
builder.
                                    * row to the current region
                                    */
                                    if
(!R_add_row_to_region_builder(&r2_dat, 0x2, !done_r1_in_row))
                                        return FALSE;
                                }
                                /*
                                * Now, we generate the output row for the input
rows.
                                */
                                if (!r_check_rgn_buf_len(dest_size + 2))
                                {
                                    return FALSE;
                                }
                                r_RgnBuf[dest_size++] = R_NEXT_IS_Y;
                                r_RgnBuf[dest_size++] = min_row;
                                rgn_bld_stat = R_CurRB->rrb_StateData;
                                rgn_bld_dat = R_CurRB->rrb_RgnData;
                                in_run = FALSE;
                                for (i = R_CurRB->rrb_Nels; i > 0; i--)
                                {
                                    m_high = (*rgn_bld_stat &
RB_CUR_STATE_MASK) >> RB_STATE_SIZE;
                                    m_low = *rgn_bld_stat &
RB_PREV_STATE_MASK;
                                    if
                                    (
                                        (
(m_low != 1 && m_high == 1)
                                        ||
(m_low == 1 && m_high != 1)
                                        )
                                    )
                                    {
                                        /*
                                        * We have to emit
a run here, if we're not already
                                        * in one..
                                        */
                                        if (!in_run)
                                        {
                                            if
                                            {
return FALSE;
                                            }
                                        }
                                    }
                                }

```

```
r_RgnBuf[dest_size++] = *rgn_bld_dat;

in_run = TRUE;

new_bbox.X.Min = min(new_bbox.X.Min, *rgn_bld_dat);
    }
    else
    {
        if (in_run)
        {
            /*
             *
            */
            We've come to the end of a run. We output the next element to end it.
            */
            (!r_check_rgn_buf_len(dest_size + 1))
            if
            {
                return FALSE;
            }

            r_RgnBuf[dest_size++] = *rgn_bld_dat;

            new_bbox.X.Max = max(new_bbox.X.Max, *rgn_bld_dat);
        }
        in_run = FALSE;
    }
    rgn_bld_stat++;
    rgn_bld_dat++;
}
if (r_RgnBuf[dest_size - 2] == R_NEXT_IS_Y)
{
    /*
     * We didn't output anything for
    */
    these input rows. Rewind..
    dest_size -= 2;
}
else
{
    if (min_row < new_bbox.Y.Min)
        new_bbox.Y.Min =
min_row;

    else if (min_row >
new_bbox.Y.Max)
        new_bbox.Y.Max =
min_row;
}
}
/*
 * We've completed constructing the data for the region.
Firstly
 * we check to see if we've emitted anything at all. If we have
 * then dest_size must be > 0. If it isn't we simply free the
 * region we created and get out, as r2 - r1 must be empty.
 */
```

```
        if (dest_size == 0)
        {
            return TRUE;
        }
        /*
        * We make a copy the constructed data from the permanent buffer
to
        * an exactly fitting buffer.
        */
        rgn->rr_RgnData = (R_Int *)malloc(++dest_size *
sizeof(R_Int));
        if (rgn->rr_RgnData == NULL)
        {
            return FALSE;
        }
        memcpy(rgn->rr_RgnData, r_RgnBuf, (dest_size - 1) *
sizeof(R_Int));
        rgn->rr_RgnData[dest_size - 1] = R_EOR;
        rgn->rr_RgnDataSize = dest_size;
        ASSERT(rgn->rr_RgnDataSize >= 9);
        /*
        * Now, copy across the bounding box..
        */
        rgn->rr_BBox = new_bbox;
        /*
        * Done! We can get out..
        */
        return TRUE;
    }
#endif /* R_USE_NEW_IMP */
/*
* r_grow_free_list
*
* This function mallocs and adds R_FREE_LIST_GROWTH_SIZE new elements
* to the front of the region growth free list.
*
* Parameters:
*
* None.
*
* Returns:
*
* TRUE on success, FALSE on failure.
*/
int
r_grow_free_list()
{
    R_RgnGrowItem *rgi;
    int i;
    /*
    * First, malloc the memory..
    */
    rgi = (R_RgnGrowItem *)malloc
        (
            R_FREE_LIST_GROWTH_SIZE *
sizeof(R_RgnGrowItem)
        );
    if (rgi == NULL)
        return FALSE;
    /*
    * Now make the whole block of memory into a list..
    */
}
```



```

    */
    for (i = 0; i < R_FREE_LIST_GROWTH_SIZE - 1; i++)
        rgi[i].rrgi_Next = &rgi[i + 1];
    /*
    * Now, add it to the front of the free list..
    */
    rgi[R_FREE_LIST_GROWTH_SIZE - 1].rrgi_Next = r_free_list;
    r_free_list = rgi;
    return TRUE;
}

/*
* R_add_row_to_region_growth_list
*
* This function adds a row from a R_Region to a linked list comprised of
* R_RgnGrowItem structures. "Adding" implies that the linked list is
* modified such that the state and coordinate information present
* in the list is updated to take into account the new row just added.
* Adding a row has the following properties:
*
*      * If a pixel run in the row does not exist in the
list before,
*
*      it is added and it's current state is tagged
with the region
*
*      to which the row belongs. The previous state
is set to 0,
*
*      indicating that it did not exist before.
*
*      * If a pixel run in the row did exist before, but
it's present state
*
*      indicates that it came from the other region
then the run
*
*      is retained but it's state is modified to
indicate that
*
*      both regions are active at this point.
*
*      * If a pixel run in the row did exist before, and
it's present
*
*      state indicates that the current region then
the region is
*
*      removed and it's state is modified to indicate
that the run is
*
*      now empty.
*
*      * If a pixel run in the row did exist before, and
it's present state
*
*      indicates that both regions are currently
active then the run
*
*      is retained, but its state is modified to
indicate that only the
*
*      other region is active in this run.
*
* Parameters:
*
*      row_ptr          A R_Int ** pointer
to the row in the region. Used
*
*      to
return the updates row pointer.
*
*      rgn_mask         A mask for the region the row
comes from. Must
*
*      be
either 1 or 2.

```

```
*
first                                Whether this is the
first region to be processed
*
the current scanline.                on
* Returns:
*
TRUE on success, FALSE on failure.
*/
int
R_add_row_to_region_growth_list
(
    R_Int      **row_ptr,
    int         rgn_mask,
    int         first
)
{
    R_Int      *row;
    R_RgnGrowItem *rgi;
    unsigned char rb_prev_run_state;
    int
row_on;

    row = *row_ptr;
    /*
     * Skip over the row's y value at the beginning.
     */
    ASSERT(*row == R_NEXT_IS_Y);
    row += 2;
    ASSERT(*row != R_NEXT_IS_Y && *row != R_EOR);

    if (r_growth_list == NULL)
    {
        R_RgnGrowItem **ptr_next_ptr;
        /*
         * The growth list is currently empty. Therefore,
         * the input row to the region growth list
         *
         */
        row_on = TRUE;
        ptr_next_ptr = &r_growth_list;
        while (R_NOT_END_OF_ROW(*row))
        {
            if (r_free_list == NULL)
            {
                if
(!r_grow_free_list())

                return FALSE;

            }
            rgi = r_free_list;
            *ptr_next_ptr = rgi;
            ptr_next_ptr = &rgi->rrgi_Next;
            r_free_list = *ptr_next_ptr;
            rgi->rrgi_RgnData = *row;
            if (row_on)
                rgi->rrgi_StateData = (rgn_mask << RB_STATE_SIZE);
            else
                return FALSE;
        }
    }
}
```

```

>rrgi_StateData = 0;

rgi-
    row_on = !row_on;
    row++;
}
*row_ptr = row;
*ptr_next_ptr = NULL;
return TRUE;
}

R_RgnGrowItem fake_item;
R_RgnGrowItem *prev_rgi;

/*
 * "fake_item" is used as the head of the list. This is so that
we _always_ have
 * a valid pointer to the previous item in the list. Only the
next pointer and
 * state data are initialised, as these are they only elements
which will be
 * referenced.
 */
fake_item.rrgi_StateData = 0;
fake_item.rrgi_Next = r_growth_list;
prev_rgi = &fake_item;

if (first)
{
    /*
     * If this is the first row to be added on this
particular scanline,
     * then we have to update the existing contents
of those elements
     * at the beginning of the growth list which
precede (in coords) the
     * first element of the row. "Updating" involves
updating the
     * previous state of each element to match the
current state. This is
     * because none of the elements were effected by
the addition of the
     * new row.
     */
    rgi = r_growth_list;
    while (rgi != NULL && *row > rgi->rrgi_RgnData)
    {
        rgi->rrgi_StateData = (rgi-
>rrgi_StateData & RB_CUR_STATE_MASK) |
        (rgi->rrgi_StateData >> RB_STATE_SIZE);
    }
    rgi->rrgi_StateData =
    r_shift_and_dup[rgi->rrgi_StateData];
    prev_rgi = rgi;
    rgi = rgi->rrgi_Next;
}
}

```

```

    }
    else
    {
        /*
        * This is the second row to be added on this
        * Therefore, we don't need to update the state
        * preceding (in coords) the first run of the row
        * they have already been updated by the first
        * this scanline. We simply skip over the
        */
        rgi = r_growth_list;
        while (rgi != NULL && *row > rgi->rrgi_RgnData)
        {
            prev_rgi = rgi;
            rgi = rgi->rrgi_Next;
        }
        if (rgi == NULL)
        {
            /*
            * We've already exhausted the current growth
            * of the next pixel run to be the max. possible
            * to be 0.
            */
            rb_prev_run_state = 0;
        }
        else
        {
            /*
            * We are still within the current growth list
            * the run info appropriately.
            */
            if (rgi == r_growth_list)
                rb_prev_run_state = 0;
            else
                rb_prev_run_state = prev_rgi->rrgi_StateData;
        }
        /*
        * We can now start merging the elements of the row with the
        * elements of the growth list.
        */
        row_on = TRUE;
        while (R_NOT_END_OF_ROW(*row))
        {
            if (rgi == NULL || *row < rgi->rrgi_RgnData)
            {
                /*

```

which we actually have to
First, we check that we
free list that we

```
(!r_grow_free_list())
```

```
return FALSE;
```

```
r_free_list;
```

```
>rrgi_Next;
```

processing the first region. Therefore, we
state of the run to the lowest
bits.

```
#ifndef RB_USE_LOOKUP
```

```
>rrgi_StateData = (rb_prev_run_state & RB_CUR_STATE_MASK) |
```

```
(rb_prev_run_state >> RB_STATE_SIZE);
```

```
#else
```

```
>rrgi_StateData = r_shift_and_dup[rb_prev_run_state];
```

```
#endif
```

processing the second region. Therefore, the state data
copied to the previous state area so we
state.

```
>rrgi_StateData = rb_prev_run_state;
```

```
* This is the only situation in  
* create a new list element.  
* actually have an element in the  
* can use in the growth list..  
*/  
if (r_free_list == NULL)  
{  
    if
```

```
}  
prev_rgi->rrgi_Next =  
prev_rgi = r_free_list;  
r_free_list = r_free_list-  
prev_rgi->rrgi_Next = rgi;  
/*  
* Now, fill in the data..  
*/  
prev_rgi->rrgi_RgnData = *row;  
if (first)  
{
```

```
/*  
* We are  
* copy the current  
* RB_STATE_SIZE  
*/
```

```
prev_rgi-
```

```
prev_rgi-
```

```
/*  
* We are  
* has already been  
* just copy the  
*/  
prev_rgi-
```

```
}
```



```

region is active at this transition, we
current contents of the new list item.
behaviour of making that region active
turns it off if it is...

/*
 * Now, if the row for the current
 * xor the region mask with the
 * This gives the desired
 * if it is not there already, but
 */
if (row_on)
    prev_rgi-
>rrgi_StateData ^= (rgn_mask << RB_STATE_SIZE);
/*
 * We now move onto the next row
 */
row++;
row_on = !row_on;
continue;
}
/*
 * If the current row transition point is equal
in x position to the current
 * list item's transition point, we advance the
row counter to
 * the next position.
 */
if (*row == rgi->rrgi_RgnData)
{
    row++;
    row_on = !row_on;
}
/*
 * We update the current list item to deal with
the affects of the
 * current row run..
 */
rb_prev_run_state = rgi->rrgi_StateData;
if (first)
{
#ifdef RB_USE_LOOKUP
    rgi->rrgi_StateData =
    (rb_prev_run_state & RB_CUR_STATE_MASK) |
    (rb_prev_run_state >> RB_STATE_SIZE);
#else
    rgi->rrgi_StateData =
    r_shift_and_dup[rb_prev_run_state];
#endif
}
if (!row_on)
    rgi->rrgi_StateData ^= (rgn_mask
<< RB_STATE_SIZE);
/*
 * We now move onto the next element in the
list..
 */

```

```

        prev_rgi = rgi;
        rgi = rgi->rrgi_Next;
    }
    /*
    * Now, simply update the remainder of the elements in the
list..
    */
    if (first)
    {
        while (rgi != NULL)
        {
#ifdef RB_USE_LOOKUP
            rgi->rrgi_StateData = (rgi-
>rrgi_StateData & RB_CUR_STATE_MASK) |
(rgi->rrgi_StateData >> RB_STATE_SIZE);
#else
            rgi->rrgi_StateData =
r_shift_and_dup[rgi->rrgi_StateData];
#endif

            rgi = rgi->rrgi_Next;
        }
    }
    /*
    * Now copy "fake_item"'s next pointer to r_growth_list, as it
will have
    * changed if something was added to the head of the list..
    */
    r_growth_list = fake_item.rrgi_Next;
    /*
    * Update the return pointer to the region data..
    */
    *row_ptr = row;
    /*
    * Everything should now be OK..
    */
    return TRUE;
}

#ifdef R_USE_NEW_IMP
/*
* r_union_test_table
*
* A 16-int lookup table which when provided with an unsigned char
* of the following form xxyy, will provide evaluate the key
* state transition test of the union construction loop.
* Note that R_STATE_SIZE _must_ be 2 for this lookup table to
* work.
*/
int r_union_test_table[16] = {
0, 1, 1, 1,
0, 0, 0,
0, 0, 0,
1,
1,

```

```

0, 0, 0
1,
};

/*
 * R_union
 *
 * This function inits a R_Region structure to represent the union
 * of it's two arguments.
 *
 * Parameters:
 *     rgn    The R_Region to be initialised.
 *     r1     A R_Region ptr
 *     representing the first region.
 *     r2     A R_Region ptr
 *     representing the second region.
 * Returns
 *     TRUE on success, FALSE on failure.
 */
int
R_union
(
    R_Region    *rgn,
    R_Region    *r1,
    R_Region    *r2
)
{
    R_Int        *r1_dat;
    R_Int        *r2_dat;
    int          overlap_flags;

    union_tot++;

    if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
    {
        /*
         * The bounding boxes don't intersect. This means
         * union very easily, simply by copying data from
         * We malloc a new region data array of size r1-
         * r2->rr_RgnDataSize - 1. This is the maximum
         * resulting region. Not all of this memory will
         * the two regions being combined have rows with
         * (R_NEXT_IS_Y marker is not duplicated).
         */
        rgn->rr_RgnDataSize = r1->rr_RgnDataSize + r2-
        rgn->rr_RgnData = (R_Int *)malloc(rgn-
        rgn->rr_RgnDataSize *
        sizeof(R_Int));

        if (rgn->rr_RgnData == NULL)

```

```

    {
        return FALSE;
    }
    /*
    * Now, check to see if the regions overlap in
y...
    */
    if (!(overlap_flags & BB_INTERSECT_OVERLAP_Y))
    {
        /*
        * The regions don't overlap in y.
        * and then another into the array
        * that r1 points to the region
        */
        if (r2->rr_BBox.Y.Min < r1->rr_BBox.Y.Min)
        {
            R_Region
            tmp = r1;
            r1 = r2;
            r2 = tmp;
        }
        memcpy
        (
            rgn->rr_RgnData,
            r1->rr_RgnData,
            (r1->rr_RgnDataSize - 1) * sizeof(R_Int)
        );
        memcpy
        (
            rgn->rr_RgnData +
            r1->rr_RgnDataSize - 1,
            r2->rr_RgnData,
            r2->rr_RgnDataSize * sizeof(R_Int)
        );
        ASSERT(rgn->rr_RgnData[rgn->rr_RgnDataSize - 1] == R_EOR);
    }
    else
    {
        R_Int      *r1_tmp;
        R_Int      *r2_tmp;
        R_Int      *dest;
        R_Int      min_row;
        int
        r1_done;
        int
        r1_consumed;
        int
        r2_consumed;
        int
        num_written;
    }

```

in x. We simply go row
memcpy the individual rows as
points to the region with

>rr_BBox.X.Min)

*tmp;

!= R_EOR)

R_NEXT_IS_Y);

R_NEXT_IS_Y);

min(r1_dat[1], r2_dat[1]);

min_row)

need to emit r1. We therefore need to find where
the next row (if any) starts. When we do this we
recall that a y value _must be followed by at least
two x values..

r1_tmp = r1_dat + 4;

while (*r1_tmp != R_NEXT_IS_Y && *r1_tmp != R_EOR)

r1_tmp++;

num_written = r1_tmp - r1_dat;

memcpy(dest, r1_dat, num_written * sizeof(R_Int));

```
/*
 * The bboxes overlap in y but not
 * by row through each region and
 * appropriate. We ensure that r1
 * the smallest x coordinate.
 */
if (r2->rr_BBox.X.Min < r1-
```

```
{
    R_Region
    tmp = r1;
    r1 = r2;
    r2 = tmp;
}
r1_dat = r1->rr_RgnData;
r2_dat = r2->rr_RgnData;
dest = rgn->rr_RgnData;
rgn->rr_RgnDataSize = 0;
r1_consumed = 0;
r2_consumed = 0;
while (*r1_dat != R_EOR && *r2_dat
```

```
{
    ASSERT(*r1_dat ==
    ASSERT(*r2_dat ==
    min_row =
    r1_done = FALSE;
    if (r1_dat[1] ==
    {
```

```
/*
 * We
 *
 *
 *
 */
```

```
dest
+=  num_written;

r1_consumed +=  num_written;

>rr_RgnDataSize +=  num_written;

r1_dat = r1_tmp;

r1_done = TRUE;

min_row)

need to emit r1. We therefore need to find where
the next row (if any) starts. When we do this we
recall that a y value _must be followed by at least
two x values. If r1's current row has already been
emitted for this y value, we do _not_ emit the
R_NEXT_IS_Y marker or the y value itself.

(r1_done)

r2_dat += 2;
r2_tmp = r2_dat + 2;
r2_consumed += 2;

r2_tmp = r2_dat + 4;

while (*r2_tmp != R_NEXT_IS_Y && *r2_tmp != R_EOR)
r2_tmp++;
num_written = r2_tmp - r2_dat;
memcpy(dest, r2_dat, num_written * sizeof(R_Int));
+= num_written;
r2_consumed += num_written;
>rr_RgnDataSize += num_written;
```

```
}
if (r2_dat[1] ==
{
/*
 * We
 *
 *
 *
 *
 *
 *
 */
if
{
}
else
{
}
}
```

r2_dat = r2_tmp;

region left standing. We memcpy
the region (including the
the destination.

== r2->rr_RgnDataSize - 1);

dest,

r1_dat,

>rr_RgnDataSize - r1_consumed) * sizeof(R_Int)

>rr_RgnDataSize += (r1->rr_RgnDataSize - r1_consumed);

region left standing. We memcpy
the region (including the
the destination.

== r1->rr_RgnDataSize - 1);

dest,

r2_dat,

>rr_RgnDataSize - r2_consumed) * sizeof(R_Int)

>rr_RgnDataSize += (r2->rr_RgnDataSize - r2_consumed);

>rr_RgnData[rgn->rr_RgnDataSize - 1] == R_EOR

```
    }  
    if (*r1_dat != R_EOR)  
    {  
        /*  
        * r1 is the last  
        * the remainder of  
        * R_EOR marker) to  
        */  
        ASSERT(r2_consumed  
memcpy  
(  
  
        (r1-  
);  
rgn-  
    }  
    else  
    {  
        /*  
        * r2 is the last  
        * the remainder of  
        * R_EOR marker) to  
        */  
        ASSERT(r1_consumed  
memcpy  
(  
  
        (r2-  
);  
rgn-  
    }  
    ASSERT  
(  
    rgn-  
);
```



```

    }
    else
    {
        min_row;
        dest_size;

        in_run;
        done_r1_in_row;

        R_Int
        int
        R_RgnGrowItem *rgi;
        R_RgnGrowItem *rgi_tail;
        int
        int

        union_full++;

        /*
         * The two regions _do_ overlap in x_and_ y. We
         * to do a bit more work in calculating the union
         * regions. We use the a list of R_RgnGrowItem
         * regarding the currently active regions as we
         * the rows of each region. After any rows
         * are added to the list, we examine the state
         * pixel run in the list. If the addition of the
         * for the y-coord have caused a transition to
         * the pixel run is emitted.
         */
        r1_dat = r1->rr_RgnData;
        r2_dat = r2->rr_RgnData;
        dest_size = 0;
        /*
         * We are now ready to loop through the data of
         * We continue building the new region whilst
         * remaining in either of the two regions.
         */
        while (*r1_dat != R_EOR || *r2_dat != R_EOR)
        {
            ASSERT(*r1_dat == R_NEXT_IS_Y ||
                *r2_dat == R_NEXT_IS_Y ||
                if (*r1_dat == R_EOR)
                    min_row =
                else if (*r2_dat == R_EOR)
                    min_row =

            both regions.

            there is data

            *r1_dat == R_EOR);
            *r2_dat == R_EOR);

            r2_dat[1];

            r1_dat[1];
    }

```



```

min(r1_dat[1], r2_dat[1]);

== min_row)

is active on this y coord. We add this
current region builder.

(!R_add_row_to_region_growth_list(&r1_dat, 0x1, TRUE))
return FALSE;
TRUE;

== min_row)

is active on this y coord. We add this
current region builder.

(!R_add_row_to_region_growth_list(&r2_dat, 0x2, !done_r1_in_row))
return FALSE;

for the input rows.

(!r_check_rgn_buf_len(dest_size + 2))

R_NEXT_IS_Y;

#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
NULL; rgi = rgi->rrgi_Next)
#endif 0

>rrgi_StateData > 0

else
    min_row =
done_r1_in_row = FALSE;
if (*r1_dat != R_EOR && r1_dat[1]
{
    /*
     * The first region
     * row to the
     */
    if
done_r1_in_row =
}
if (*r2_dat != R_EOR && r2_dat[1]
{
    /*
     * The first region
     * row to the
     */
    if
}
/*
 * Now, we generate the output row
 */
if
{
    return FALSE;
}
r_RgnBuf[dest_size++] =
r_RgnBuf[dest_size++] = min_row;
in_run = FALSE;
for (rgi = r_growth_list; rgi !=
{
    if
    (
        rgi-
        &&
        (

```

```

(rgi->rrgi_StateData & RB_CUR_STATE_MASK) == 0
    ||

(rgi->rrgi_StateData & RB_PREV_STATE_MASK) == 0
    )

#else
    if

(r_union_test_table[rgi->rrgi_StateData])
    {

#endif
        /*
         * We
         * in

        */
        if

        {
            if

            {

            }

        }
    }
else
{
    if

    {
        /*
         *

        */
        if

        {

        }

    }

}

/*
 * Not efficient,
 */

get rid of it..

```

```

== NULL)
rgi_tail = rgi;
}
#else
rgi = r_growth_list;
rgi_tail = rgi;
while (rgi != NULL && !r_union_test_table[rgi-
>rrgi_StateData])
    rgi = rgi->rrgi_Next;
while (rgi != NULL)
{
    if
    return FALSE;
    r_RgnBuf[dest_size++] = rgi-
do
{
    rgi = rgi-
} while (rgi != NULL &&
rgi_tail = rgi;
r_RgnBuf[dest_size++] = rgi-
do
{
    rgi = rgi-
} while (rgi != NULL &&
!r_union_test_table[rgi->rrgi_StateData]);
}
#endif
R_NEXT_IS_Y)
if (r_RgnBuf[dest_size - 2] ==
{
    /*
    * We didn't output
    */
    dest_size -= 2;
}
}
/*
* Now, we've completed using the growth list for
constructing this
* region. Therefore, we add it to the front of
the free list, to
* be re-used later.
*/
#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
while (rgi_tail->rrgi_Next != NULL)
    rgi_tail = rgi_tail->rrgi_Next;
#endif
rgi_tail->rrgi_Next = r_free_list;
r_free_list = r_growth_list;

```

```

        r_growth_list = NULL;
        /*
        * We've completed constructing the data for the
region. We
        * make a copy the constructed data from the
permanent buffer to
        * an exactly fitting buffer.
        */
        rgn->rr_RgnData = (R_Int *)malloc(++dest_size *
sizeof(R_Int));
        if (rgn->rr_RgnData == NULL)
        {
            return FALSE;
        }
        memcpy(rgn->rr_RgnData, r_RgnBuf, (dest_size -
1) * sizeof(R_Int));
        rgn->rr_RgnData[dest_size - 1] = R_EOR;
        rgn->rr_RgnDataSize = dest_size;
        ASSERT(rgn->rr_RgnDataSize >= 9);
    }
    /*
    * We now do a bounding box union of the two component bboxes
and place
    * the result in the new region.
    */
    BB_union(&r1->rr_BBox, &r2->rr_BBox, &rgn->rr_BBox);
    /*
    * Done! We can get out..
    */
    return TRUE;
}

/*
* R_union_equals
*
* This function basically implements a r1 union= r2 type operation. Ie
* r1 union r2 is calculated and the result returned in r1.
*
* Parameters:
*
*           r1                A pointer to an
R_Region. This represents
*                           the first half of
the union, and is also used to return
*                           the eventual
result.
*           r2                A pointer to an
R_Region. This represents the second
*                           half of the union.
*
* Returns:
*
*           TRUE on success, FALSE on failure.
*/
int
R_union_equals
(
    R_Region    *r1,
    R_Region    *r2
)
{

```

```

        R_Region    new_rgn;
        if (r1->rr_RgnData == NULL)
            return R_init_region_with_region(r1, r2);
        if (!R_union(&new_rgn, r1, r2))
            return FALSE;
        R_empty_region(r1);
        *r1 = new_rgn;
        return TRUE;
    }

/*
 * r_intersection_test_table
 *
 * A 16-int lookup table which when provided with an unsigned char
 * of the following form xxyy, will provide evaluate the key
 * state transition test of the intersection construction loop.
 * Note that R_STATE_SIZE _must_ be 2 for this lookup table to
 * work.
 */
int r_intersection_test_table[16] = {
    0,
    0, 0, 1,
    0,
    0, 0, 1,
    0,
    0, 0, 1,
    1,
    1, 1, 0
};

/*
 * R_intersection
 *
 * This function inits a R_Region structure to represent the intersection
 * of it's two arguments.
 *
 * Parameters:
 *     rgn      A R_Region ptr to the R_Region structure to be initialised.
 *     r1      A R_Region ptr
 *     representing the first region.
 *     r2      A R_Region ptr
 *     representing the second region.
 * Returns
 *     TRUE on success, FALSE on failure.
 */
int
R_intersection
(
    R_Region    *rgn,
    R_Region    *r1,
    R_Region    *r2
)
{
    R_Int        *r1_dat;
    R_Int        *r2_dat;
    int          overlap_flags;

```

```

        int_tot++;

        rgn->rr_RgnData = NULL;
        if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
        {
            /*
            * The bounding boxes don't intersect. This means
            * don't intersect. Therefore, we simply set rgn-
            * (signifying an empty region) and get out..
            */
            return TRUE;
        }
        R_Int
min_row;
        int
dest_size;
        R_RgnGrowItem      *rgi;
        R_RgnGrowItem      *rgi_tail;
        int
in_run;
        int
done_r1_in_row;
        IntXYMinMax
                                new_bbox;

        int_full++;

        /*
        * The two regions _do_ overlap in x _and_ y. We therefore have
        * to do a bit more work in calculating the intersection of the
two
        * regions. We use the R_RegionBuilder struct to store state
        * regarding the currently active regions as we progress
through
        * the rows of each region. After any rows relevant to a y-coord
        * are added to the region builder, we examine the state of each
        * pixel run in the region builder. If the addition of the
row(s)
        * for the y-coord have caused a transition to or from 0x3, then
        * the pixel run is emitted.
        */
        /*
        * Initialise the new_bbox structure for determining the new
bounding box.
        */
        new_bbox.X.Min = R_INT_MAX_VALUE;
        new_bbox.Y.Min = R_INT_MAX_VALUE;
        new_bbox.X.Max = R_INT_MIN_VALUE;
        new_bbox.Y.Max = R_INT_MIN_VALUE;

        /*
        * The next thing we do is ensure the current region builder
is empty,
        * and set up pointers into the region data of the two regions.
        */
        r1_dat = r1->rr_RgnData;

```

```

r2_dat = r2->rr_RgnData;
dest_size = 0;
/*
 * We are now ready to loop through the data from both regions.
Notice
 * that we only keep looping whilst _both_ regions have some
data left
 * to give. As soon as either of the region's data has been
exhausted,
 * then we stop as the intersection region has already been
calculated
 * and is sitting in the rgn_buf.
 */
while (*r1_dat != R_EOR && *r2_dat != R_EOR)
{
    ASSERT(*r1_dat == R_NEXT_IS_Y || *r1_dat ==
R_EOR);
    ASSERT(*r2_dat == R_NEXT_IS_Y || *r2_dat ==
R_EOR);

    if (*r1_dat == R_EOR)
        min_row = r2_dat[1];
    else if (*r2_dat == R_EOR)
        min_row = r1_dat[1];
    else
        min_row = min(r1_dat[1],
r2_dat[1]);

    done_r1_in_row = FALSE;
    if (*r1_dat != R_EOR && r1_dat[1] == min_row)
    {
        /*
        * The first region is active on
this y coord. We add this
builder.
        * row to the current region
        */
        if
(!R_add_row_to_region_growth_list(&r1_dat, 0x1, TRUE))
            return FALSE;
        done_r1_in_row = TRUE;
    }
    if (*r2_dat != R_EOR && r2_dat[1] == min_row)
    {
        /*
        * The first region is active on
this y coord. We add this
builder.
        * row to the current region
        */
        if
(!R_add_row_to_region_growth_list(&r2_dat, 0x2, !done_r1_in_row))
            return FALSE;
    }
    /*
    * Now, we generate the output row for the input
rows.
    */
    if (!r_check_rgn_buf_len(dest_size + 2))
    {

```

```

        return FALSE;
    }
    r_RgnBuf[dest_size++] = R_NEXT_IS_Y;
    r_RgnBuf[dest_size++] = min_row;
    in_run = FALSE;
#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
    for (rgi = r_growth_list; rgi != NULL; rgi = rgi-
>rrgi_Next)
    {
        #if 0
            if
            (
                rgi-
>rrgi_StateData != (3 | (3 << RB_STATE_SIZE))
                &&
                (
                    (rgi->rrgi_StateData & RB_PREV_STATE_MASK) == 3
                    ||
                    (rgi->rrgi_StateData & RB_CUR_STATE_MASK) == (3 << RB_STATE_SIZE)
                )
            )
        #else
            if
            (r_intersection_test_table[rgi->rrgi_StateData])
        #endif
        {
            /*
             * We have to emit
             * in one..
             */
            if (!in_run)
            {
                if
                {
                    return FALSE;
                }

                r_RgnBuf[dest_size++] = rgi->rrgi_RgnData;
                in_run = TRUE;
                new_bbox.X.Min = min(new_bbox.X.Min, rgi->rrgi_RgnData);
            }
            else
            {
                if (in_run)
                {
                    /*
                     *
                     * We've come to the end of a run. We output the next element to end it.
                     */

```



```

                                                                    if
(!r_check_rgn_buf_len(dest_size + 1))
                                                                    {

return FALSE;
                                                                    }

r_RgnBuf[dest_size++] = rgi->rrgi_RgnData;

new_bbox.X.Max = max(new_bbox.X.Max, rgi->rrgi_RgnData);
                                                                    }
                                                                    in_run = FALSE;
                                                                    }
                                                                    /*
                                                                    * Not efficient, get rid of it..
                                                                    */
                                                                    if (rgi->rrgi_Next == NULL)
                                                                    rgi_tail = rgi;
                                                                    }
#else
                                                                    rgi = r_growth_list;
                                                                    rgi_tail = rgi;
                                                                    while (rgi != NULL &&
!r_intersection_test_table[rgi->rrgi_StateData])
                                                                    rgi = rgi->rrgi_Next;
                                                                    while (rgi != NULL)
                                                                    {
                                                                    if
(!r_check_rgn_buf_len(dest_size + 2))
                                                                    return FALSE;
                                                                    r_RgnBuf[dest_size++] = rgi-
>rrgi_RgnData;
                                                                    new_bbox.X.Min =
do
                                                                    {
                                                                    rgi = rgi-
>rrgi_Next;
                                                                    } while (rgi != NULL &&
r_intersection_test_table[rgi->rrgi_StateData]);
                                                                    rgi_tail = rgi;
                                                                    r_RgnBuf[dest_size++] = rgi-
>rrgi_RgnData;
                                                                    new_bbox.X.Max =
do
                                                                    {
                                                                    rgi = rgi-
>rrgi_Next;
                                                                    } while (rgi != NULL &&
!r_intersection_test_table[rgi->rrgi_StateData]);
                                                                    }
#endif
                                                                    if (r_RgnBuf[dest_size - 2] == R_NEXT_IS_Y)
                                                                    {
                                                                    /*
                                                                    * We didn't output anything for
these input rows. Rewind..

```

```

        */
        dest_size -= 2;
    }
    else
    {
        if (min_row < new_bbox.Y.Min)
            new_bbox.Y.Min =
min_row;

        else if (min_row >
new_bbox.Y.Max)
            new_bbox.Y.Max =
min_row;
    }
}
/*
 * Now, we've completed using the growth list for constructing
this
 * region. Therefore, we add it to the front of the free list,
to
 * be re-used later.
 */
#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
    while (rgi_tail->rrgi_Next != NULL)
        rgi_tail = rgi_tail->rrgi_Next;
#endif

rgi_tail->rrgi_Next = r_free_list;
r_free_list = r_growth_list;
r_growth_list = NULL;
/*
 * We've completed constructing the data for the region.
Firstly
 * we check to see if we've emitted anything at all. If we have
 * then dest_size must be > 0. If it isn't we simply free the
 * region we created and get out, as the regions don't really
 * intersect, in spite of their intersecting bounding boxes.
 */
if (dest_size == 0)
{
    return TRUE;
}
/*
 * We make a copy the constructed data from the permanent buffer
to
 * an exactly fitting buffer.
 */
rgn->rr_RgnData = (R_Int *)malloc(++dest_size *
sizeof(R_Int));
if (rgn->rr_RgnData == NULL)
{
    return FALSE;
}
memcpy(rgn->rr_RgnData, r_RgnBuf, (dest_size - 1) *
sizeof(R_Int));
rgn->rr_RgnData[dest_size - 1] = R_EOR;
rgn->rr_RgnDataSize = dest_size;
ASSERT(rgn->rr_RgnDataSize >= 9);
/*

```

```

        * Now, copy across the bounding box.. Before we do this, we
subtract
        * 1 from X.Max and Y.Max because of the region format.
        */
        new_bbox.X.Max--;
        new_bbox.Y.Max--;
        rgn->rr_BBox = new_bbox;
        /*
        * Done! We can get out..
        */
        return TRUE;
}

/*
* r_difference_test_table
*
* A 16-int lookup table which when provided with an unsigned char
* of the following form xxyy, will provide evaluate the key
* state transition test of the difference construction loop.
* Note that R_STATE_SIZE _must_ be 2 for this lookup table to
* work.
*/
int r_difference_test_table[16] = {

0, 1, 0, 0,
0, 1, 1,
1, 0, 0,
1, 0, 0

1,
0,
0,
0,
};

/*
* R_difference
*
* This function inits a R_Region structure to represent the difference of
* it's two arguments. It essentially calculates r1 - r2
*
* Parameters:
*     rgn      A R_Region ptr representing the R_Region to be initied.
*     r1      A R_Region ptr
representing the first region.
*     r2      A R_Region ptr
representing the second region.
* Returns
*     TRUE on success, FALSE on failure.
*/
int
R_difference
(
    R_Region    *rgn,
    R_Region    *r1,
    R_Region    *r2
)
{
    R_Int        *r1_dat;
    R_Int        *r2_dat;

```

```

int                                     overlap_flags;

diff_tot++;

rgn->rr_RgnData = NULL;
if (!BB_intersect_test(&r1->rr_BBox, &r2->rr_BBox,
&overlap_flags))
{
    /*
    * The bounding boxes don't intersect. This means
    that r1 - r2
    * simply equals r1. We make a copy of the
    relevant bits and get out..
    */
    rgn->rr_BBox = r1->rr_BBox;
    rgn->rr_RgnDataSize = r1->rr_RgnDataSize;
    rgn->rr_RgnData = (R_Int *)malloc(r1-
>rr_RgnDataSize * sizeof(R_Int));
    if (rgn->rr_RgnData == NULL)
    {
        return FALSE;
    }
    memcpy
    (
        rgn->rr_RgnData,
        r1->rr_RgnData,
        r1->rr_RgnDataSize *
sizeof(R_Int)
    );
    return TRUE;
}
R_Int                                     min_row;
int
dest_size;
R_RgnGrowItem *rgi;
R_RgnGrowItem *rgi_tail;
int
in_run;
int
done_r1_in_row;
unsigned char m_high;
unsigned char m_low;
IntXYMinMax                                     new_bbox;

diff_full++;

/*
* The two regions _do_ overlap in x _and_ y. We therefore have
* to do a bit more work in calculating the difference of the
two
* regions. We use the R_RegionBuilder struct to store state
* regarding the currently active regions as we progress
through
* the rows of each region. After any rows relevant to a y-coord
* are added to the region builder, we examine the state of each
* pixel run in the region builder. If the addition of the
row(s)

```

```

* for the y-coord have caused the following transitions -
*
*           r1      -> 0
*           0       -> r2
*           r1 + r2 -> r2
*           r2      -> r1 + r2
* ..then the relevent runs are emitted. Firstly, though,
* we ensure the current region builder is empty,
* and set up pointers into the region data of the two regions.
*/
r1_dat = r1->rr_RgnData;
r2_dat = r2->rr_RgnData;
dest_size = 0;

/*
* Initialise the new_bbox structure for determining the new
bounding box.
*/
new_bbox.X.Min = 32767;
new_bbox.Y.Min = 32767;
new_bbox.X.Max = -32768;
new_bbox.Y.Max = -32768;

/*
* We are now ready to loop through the data from both regions.
Notice
* that we only keep looping whilst r1 has data outstanding.
When
* r1's data is consumed, then any transitions made by r2 are
* irrelevant.
*/
while (*r1_dat != R_EOR)
{
    ASSERT(*r1_dat == R_NEXT_IS_Y || *r1_dat ==
R_EOR);
    ASSERT(*r2_dat == R_NEXT_IS_Y || *r2_dat ==
R_EOR);

    if (*r1_dat == R_EOR)
        min_row = r2_dat[1];
    else if (*r2_dat == R_EOR)
        min_row = r1_dat[1];
    else
        min_row = min(r1_dat[1],
r2_dat[1]);

    done_r1_in_row = FALSE;
    if (*r1_dat != R_EOR && r1_dat[1] == min_row)
    {
        /*
        * The first region is active on
this y coord. We add this
builder.
        */
        if
(!R_add_row_to_region_growth_list(&r1_dat, 0x1, TRUE))
            return FALSE;
        done_r1_in_row = TRUE;
    }
    if (*r2_dat != R_EOR && r2_dat[1] == min_row)

```

```

        {
            /*
            * The first region is active on
            this y coord. We add this
            builder.
            * row to the current region
            */
            if
            (!R_add_row_to_region_growth_list(&r2_dat, 0x2, !done_r1_in_row))
                return FALSE;
        }
        /*
        * Now, we generate the output row for the input
        rows.
        */
        if (!r_check_rgn_buf_len(dest_size + 2))
        {
            return FALSE;
        }
        r_RgnBuf[dest_size++] = R_NEXT_IS_Y;
        r_RgnBuf[dest_size++] = min_row;
        in_run = FALSE;
#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
        for (rgi = r_growth_list; rgi != NULL; rgi = rgi-
>rrgi_Next)
        {
            #if 0
            RB_CUR_STATE_MASK) >> RB_STATE_SIZE;
            RB_PREV_STATE_MASK;

            m_high = (rgi->rrgi_StateData &
            m_low = rgi->rrgi_StateData &

            if
            (
                (
                    (m_low != 1 && m_high == 1)
                    ||
                    (m_low == 1 && m_high != 1)
                )
            )
            #else
            if (r_difference_test_table[rgi-
>rrgi_StateData])
            #endif
            {
                /*
                * We have to emit
                * in one..
                */
                if (!in_run)
                {
                    if
                    {
                        (!r_check_rgn_buf_len(dest_size + 1))

                        return FALSE;
                    }
                }
            }
        }
    }

```

```

}

r_RgnBuf[dest_size++] = rgi->rrgi_RgnData;

in_run = TRUE;

new_bbox.X.Min = min(new_bbox.X.Min, rgi->rrgi_RgnData);
}
else
{
    if (in_run)
    {
        /*
        *
        * We've come to the end of a run. We output the next element to end it.
        */
        if
        {
            (!r_check_rgn_buf_len(dest_size + 1))

            return FALSE;
        }

        r_RgnBuf[dest_size++] = rgi->rrgi_RgnData;

        new_bbox.X.Max = max(new_bbox.X.Max, rgi->rrgi_RgnData);
    }
    in_run = FALSE;
}
/*
 * Not efficient, get rid of it..
 */
if (rgi->rrgi_Next == NULL)
    rgi_tail = rgi;
}

#else
    rgi = r_growth_list;
    rgi_tail = rgi;
    while (rgi != NULL &&
    !r_difference_test_table[rgi->rrgi_StateData])
        rgi = rgi->rrgi_Next;
    while (rgi != NULL)
    {
        if
        {
            (!r_check_rgn_buf_len(dest_size + 2))

            return FALSE;
            r_RgnBuf[dest_size++] = rgi-
            >rrgi_RgnData;

            new_bbox.X.Min =
            min(new_bbox.X.Min, rgi->rrgi_RgnData);

            do
            {
                rgi = rgi-
                >rrgi_Next;

            } while (rgi != NULL &&
            r_difference_test_table[rgi->rrgi_StateData]);
            rgi_tail = rgi;
        }
    }
}

```

```

>rrgi_RgnData;
max(new_bbox.X.Max, rgi->rrgi_RgnData);

>rrgi_Next;
!r_difference_test_table[rgi->rrgi_StateData]);
#endif
    if (r_RgnBuf[dest_size - 2] == R_NEXT_IS_Y)
    {
        /*
         * We didn't output anything for
these input rows. Rewind..
        */
        dest_size -= 2;
    }
    else
    {
        if (min_row < new_bbox.Y.Min)
            new_bbox.Y.Min =
min_row;
        else if (min_row >
new_bbox.Y.Max)
            new_bbox.Y.Max =
min_row;
    }
}
/*
 * Now, we've completed using the growth list for constructing
this
 * region. Therefore, we add it to the front of the free list,
to
 * be re-used later.
 */
#ifdef R_NEW_IMP_CONSTRUCTION_LOOP
    while (rgi_tail->rrgi_Next != NULL)
        rgi_tail = rgi_tail->rrgi_Next;
#endif
    rgi_tail->rrgi_Next = r_free_list;
    r_free_list = r_growth_list;
    r_growth_list = NULL;
/*
 * We've completed constructing the data for the region.
Firstly
 * we check to see if we've emitted anything at all. If we have
 * then dest_size must be > 0. If it isn't we simply free the
 * region we created and get out, as r2 - r1 must be empty.
 */
if (dest_size == 0)
{
    return TRUE;
}
/*

```



```

        * We make a copy the constructed data from the permanent buffer
to
        * an exactly fitting buffer.
        */
    rgn->rr_RgnData = (R_Int *)malloc(++dest_size *
sizeof(R_Int));
    if (rgn->rr_RgnData == NULL)
    {
        return FALSE;
    }
    memcpy(rgn->rr_RgnData, r_RgnBuf, (dest_size - 1) *
sizeof(R_Int));
    rgn->rr_RgnData[dest_size - 1] = R_EOR;
    rgn->rr_RgnDataSize = dest_size;
    ASSERT(rgn->rr_RgnDataSize >= 9);
    /*
    * Now, copy across the bounding box..
    */
    rgn->rr_BBox = new_bbox;
    /*
    * Done! We can get out..
    */
    return TRUE;
}
#endif /* R_USE_NEW_IMP */

/*
 * R_compare
 *
 * This function compares two regions and determines if they are the same.
 *
 * Parameters:
 *
 *          rgn1          The first R_Region.
 *          rgn2          The second R_Region.
 *
 * Returns:
 *
 *          TRUE if they are the same, FALSE if they aren't.
 */
int
R_compare
(
    R_Region    *rgn1,
    R_Region    *rgn2
)
{
    /*
    * If their region data sizes don't agree, then they aren't the
same.
    */
    if (rgn1->rr_RgnDataSize != rgn2->rr_RgnDataSize)
        return FALSE;
    if
    (
        memcmp
        (
            rgn1->rr_RgnData,
            rgn2->rr_RgnData,
            rgn1->rr_RgnDataSize *
sizeof(R_Int)

```

```

        )
        ==
        0
    )
    return TRUE;
return FALSE;
}

/*
 * r_check_rect_buf_len
 *
 * This function checks to see if the static rectangle buffer is large
 * enough. If it isn't then it is reallocated to make it large enough.
 *
 * Parameters:
 *          size          The required size of the r_RectBuf
array.
 * Returns:
 *          TRUE on success, FALSE on failure.
 */
static int
r_check_rect_buf_len
(
    int          size
)
{
    ASSERT(size >= 0);
    if (size > r_RectBufSize)
    {
        int
new_buf_size;

        IntXYMinMax    *new_buf;
        new_buf_size = max(size, r_RectBufSize * 2);
        new_buf = (IntXYMinMax *)malloc
                                (new_buf_size
 * sizeof(IntXYMinMax)
                                );
        if (new_buf == NULL)
            return FALSE;
        if (r_RectBuf != NULL)
        {
            memcpy(new_buf, r_RectBuf,
r_RectBufSize * sizeof(IntXYMinMax));
            free(r_RectBuf);
        }
        r_RectBuf = new_buf;
        r_RectBufSize = new_buf_size;
    }
    return TRUE;
}

#ifdef R_USE_NEW_IMP
/*
 * R_rects_from_region
 *
 * This function returns a group of non-overlapping rectangles which
 * together constitute the region. The group of rectangles returned is
 * currently non-optimal as the function uses the R_RegionBuilder structure

```

```

* to store state. A more specific data structure will be required to
* make the rectangles produced more optimal.
*
* Parameters:
*           rgn                The region from
which a rectangle array is required.
*           rects            A pointer to a pointer to a
IntXYMinMax structure. Used
*                           to return the
array.
*           num_rects        A pointer to an int. Used to
return the number of
*                           elements in the
array.
*           static_ok        This boolean arg is passed as TRUE
if a pointer to
*                           the r_RectBuf is
sufficient. This is TRUE if usefulness of
*                           the rectangle data
obtained ends before the next call to
*
R_rects_from_region (for any region). FALSE is passed if
*                           a newly malloced
copy is required. Basically is TRUE is
*                           passed the pointer
returned must _not_ be freed.
* Returns:
*                           TRUE on success, FALSE on failure.
*/
int
R_rects_from_region
(
    R_Region    *rgn,
    IntXYMinMax **rects,
    int         *num_rects,
    int         static_ok
)
{
    R_Int      *rgn_data;
    int
dest_index;
    int
prev_y;
    int
prev_x;
    unsigned char *rgn_bld_stat;
    R_Int      *rgn_bld_dat;
    int
    int
in_run;

    /*
     * Give "nice" defaults for return stuff in cause we fail..
     */
    *rects = NULL;
    *num_rects = 0;
    /*

```

```

ensure
    * We grab a pointer to the region data for the region and
    * that the current region builder is empty..
    */
    rgn_data = rgn->rr_RgnData;
    if (rgn_data == NULL)
        /*
        * This is an empty region.. Get out..
        */
        return TRUE;
    ASSERT(*rgn_data == R_NEXT_IS_Y);
    R_CurRB->rrb_Nels = 0;
    /*
    * We add the first row of the region to the region builder.
    * store the y-coord of this first row.
    */
    prev_y = rgn_data[1];
    if (!R_add_row_to_region_builder(&rgn_data, 0x1, TRUE))
        return FALSE;
    ASSERT(*rgn_data == R_NEXT_IS_Y);
    ASSERT(*rgn_data != R_EOR);
    /*
    * We are now in a position to loop through the data of the
    * We continue until the region data runs out. Basically, we
    * the runs in the current region builder out as rectangles.
    * x-coords from the region builder and y coords of the rows.
    * we add then next row to the region builder.
    */
    dest_index = 0;
    while (*rgn_data != R_EOR)
    {
        ASSERT(*rgn_data == R_NEXT_IS_Y);

        rgn_bld_stat = R_CurRB->rrb_StateData;
        rgn_bld_dat = R_CurRB->rrb_RgnData;
        in_run = FALSE;
        for (i = R_CurRB->rrb_Nels; i > 0; i--)
        {
            if ((*rgn_bld_stat &
RB_CUR_STATE_MASK) > 0)
            {
                /*
                * We have to emit
                * in one..
                */
                if (!in_run)
                {
                    prev_x = *rgn_bld_dat;
                    in_run = TRUE;
                }
            }
        }
    }

```

```

    }
    else
    {
        if (in_run)
        {
            /*
            *
            *
            */
            if
            (!r_check_rect_buf_len(dest_index + 1))
            return FALSE;

            r_RectBuf[dest_index].X.Min = prev_x;
            r_RectBuf[dest_index].Y.Min = prev_y;
            r_RectBuf[dest_index].X.Max = *rgn_bld_dat - 1;
            r_RectBuf[dest_index++].Y.Max = rgn_data[1] - 1;

            }
            in_run = FALSE;
        }
        rgn_bld_stat++;
        rgn_bld_dat++;
    }
    /*
    * Now, we advance onto the next row..
    */
    prev_y = rgn_data[1];
    if (!R_add_row_to_region_builder(&rgn_data,
0x1, TRUE))
        return FALSE;
}
/*
* Ok, we have the array of rectangles sitting around. If
static_ok
* is TRUE then we simply set the return pointers and get out.
* Otherwise, we need to malloc a copy of the r_RectBuf.
*/
*num_rects = dest_index;
if (static_ok)
{
    *rects = r_RectBuf;
}
else
{
    *rects = (IntXYMinMax *)malloc(dest_index *
sizeof(IntXYMinMax));
    if (*rects == NULL)
        return FALSE;
    memcpy(*rects, r_RectBuf, dest_index *
sizeof(IntXYMinMax));
}
return TRUE;

```

```

}
#else
/*
 * R_rects_from_region
 *
 * This function returns a group of non-overlapping rectangles which
 * together constitute the region. The group of rectangles returned is
 * currently non-optimal as the function uses the R_RegionBuilder structure
 * to store state. A more specific data structure will be required to
 * make the rectangles produced more optimal.
 *
 * Parameters:
 *
 *          rgn                      The region from
which a rectangle array is required.
 *          rects                  A pointer to a pointer to a
IntXYMinMax structure. Used
 *                               to return the
array.
 *          num_rects              A pointer to an int. Used to
return the number of
 *                               elements in the
array.
 *          static_ok              This boolean arg is passed as TRUE
if a pointer to
 *                               the r_RectBuf is
sufficient. This is TRUE if usefulness of
 *                               the rectangle data
obtained ends before the next call to
 *
R_rects_from_region (for any region). FALSE is passed if
 *                               a newly malloced
copy is required. Basically is TRUE is
 *                               passed the pointer
returned must _not_ be freed.
 * Returns:
 *
 *          TRUE on success, FALSE on failure.
 */
int
R_rects_from_region
(
    R_Region      *rgn,
    IntXYMinMax  **rects,
    int           *num_rects,
    int           static_ok
)
{
    R_Int          *rgn_data;
    int
dest_index;
    R_RgnGrowItem *rgi;
    R_RgnGrowItem *rgi_tail;
    int
prev_y;
    int
prev_x;
    int
in_run;

```

```

/*
 * Give "nice" defaults for return stuff in cause we fail..
 */
*rects = NULL;
*num_rects = 0;
/*
 * We grab a pointer to the region data for the region and
ensure
 * that the current region builder is empty..
 */
rgn_data = rgn->rr_RgnData;
if (rgn_data == NULL)
    /*
     * This is an empty region.. Get out..
     */
    return TRUE;
ASSERT(*rgn_data == R_NEXT_IS_Y);
/*
 * We add the first row of the region to the region builder.
We also
 * store the y-coord of this first row.
 */
prev_y = rgn_data[1];
if (!R_add_row_to_region_growth_list(&rgn_data, 0x1, TRUE))
    return FALSE;
ASSERT(*rgn_data == R_NEXT_IS_Y);
ASSERT(*rgn_data != R_EOR);
/*
 * We are now in a position to loop through the data of the
region.
 * We continue until the region data runs out. Basically, we
output
 * the runs in the current region builder out as rectangles.
Using
 * x-coords from the region builder and y coords of the rows.
Then,
 * we add then next row to the region builder.
 */
dest_index = 0;
while (*rgn_data != R_EOR)
{
    ASSERT(*rgn_data == R_NEXT_IS_Y);

    in_run = FALSE;
    for (rgi = r_growth_list; rgi != NULL; rgi = rgi-
>rrgi_Next)
    {
        if ((rgi->rrgi_StateData &
RB_CUR_STATE_MASK) > 0)
        {
            /*
             * We have to emit
             * in one..
             */
            if (!in_run)
            {
a run here, if we're not already

```

```
prev_x = rgi->rrgi_RgnData;
```

```
in_run = TRUE;
```

```

    }
    else
    {
        if (in_run)
        {
            /*
            *
            *
            */
        }
    }

```

```

We've come to the end of a run. We output the rectangle
right here..

```

```
(!r_check_rect_buf_len(dest_index + 1))
```

```
return FALSE;
```

```
r_RectBuf[dest_index].X.Min = prev_x;
```

```
r_RectBuf[dest_index].Y.Min = prev_y;
```

```
r_RectBuf[dest_index].X.Max = rgi->rrgi_RgnData - 1;
```

```
r_RectBuf[dest_index++].Y.Max = rgn_data[1] - 1;
```

```

    }
    in_run = FALSE;
}

```

```

/*
 * Not efficient, get rid of it..
 */
if (rgi->rrgi_Next == NULL)
    rgi_tail = rgi;
}

```

```

/*
 * Now, we advance onto the next row..
 */
prev_y = rgn_data[1];
if (!R_add_row_to_region_growth_list(&rgn_data,
0x1, TRUE))
    return FALSE;
}
/*
 * Now, we've completed using the growth list for constructing
the
 * rect list. Therefore, we add it to the front of the free
list, to
 * be re-used later.
 */
rgi_tail->rrgi_Next = r_free_list;
r_free_list = r_growth_list;
r_growth_list = NULL;
/*
 * Ok, we have the array of rectangles sitting around. If
static_ok

```



```

        * is TRUE then we simply set the return pointers and get out.
        * Otherwise, we need to malloc a copy of the r_RectBuf.
        */
        *num_rects = dest_index;
        if (static_ok)
        {
            *rects = r_RectBuf;
        }
        else
        {
            *rects = (IntXYMinMax *)malloc(dest_index *
sizeof(IntXYMinMax));
            if (*rects == NULL)
                return FALSE;
            memcpy(*rects, r_RectBuf, dest_index *
sizeof(IntXYMinMax));
        }
        return TRUE;
    }
#endif /* R_USE_NEW_IMP */

/*
 * R_translate_region
 *
 * This function simply translates a region by the delta provided.
 *
 * Parameters:
 *     rgn      A ptr to the R_Region to be translated.
 *     delta    An IntXY ptr representing the amount to translate
 *              in x and y.
 *
 * Returns:
 *     Nothing.
 */
void
R_translate_region
(
    R_Region    *rgn,
    IntXY       *delta
)
{
    R_Int    *rgn_data;
    BB_translate(&rgn->rr_BBox, delta);
    rgn_data = rgn->rr_RgnData;
    for (int i = 0; i < rgn->rr_RgnDataSize - 1; i++)
    {
        if (rgn_data[i] == R_NEXT_IS_Y)
        {
            i++;
            rgn_data[i] += + delta->Y;
            continue;
        }
        rgn_data[i] += delta->X;
    }
}

/*
 * R_output_region_as_debug_string
 */

```

```

* This function simply outputs a region's data using the debug string
* functionality.
*
* Parameters:
*           rgn_name      A string used to output a user-
defined name for the
*           rgn           The
region to be output.
* Returns:
*           Nothing.
*/

```

```

void
R_output_region_as_debug_string
(
    char          *rgn_name,
    R_Region      *rgn
)
{
    char          buffer[128];
    int           index;
    int           line_len;

    sprintf(buffer, "\n+Rgn : %s\n", rgn_name);
    OutputDebugString(buffer);
    if (rgn == NULL)
    {
        sprintf(buffer, "+----End %s (Empty)\n",
rgn_name);
        OutputDebugString(buffer);
        return;
    }
    sprintf
    (
        buffer,
        "+----BBox:(%d, %d, %d, %d)\n",
        rgn->rr_BBox.X.Min,
        rgn->rr_BBox.Y.Min,
        rgn->rr_BBox.X.Max,
        rgn->rr_BBox.Y.Max
    );
    OutputDebugString(buffer);
    sprintf(buffer, "+----Nels: %d\n", rgn->rr_RgnDataSize);
    OutputDebugString(buffer);
    sprintf(buffer, "+----Data: ...");
    OutputDebugString(buffer);
    for (index = 0; index < rgn->rr_RgnDataSize; index++)
    {
        if (rgn->rr_RgnData[index] == R_NEXT_IS_Y)
        {
            sprintf(buffer, "\n|
Y:%3d--> ", rgn->rr_RgnData[++index]);
            line_len = strlen(buffer);
            OutputDebugString(buffer);
        }
        else if (rgn->rr_RgnData[index] == R_EOR)
        {

```

```

        %s\n", rgn_name);

        }
        else
        {
            sprintf(buffer, "%3d, ", rgn-
            >rr_RgnData[index]);
            if (strlen(buffer) + line_len >
            80)
            {
                OutputDebugString("\n|");
                line_len =
                strlen("\n|");
            }
            OutputDebugString(buffer);
            line_len += strlen(buffer);
        }
    }
}

#define NUM_ITERATIONS200

int
R_test_new_region_arithmetic()
{
    R_Region          rgn1;
    R_Region          rgn2;
    R_Region          rgn3;
    R_Region          rgn4;
    R_Region          rgn5;
    R_Region          rgn6;
    IntXYMinMax       rect;
    int               i;
    IntXY             delta;
    char              buf[256];
    unsigned long      ticks_new;
    unsigned long      ticks_old;

    #if 0
        /*
         * Union Test.
         */
        ticks_new = GetTickCount();
        rect.X.Min = 50;
        rect.Y.Min = 50;
        rect.X.Max = 100;
        rect.Y.Max = 100;
        if (!R_init_region_with_rect(&rgn1, &rect))
            return FALSE;
        rect.X.Min = 70;
        rect.Y.Min = 70;
        rect.X.Max = 120;
        rect.Y.Max = 120;
        if (!R_init_region_with_rect(&rgn2, &rect))
            return FALSE;
        if (!R_union_list_equals(&rgn1, &rgn2))
    
```

```

        return FALSE;
    delta.X = 5;
    delta.Y = 5;
    for (i = 0; i < NUM_ITERATIONS; i++)
    {
        R_translate_region(&rgn2, &delta);
        if (!R_union_list_equals(&rgn1, &rgn2))
            return FALSE;
    }
    ticks_new = GetTickCount() - ticks_new;

    ticks_old = GetTickCount();
    rect.X.Min = 50;
    rect.Y.Min = 50;
    rect.X.Max = 100;
    rect.Y.Max = 100;
    if (!R_init_region_with_rect(&rgn3, &rect))
        return FALSE;
    rect.X.Min = 70;
    rect.Y.Min = 70;
    rect.X.Max = 120;
    rect.Y.Max = 120;
    if (!R_init_region_with_rect(&rgn4, &rect))
        return FALSE;
    if (!R_union_equals(&rgn3, &rgn4))
        return FALSE;
    delta.X = 5;
    delta.Y = 5;
    for (i = 0; i < NUM_ITERATIONS; i++)
    {
        R_translate_region(&rgn4, &delta);
        if (!R_union_equals(&rgn3, &rgn4))
            return FALSE;
    }
    ticks_old = GetTickCount() - ticks_old;

    if (R_compare(&rgn1, &rgn3))
        sprintf(buf, "New & Old Region Implementations
match.\n");
    else
        sprintf(buf, "New & Old Region Implementations
DO NOT match.\n");
    OutputDebugString(buf);
    sprintf(buf, "Union Timings - New=%d vs Old=%d\n", ticks_new,
ticks_old);
    OutputDebugString(buf);
    //R_output_region_as_debug_string("New Region Description",
&rgn1);
    //R_output_region_as_debug_string("Old Region Description",
&rgn3);

    /*
     * Intersection Test.
     */
    R_empty_region(&rgn2);
    R_empty_region(&rgn4);
    rect.X.Min = 70;
    rect.Y.Min = 70;

```

```

rect.X.Max = 120;
rect.Y.Max = 120;
if (!R_init_region_with_rect(&rgn2, &rect))
    return FALSE;
delta.X = 5;
delta.Y = 5;
for (i = 0; i < NUM_ITERATIONS; i++)
{
    if (!R_intersection_list(&rgn5, &rgn1, &rgn2))
        return FALSE;
    if (!R_intersection(&rgn6, &rgn3, &rgn2))
        return FALSE;
    if (!R_compare(&rgn5, &rgn6))
    {
        sprintf(buf, "New & Old Region
Implementations DO NOT match.\n");
        OutputDebugString(buf);
    }
    R_empty_region(&rgn5);
    R_empty_region(&rgn6);
    R_translate_region(&rgn2, &delta);
}

ticks_new = GetTickCount();
R_empty_region(&rgn2);
rect.X.Min = 70;
rect.Y.Min = 70;
rect.X.Max = 120;
rect.Y.Max = 120;
if (!R_init_region_with_rect(&rgn2, &rect))
    return FALSE;
delta.X = 5;
delta.Y = 5;
for (i = 0; i < NUM_ITERATIONS; i++)
{
    if (!R_intersection_list(&rgn5, &rgn1, &rgn2))
        return FALSE;
    R_empty_region(&rgn5);
    R_translate_region(&rgn2, &delta);
}
ticks_new = GetTickCount() - ticks_new;

ticks_old = GetTickCount();
R_empty_region(&rgn2);
rect.X.Min = 70;
rect.Y.Min = 70;
rect.X.Max = 120;
rect.Y.Max = 120;
if (!R_init_region_with_rect(&rgn2, &rect))
    return FALSE;
delta.X = 5;
delta.Y = 5;
for (i = 0; i < NUM_ITERATIONS; i++)
{
    if (!R_intersection(&rgn6, &rgn3, &rgn2))
        return FALSE;
    R_empty_region(&rgn6);
    R_translate_region(&rgn2, &delta);
}

```

```

    }
    ticks_old = GetTickCount() - ticks_old;

    sprintf(buf, "Intersection Timings - New=%d vs Old=%d\n",
ticks_new, ticks_old);
    OutputDebugString(buf);
    //R_output_region_as_debug_string("New Region Description",
&rgn1);
    //R_output_region_as_debug_string("Old Region Description",
&rgn3);

    R_empty_region(&rgn1);
    R_empty_region(&rgn2);
    R_empty_region(&rgn3);
    R_empty_region(&rgn4);
    OutputDebugString("Done!!\n");
#endif
    return TRUE;
}

```

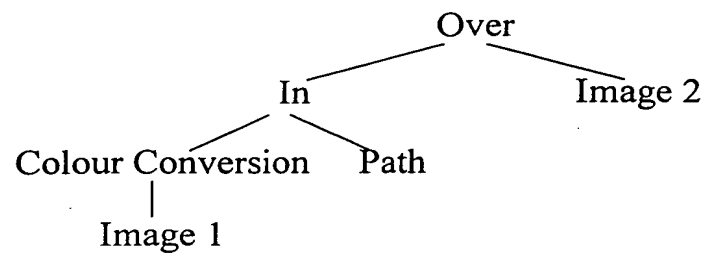


Figure 1.

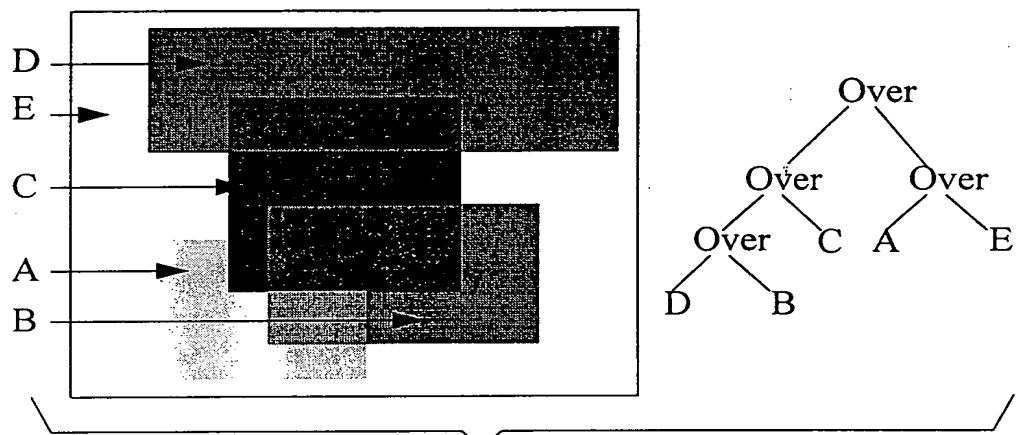


Figure 2.

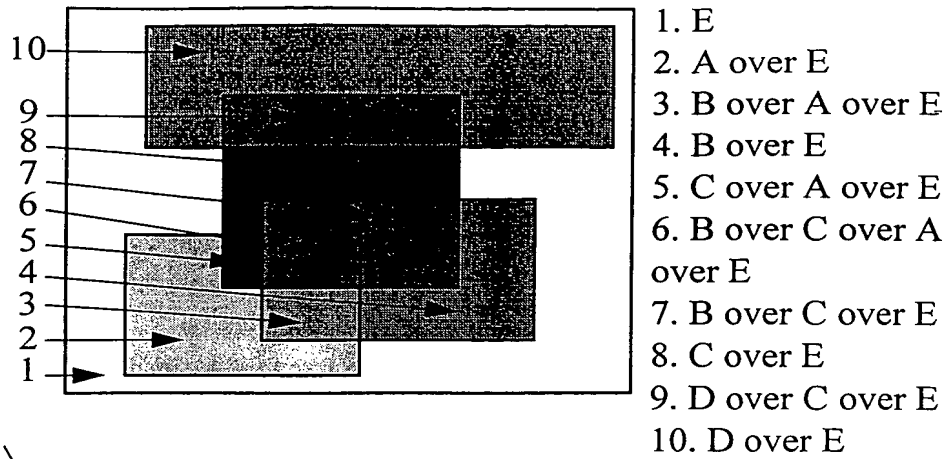


Figure 3.

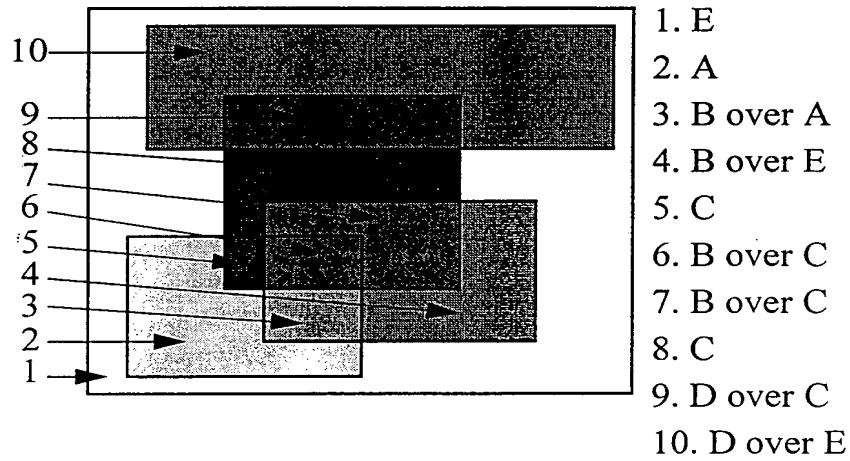


Figure 4.

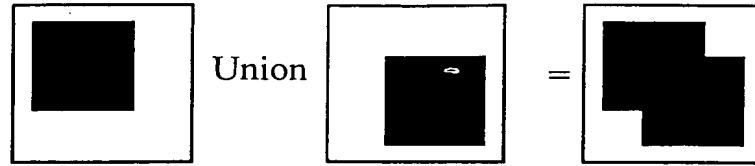


Figure 5.

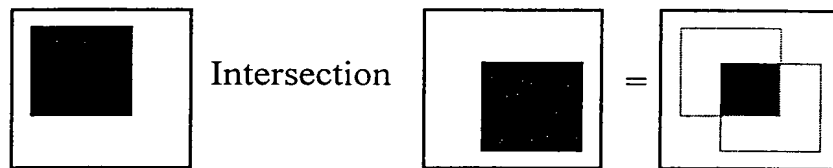


Figure 6.

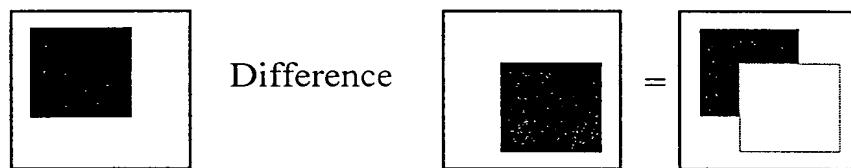
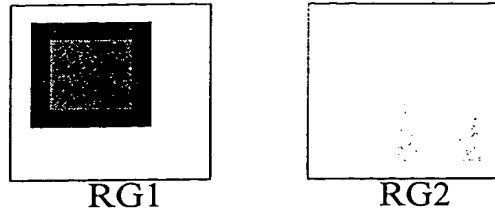
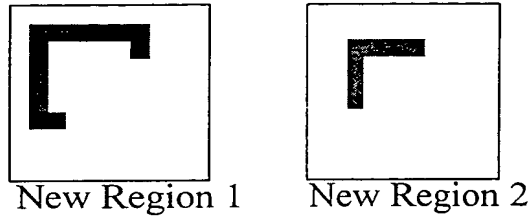
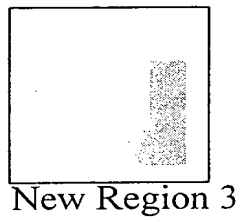
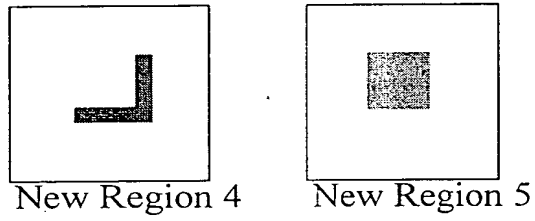
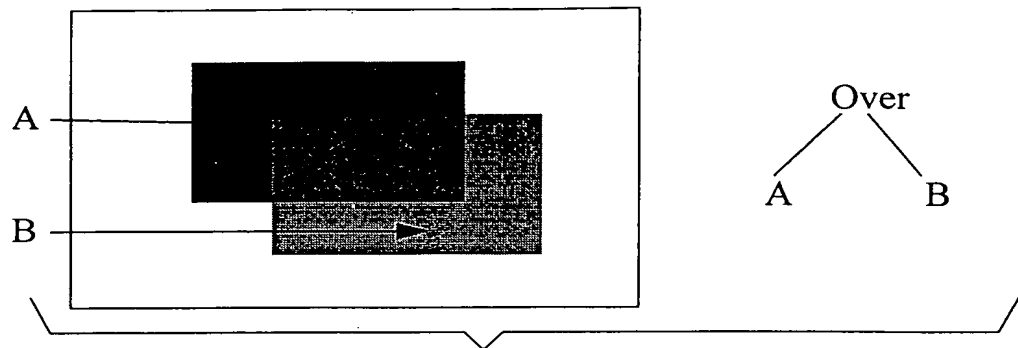
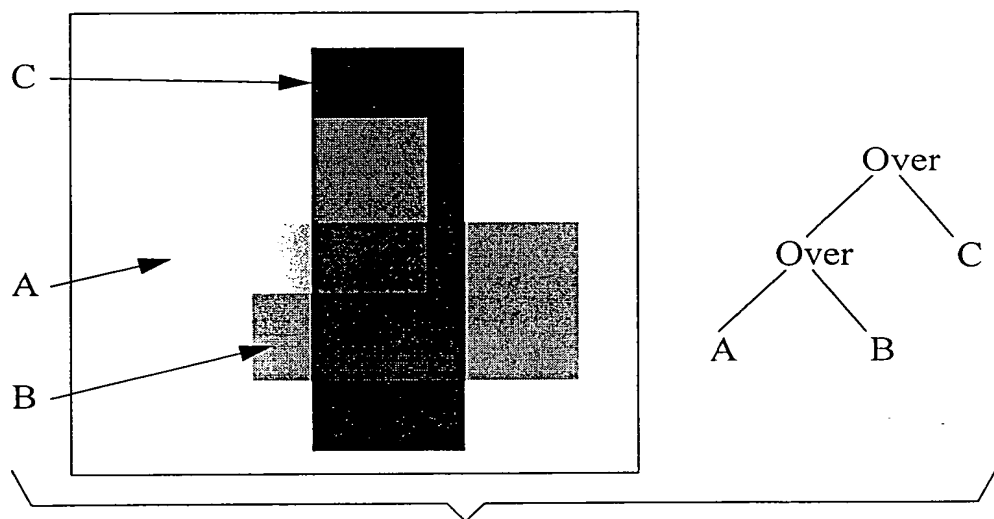


Figure 7.

Fig. 8A**Fig. 8B****Fig. 8C****Fig. 8D****Figure 8.**

**Figure 9.****Figure 10.**

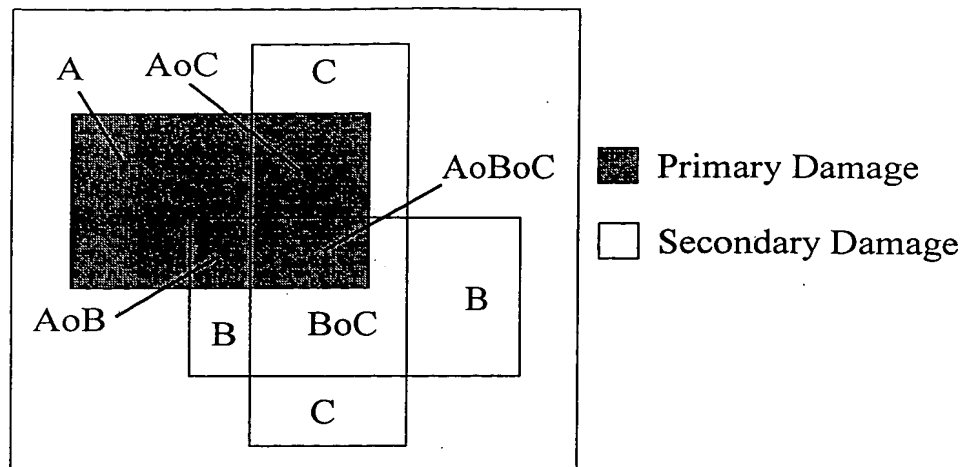


Figure 11.

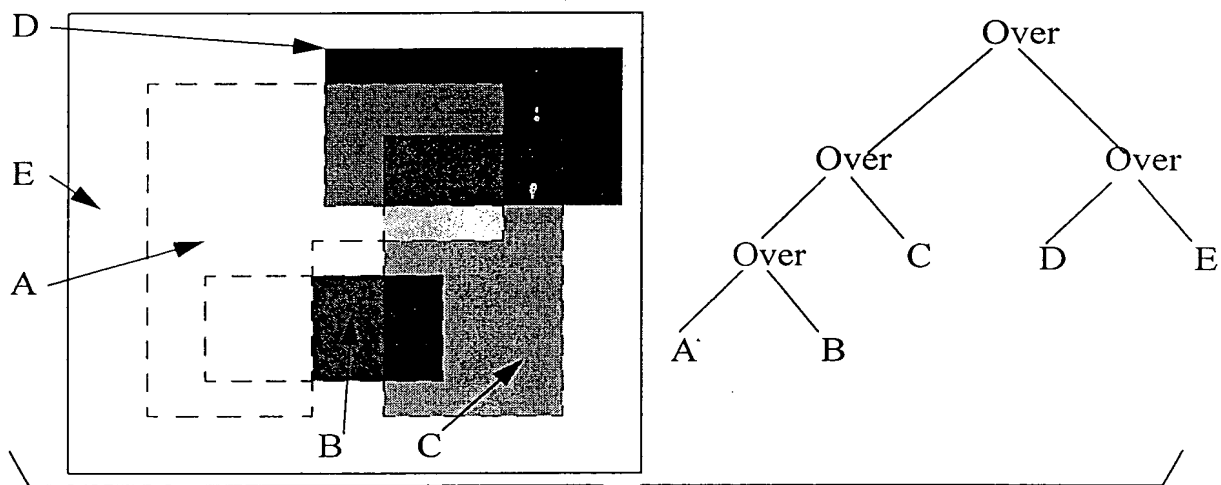
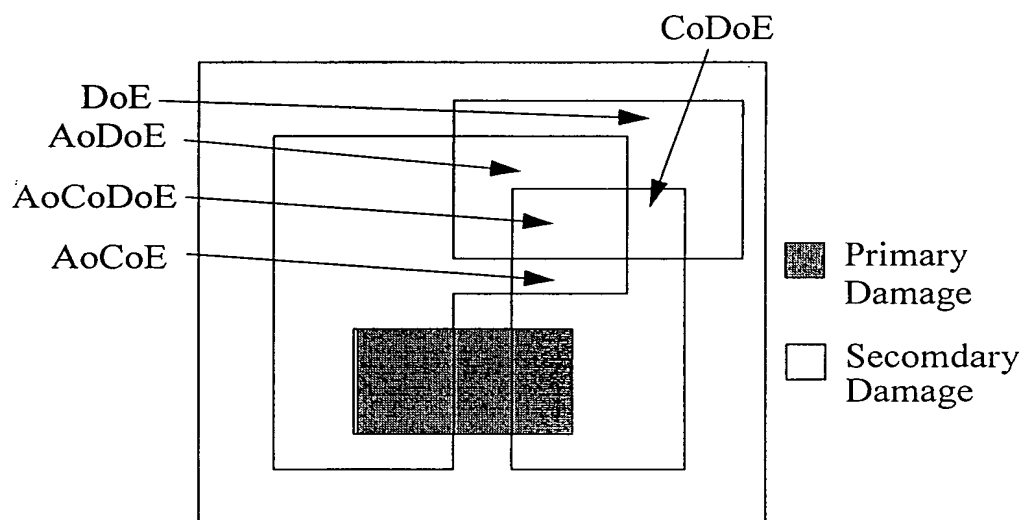
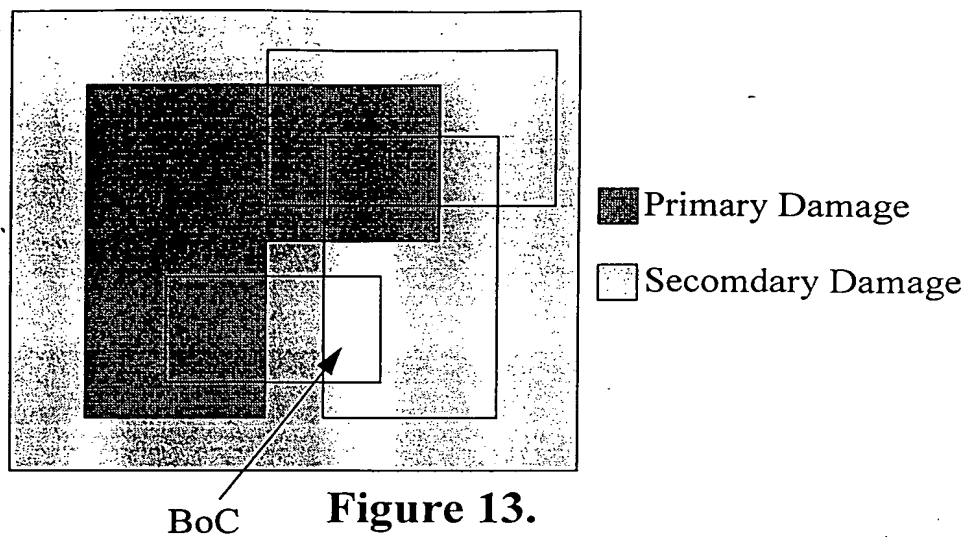


Figure 12.



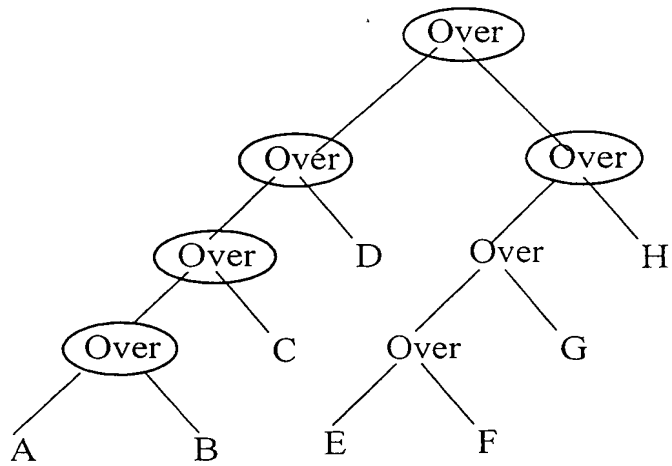
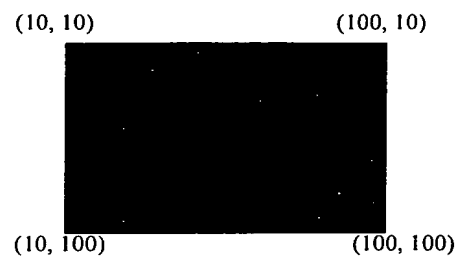
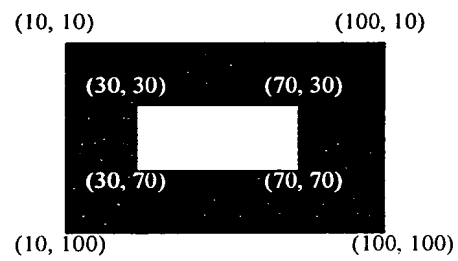


Figure 15.

**Figure. 16****Figure. 17**

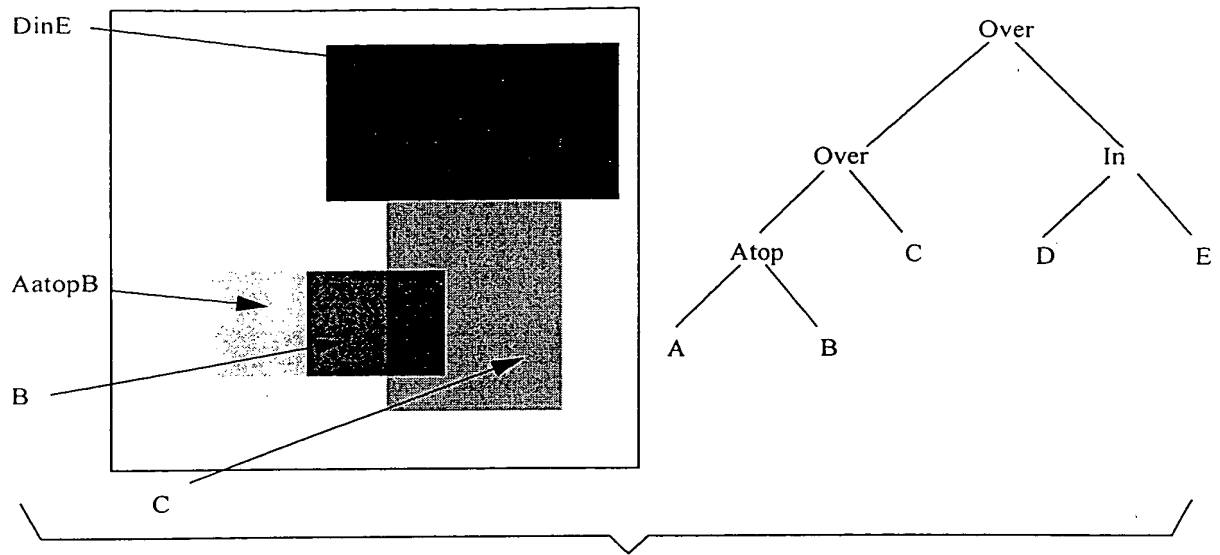
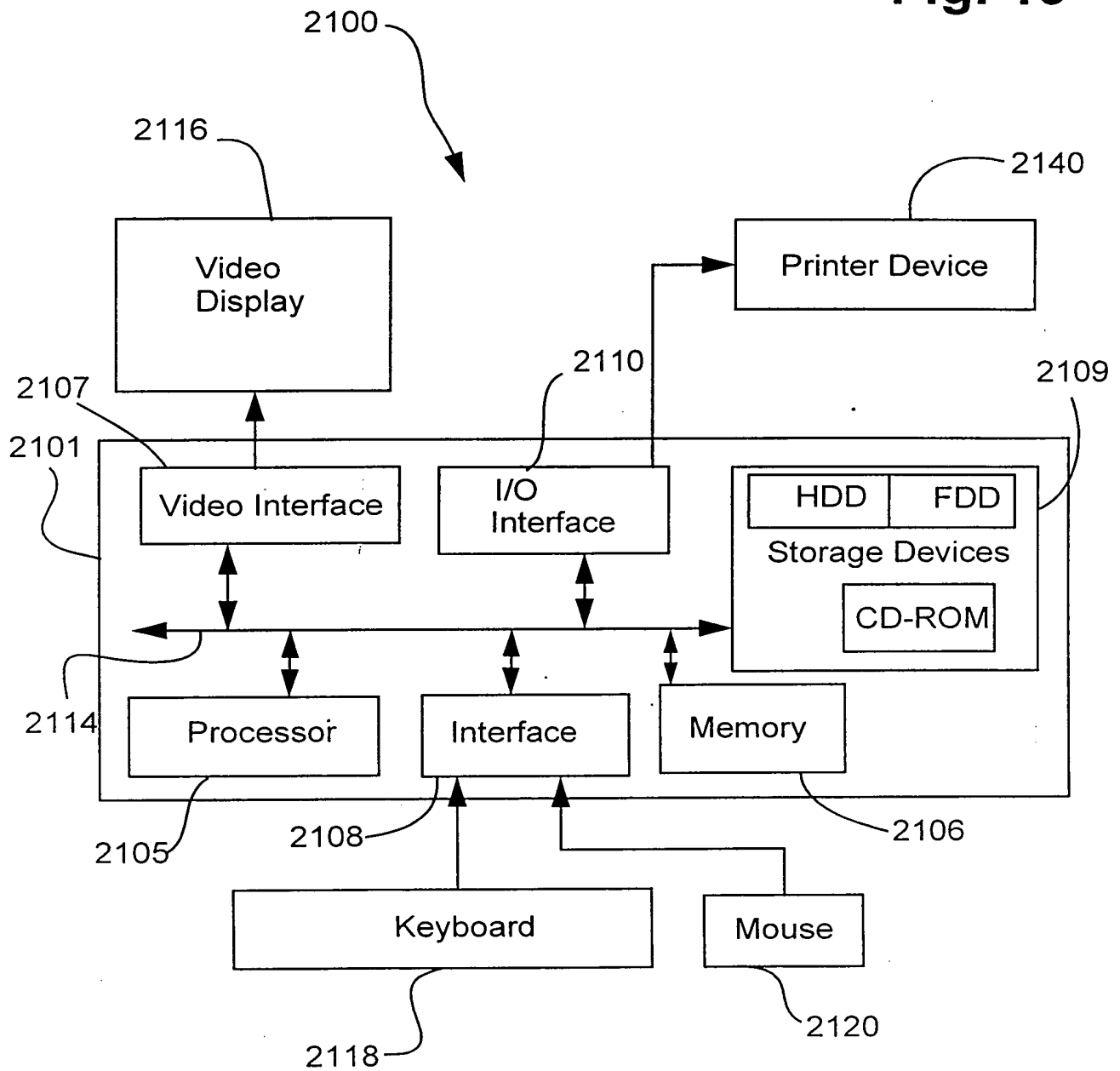
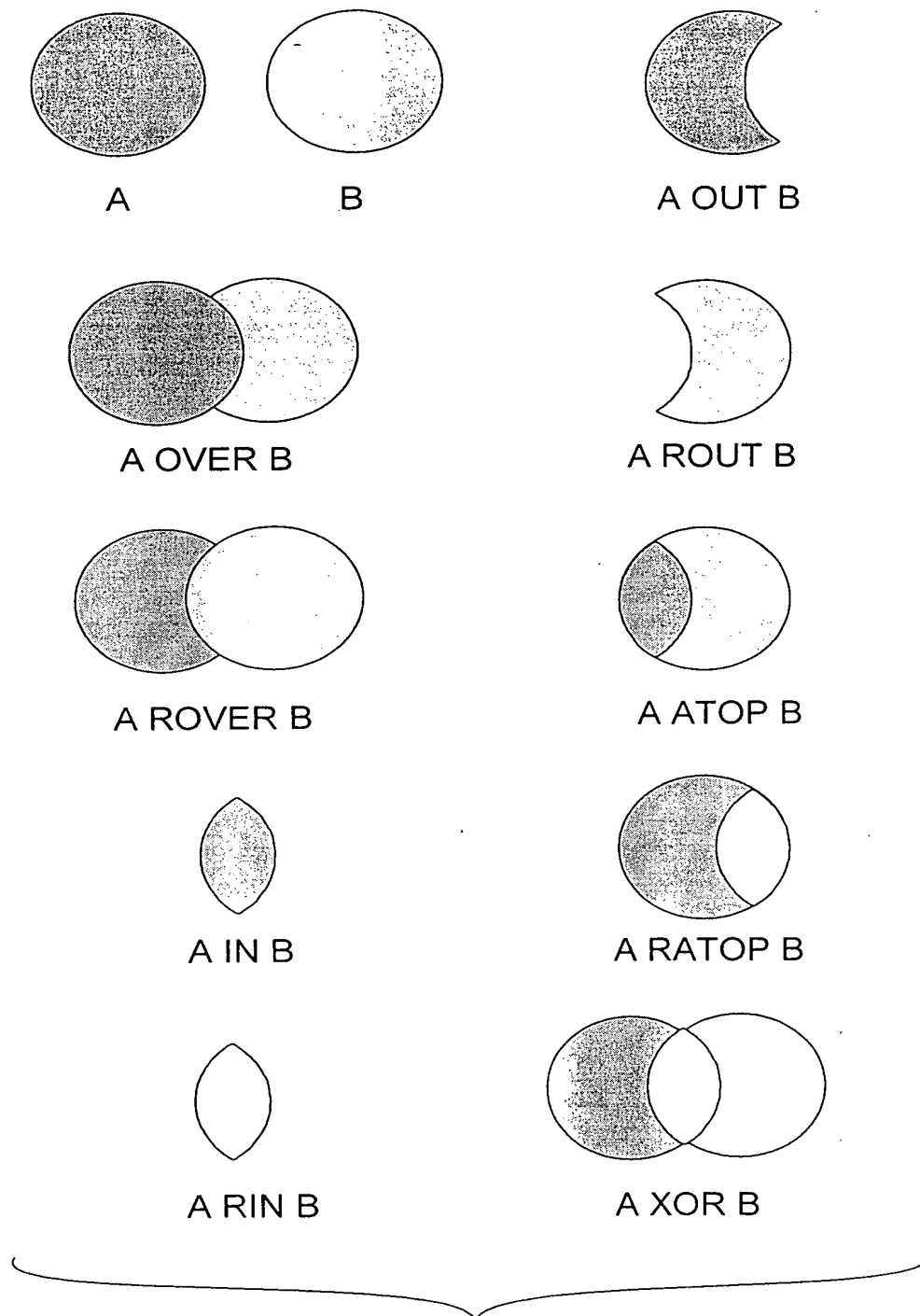
**Figure. 18**

Fig. 19

**Fig. 20**

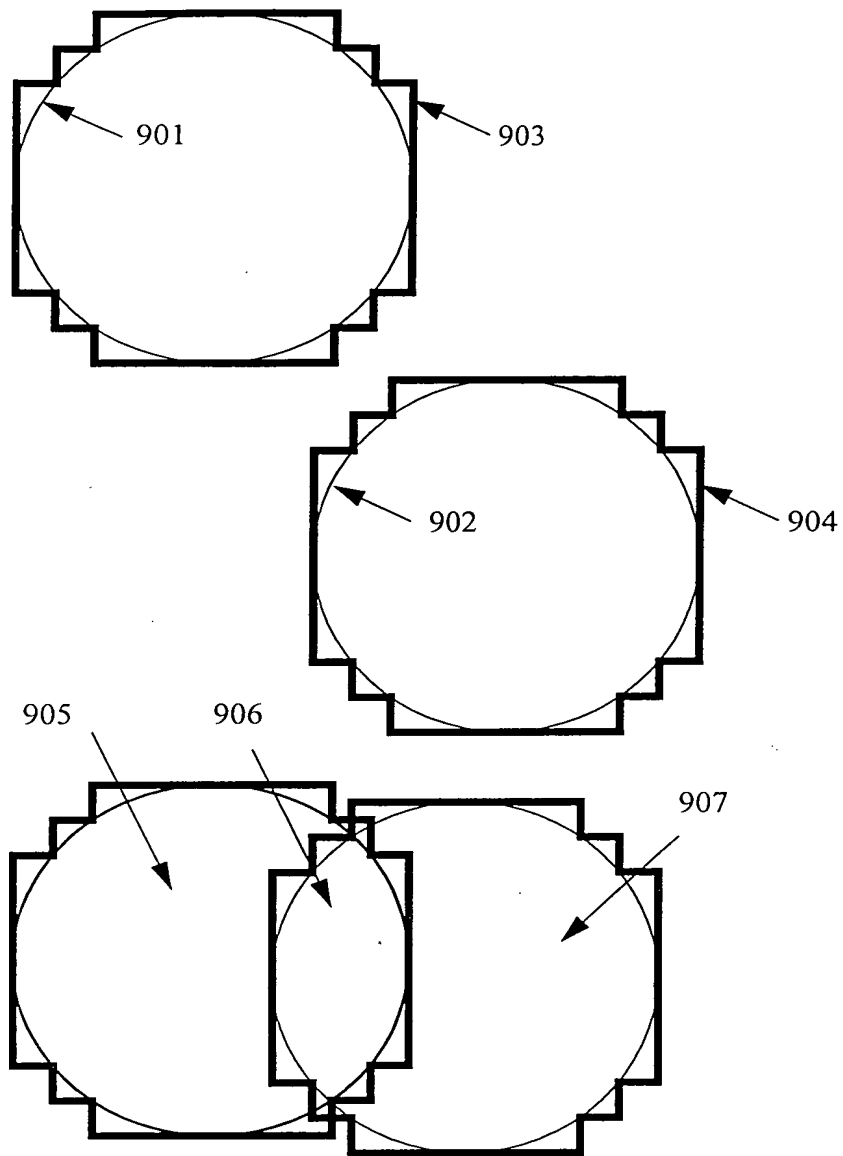


Figure 21

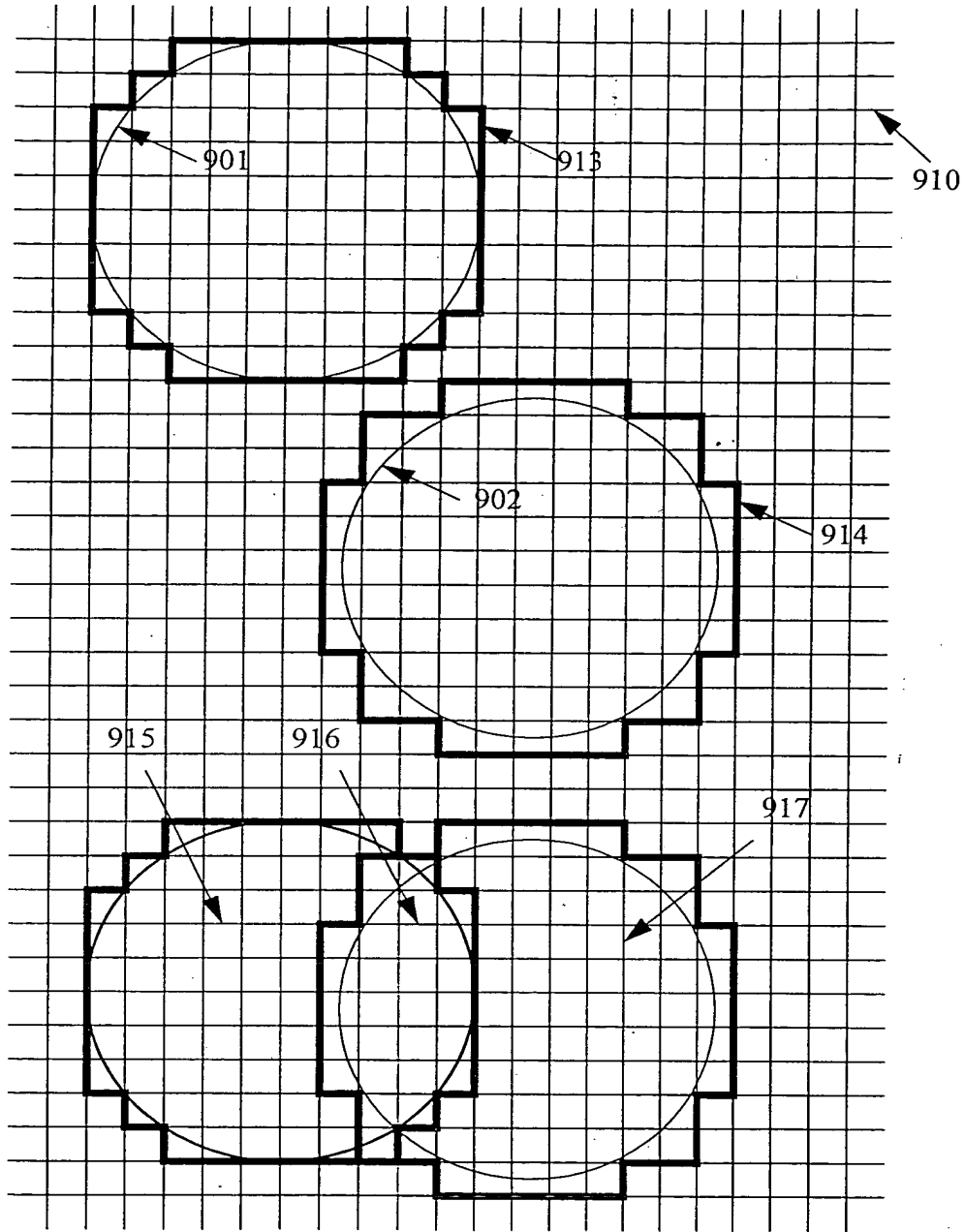


Figure 22